

Discrete Optimization

Lecture 2

Part 1

Minimum Spanning Tree

Slides online: <https://project.inria.fr/2015ma2827>

Slides courtesy of M. Pawan Kumar

Recap of previous class

- Basics of Graphs (directed, undirected)
 - Walks
 - Paths
 - Circuits
- Shortest Path Algorithms
 - 4 of them
- Assignment 1 given
 - Register team. Deadline: 7/4 !

Outline

- Chow-Liu Tree
- Minimum Spanning Tree Problem
- Kruskal's Method
- Prim's Method

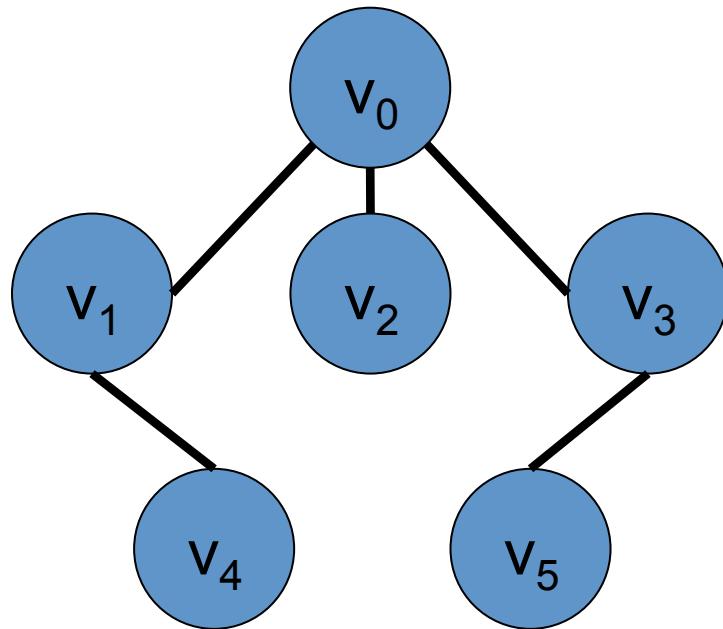
Distribution

Eg: Pose estimation – Estimate joint probability of body parts



$$P(x) = P(x_0, x_1, \dots, x_{n-1}) = P(x_0)P(x_1|x_0)\dots P(x_{n-1}|x_1, \dots, x_{n-2})$$

Known Tree Structure



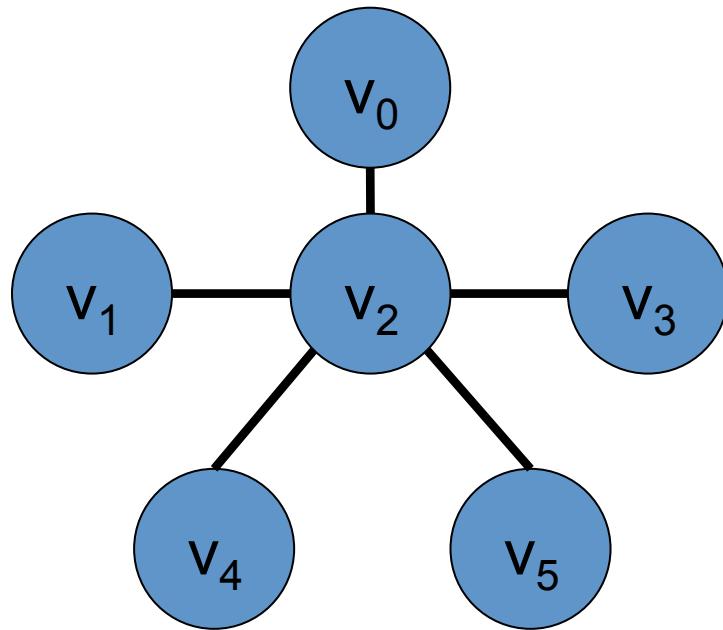
Distribution $P_T(x)$

$v_{p(a)}$ = “parent” of v_a

$$P_T(x_5|x_3)P_T(x_4|x_1)P_T(x_3|x_0)P_T(x_2|x_0)P_T(x_1|x_0)P_T(x_0)$$

Estimate $P_T(x_a|x_{p(a)}) = P(x_a|x_{p(a)})$

Known Tree Structure



Distribution $P_T(x)$

$v_{p(a)}$ = “parent” of v_a

$$P_T(x_5|x_2)P_T(x_4|x_2)P_T(x_3|x_2)P_T(x_2|x_0)P_T(x_1|x_2)P_T(x_0)$$

Estimate $P_T(x_a|x_{p(a)}) = P(x_a|x_{p(a)})$ **Which tree?**

Kullback-Leibler Divergence

$$\text{KL}(P_1 \parallel P_2) = - \sum_x P_1(x) \log P_2(x) + \sum_x P_1(x) \log P_1(x)$$

Constant

$$KL(P_1 \parallel P_2) \geq 0$$

$$KL(P_1 \parallel P_1) = 0$$

Substitute $P_1 = P$ and $P_2 = P_T$. Minimize $KL(P \parallel P_T)$

Estimating the Tree Structure

$$\min - \sum_x P(x) \log P_T(x)$$

Estimating the Tree Structure

$$\min - \sum_x P(x) \log \prod_a P_T(x_a | x_{p(a)})$$

Estimating the Tree Structure

$$\min - \sum_x P(x) \sum_a \log P_T(x_a | x_{p(a)})$$

Estimating the Tree Structure

$$\min - \sum_x P(x) \sum_a \log \frac{P_T(x_a, x_{p(a)})}{P_T(x_{p(a)})}$$

Estimating the Tree Structure

$$\min - \sum_x P(x) \sum_a \log \frac{P_T(x_a, x_{p(a)}) P(x_a)}{P_T(x_{p(a)}) P(x_a)}$$

Estimating the Tree Structure

$$\min - \sum_x P(x) \sum_a \log \frac{P_T(x_a, x_{p(a)})}{P_T(x_{p(a)}) P(x_a)}$$

$$-\sum_x P(x) \sum_a \log P(x_a)$$

Independent of the tree structure

Estimating the Tree Structure

$$\min - \sum_a \sum_{x_a} \sum_{x_{p(a)}} P(x_a, x_{p(a)}) \log \frac{P_T(x_a, x_{p(a)})}{P_T(x_{p(a)}) P(x_a)}$$

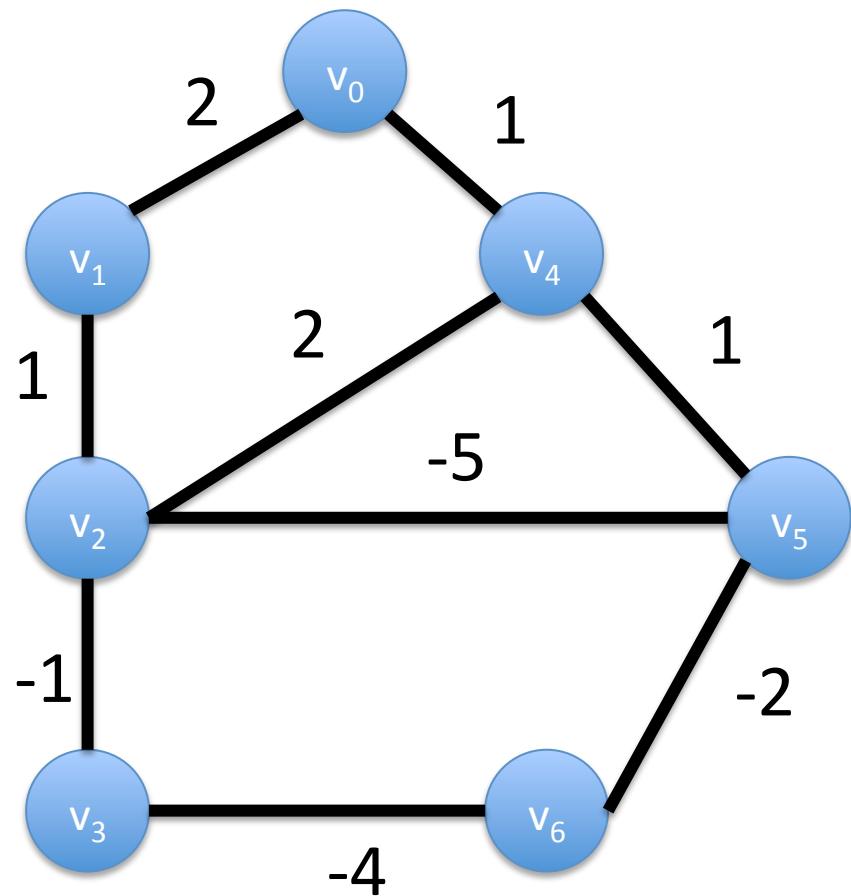
$$\min - \sum_a I(x_a, x_{p(a)})$$

Mutual Information

Outline

- Chow-Liu Tree
- **Minimum Spanning Tree Problem**
- Kruskal's Method
- Prim's Method

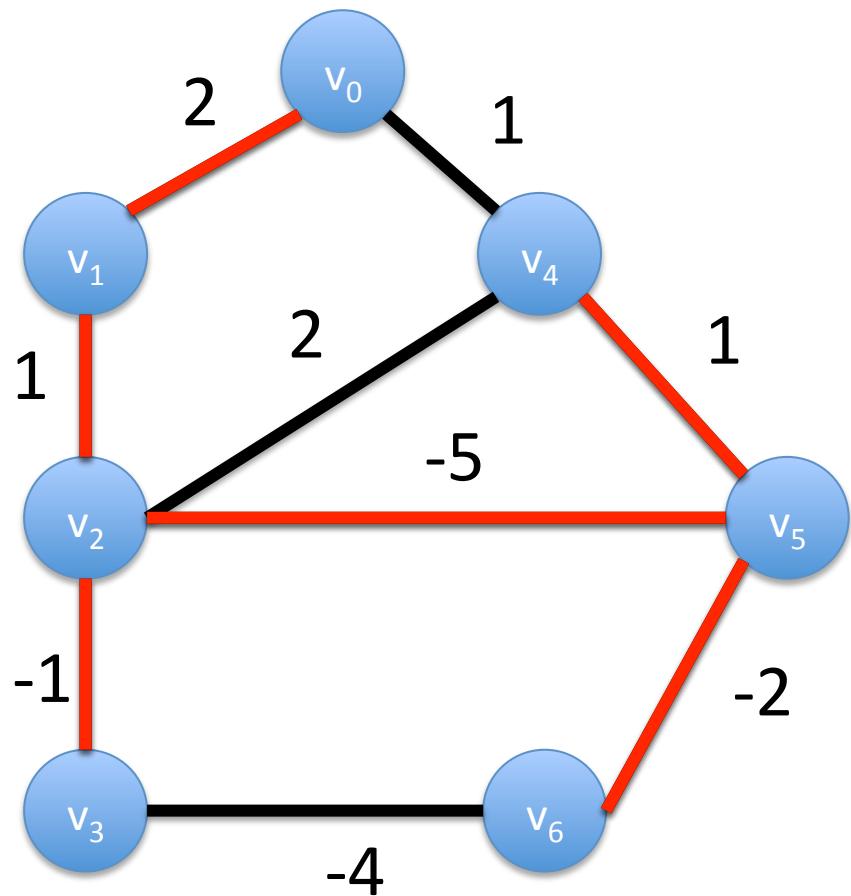
Undirected Connected Simple Graph



$$G = (V, E)$$

$$l: E \rightarrow \mathbb{R}$$

Spanning Tree



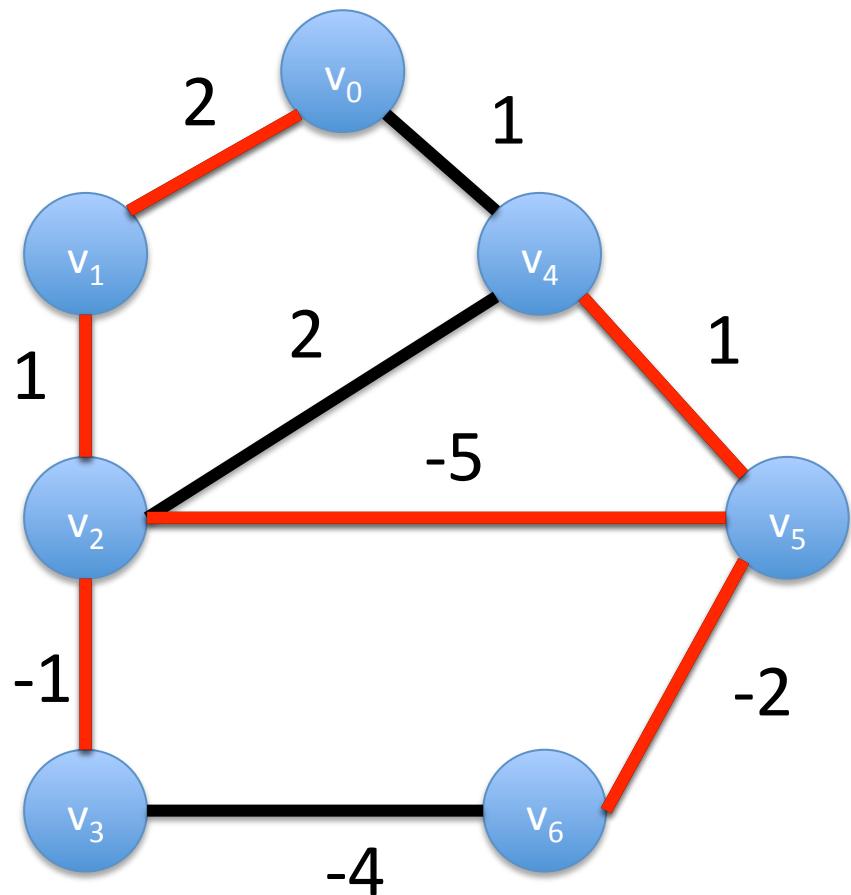
$$G = (V, E)$$

$$T = (V, E_T)$$

E_T is a subset of E

Graph T is a tree

Weight of a Spanning Tree



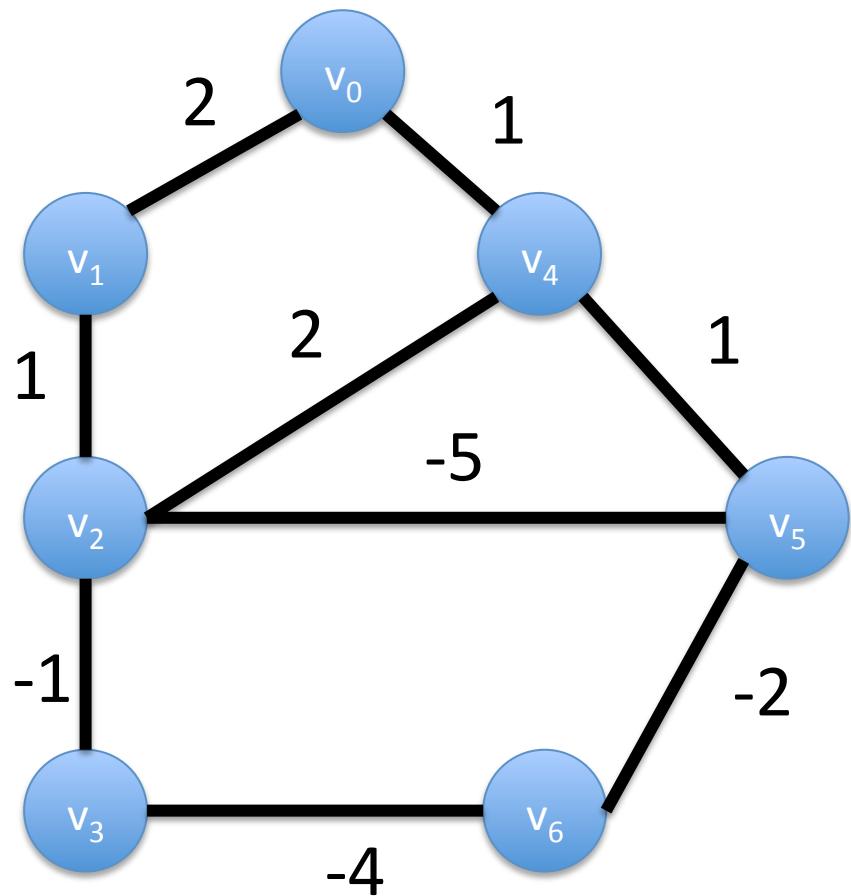
$$G = (V, E)$$

$$T = (V, E_T)$$

$w(T) = \text{Sum of the length of all edges in } E_T$

$$w(T) = 2 + 1 - 1 - 5 + 1 - 2 = -4$$

Minimum Spanning Tree Problem



$$G = (V, E)$$

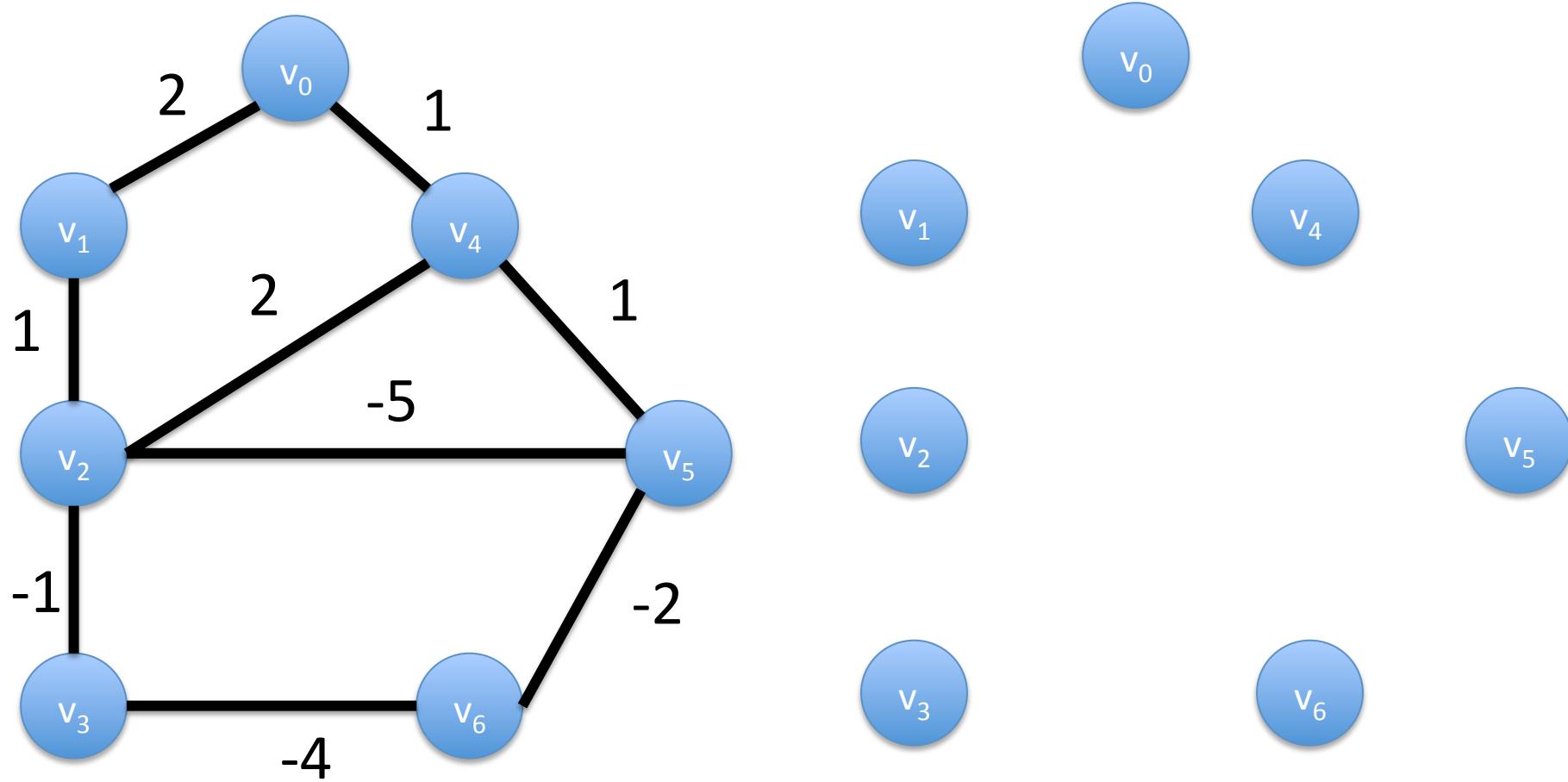
$$T^* = \operatorname{argmin}_T w(T)$$

Find a tree with the minimum weight

Outline

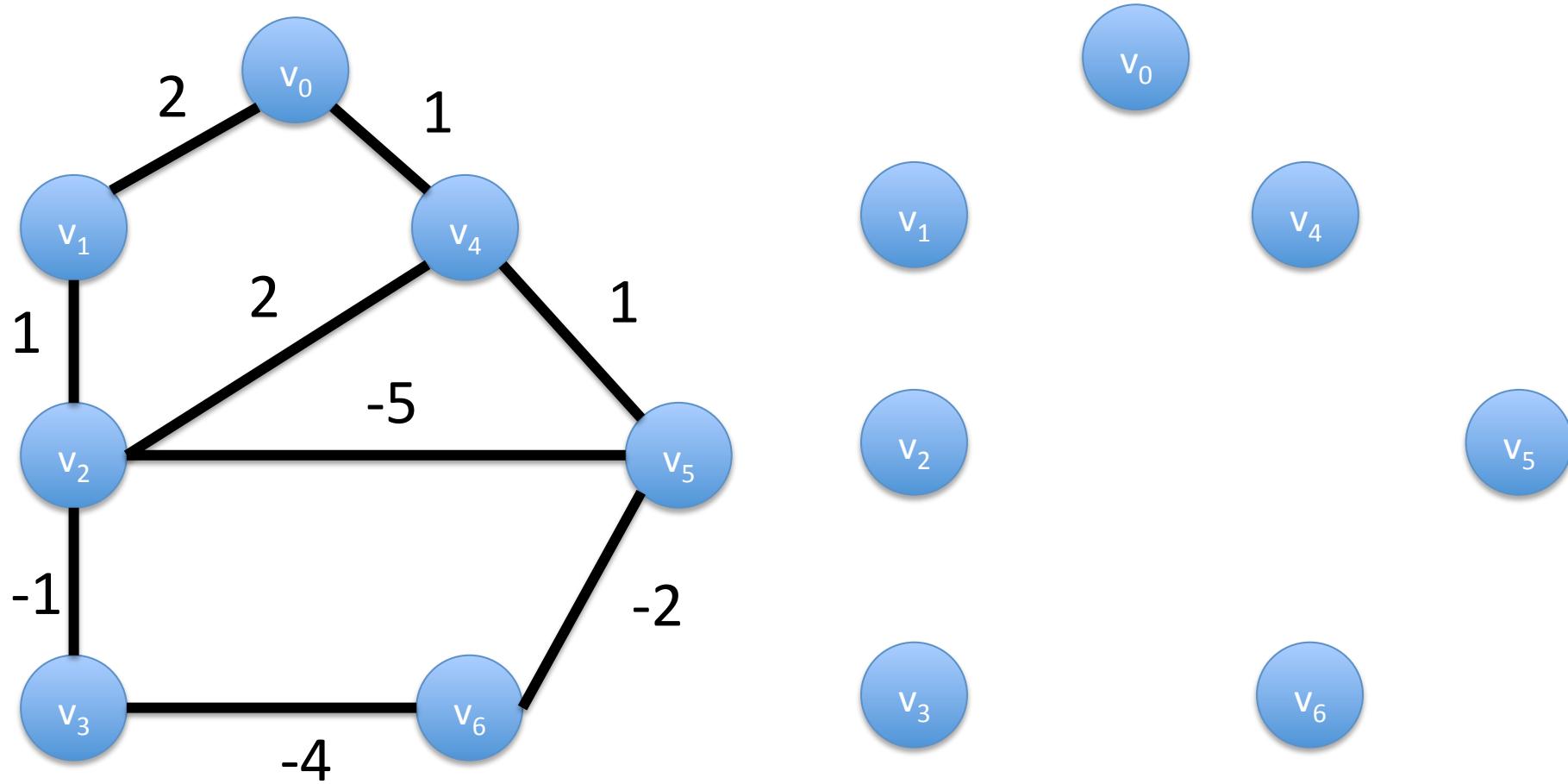
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Kruskal's Method



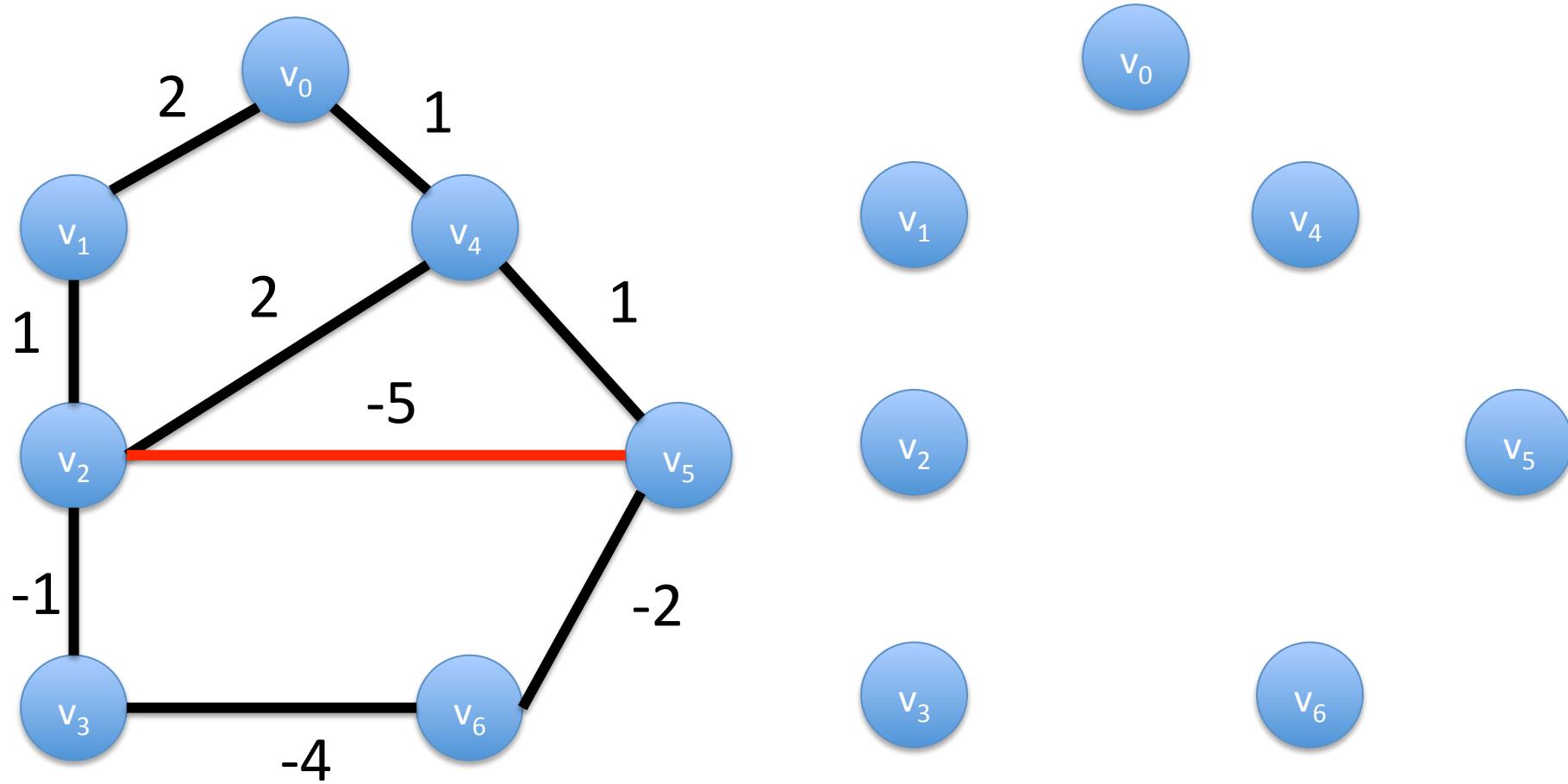
Start with a forest where every vertex is a tree

Kruskal's Method



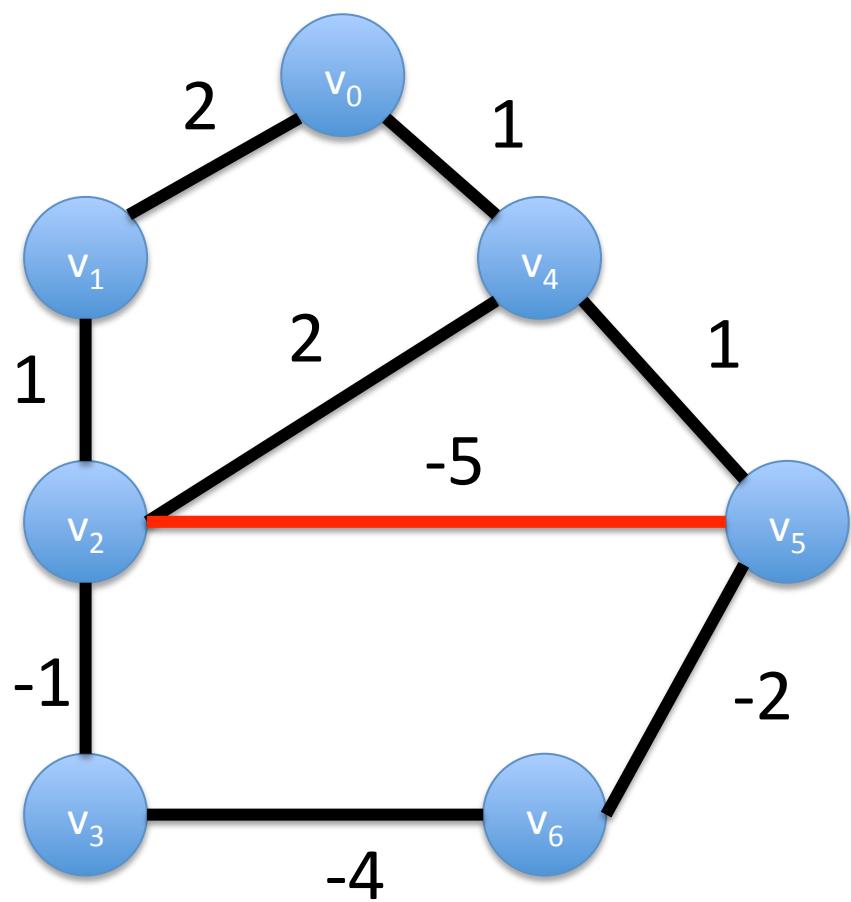
Select the edge with the minimum length

Kruskal's Method



Select the edge with the minimum length

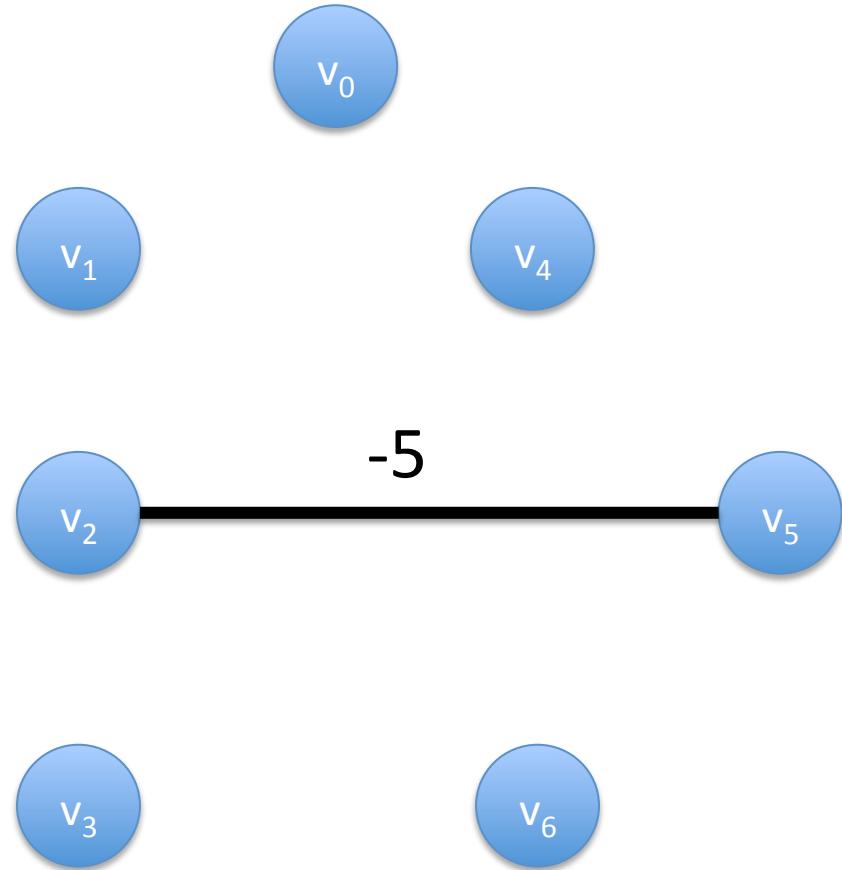
Kruskal's Method



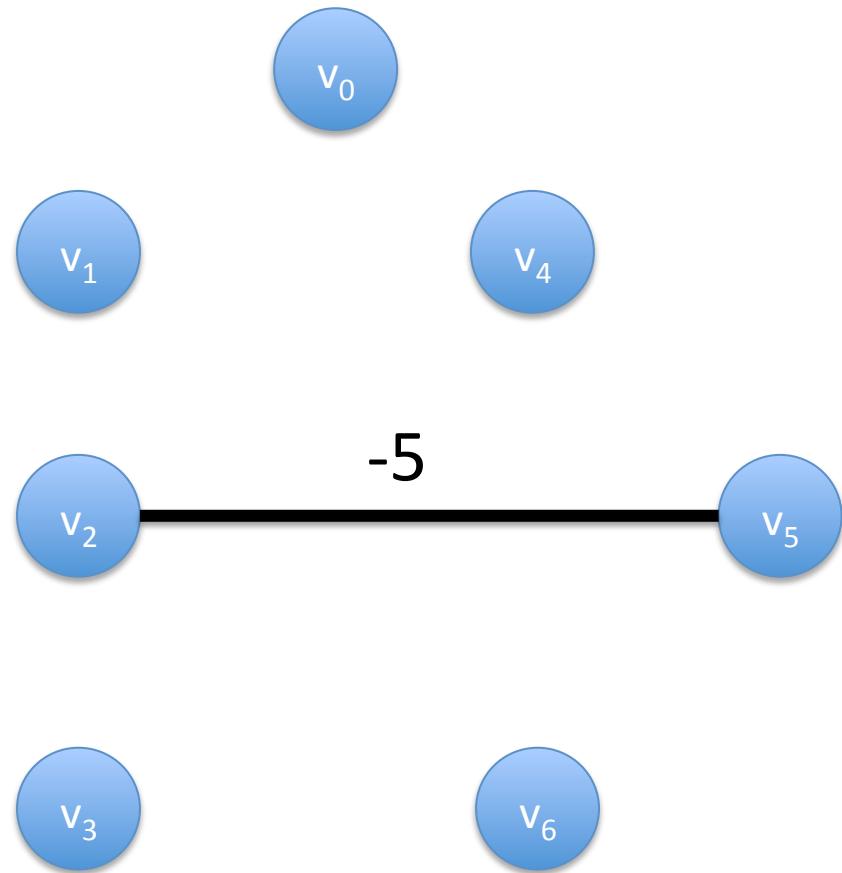
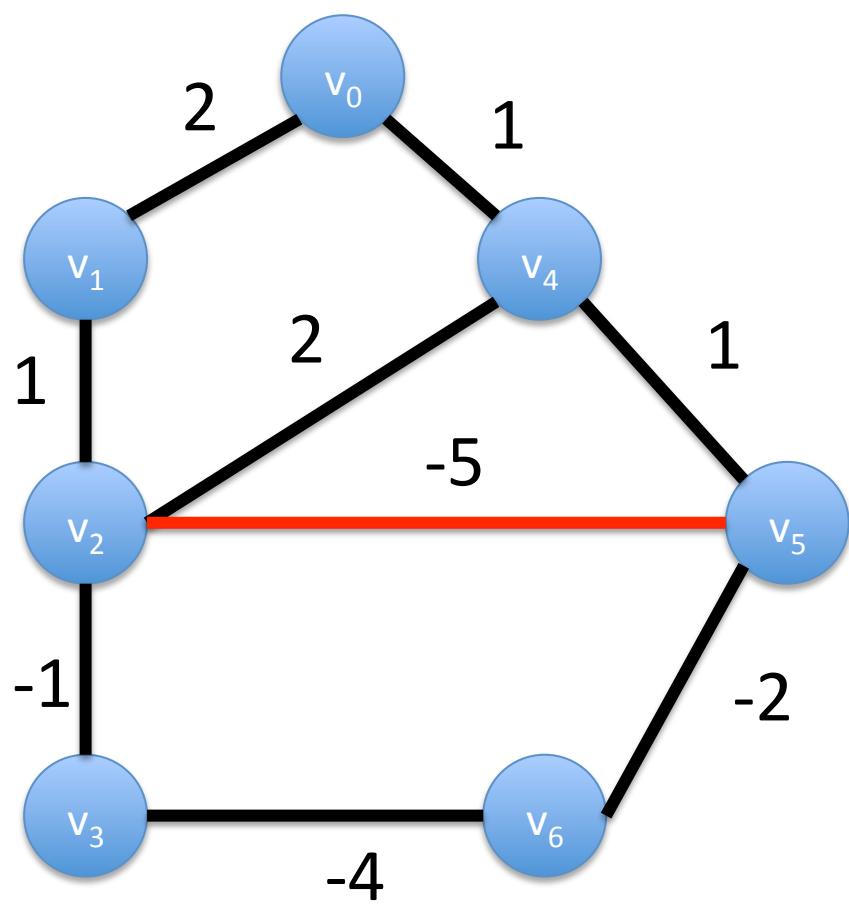
YES

Add the edge to the forest

Does this edge connect two different trees?

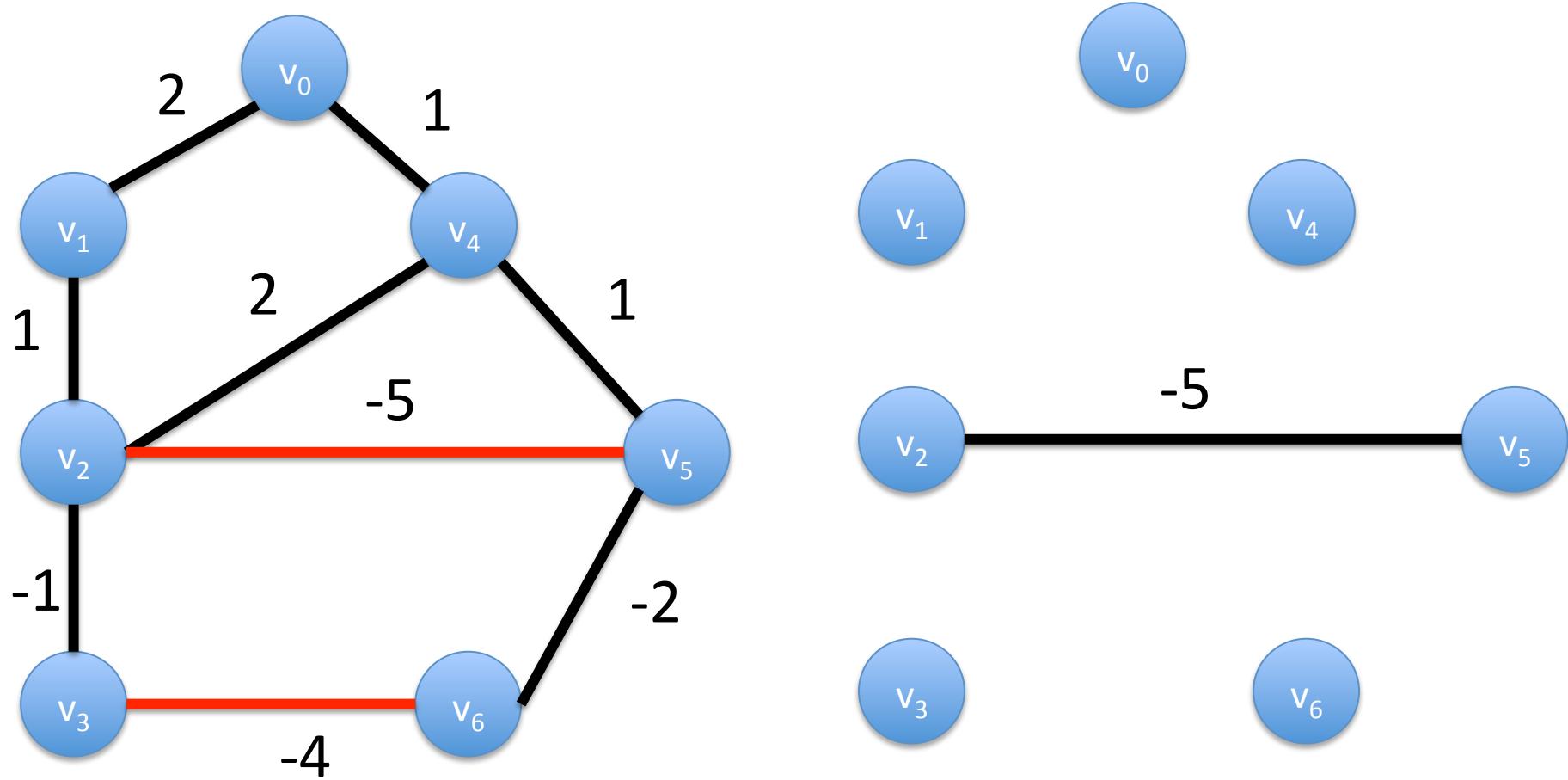


Kruskal's Method



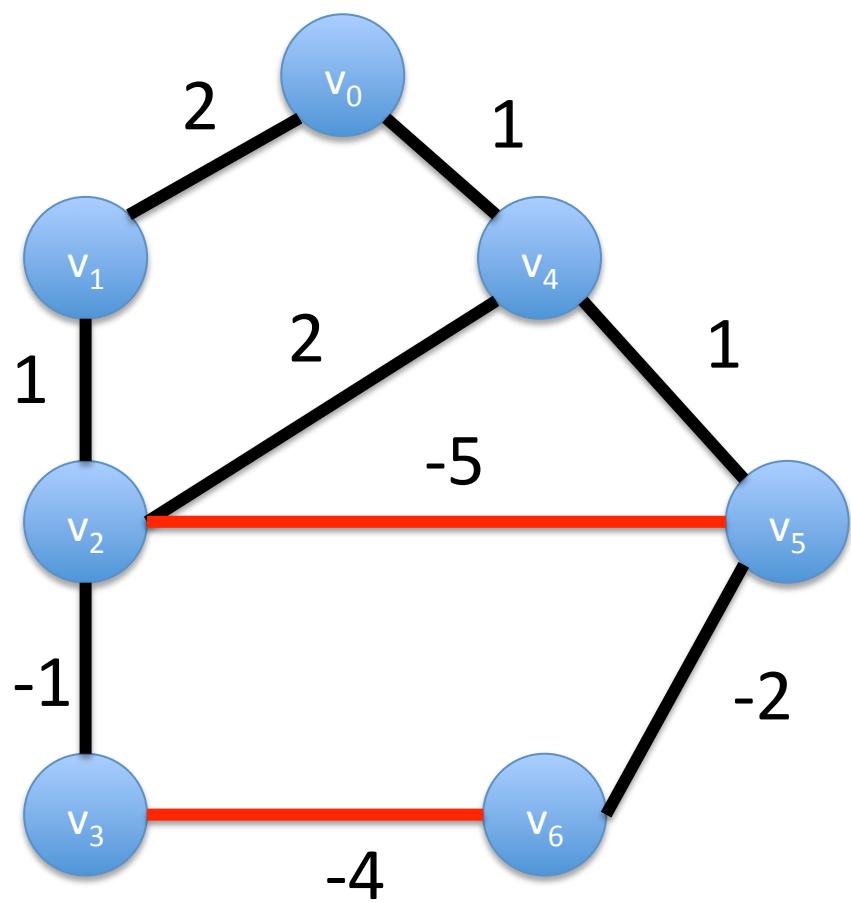
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Kruskal's Method



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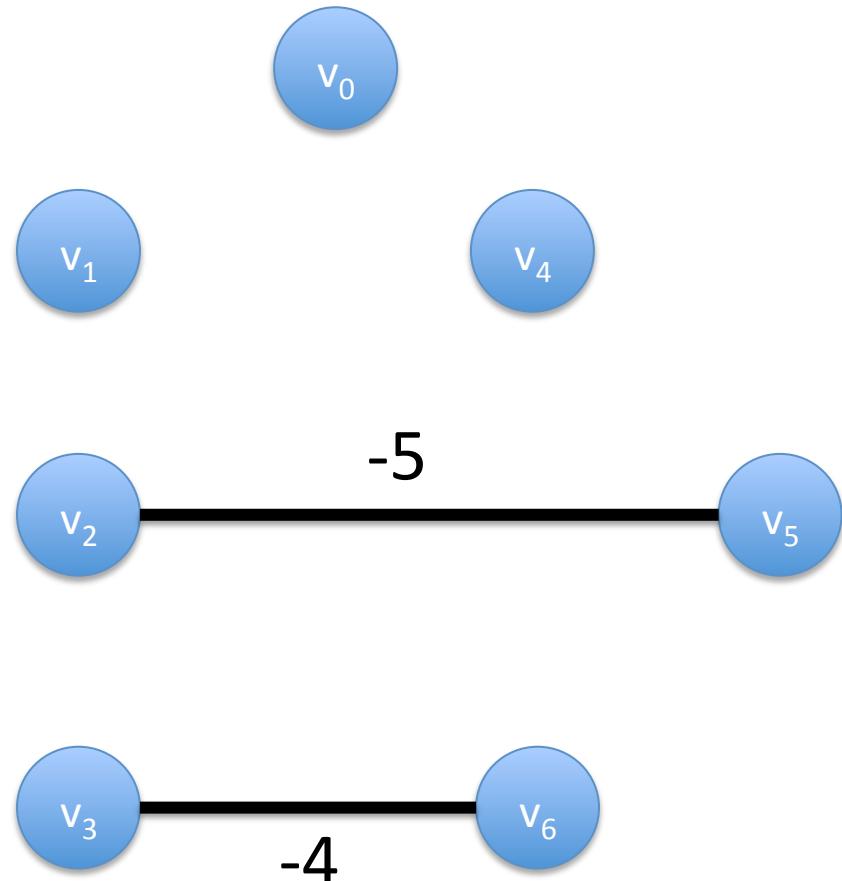
Kruskal's Method



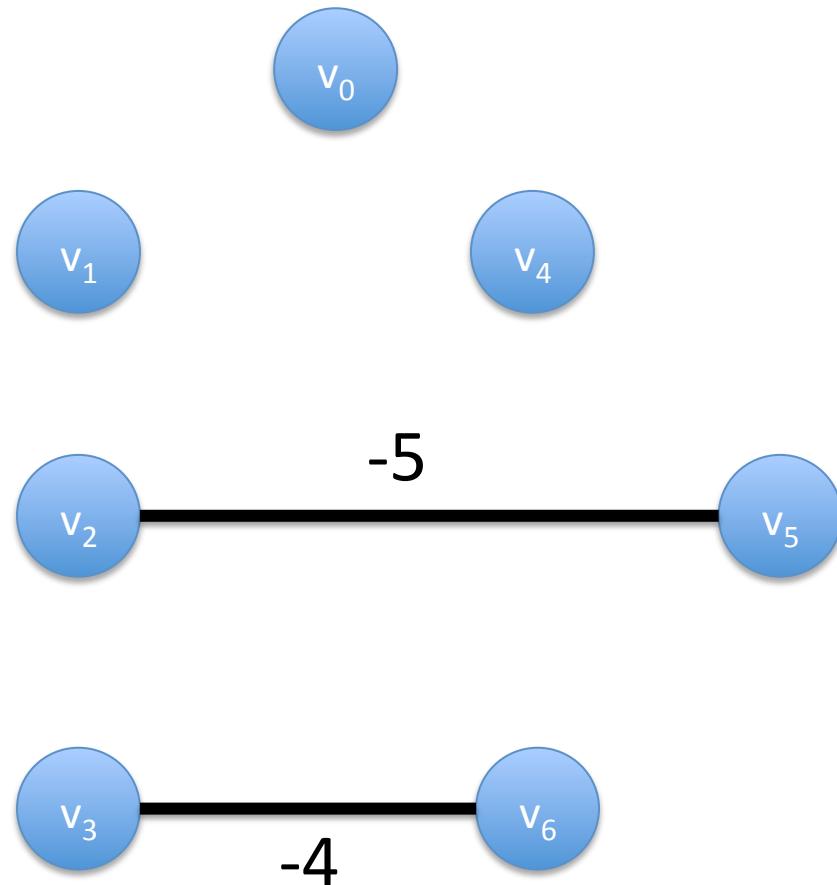
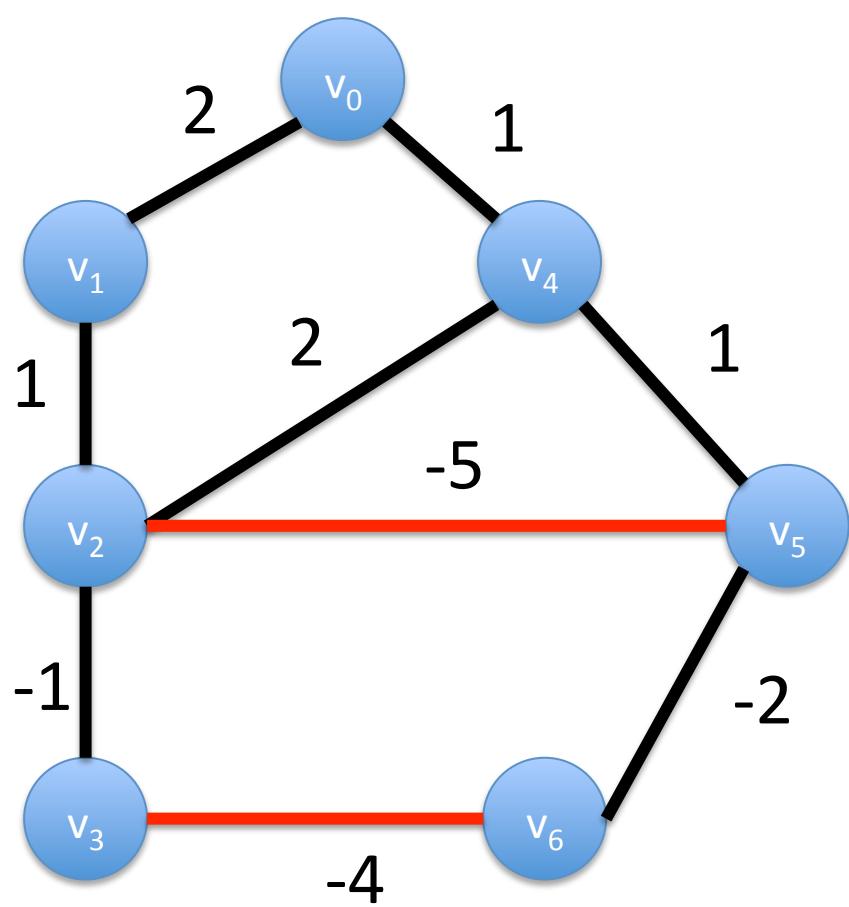
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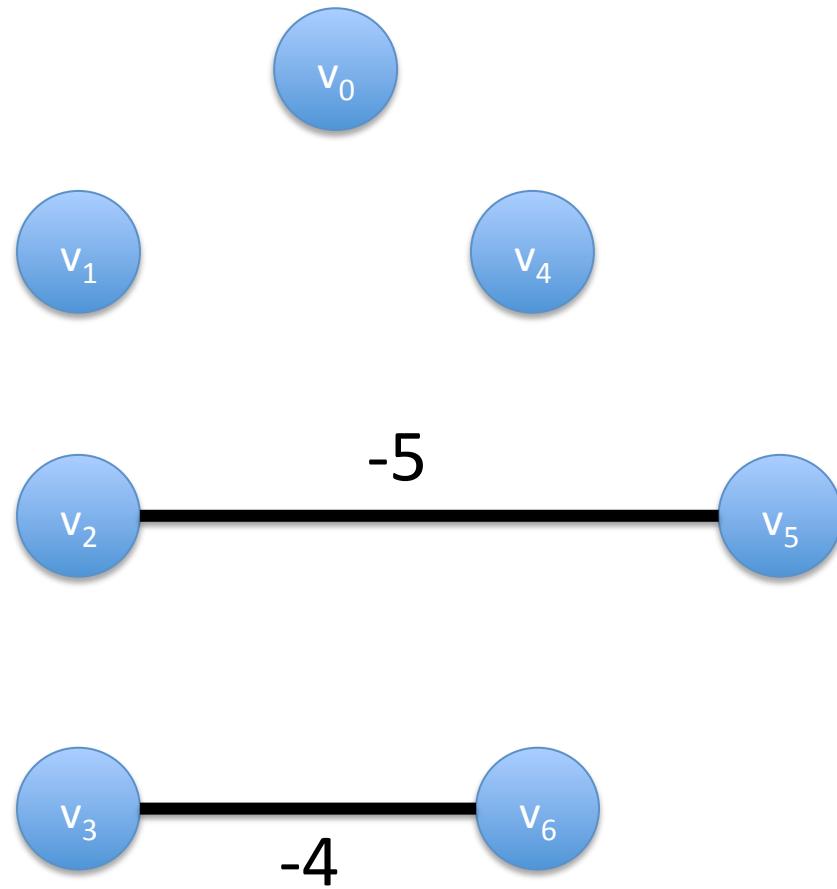
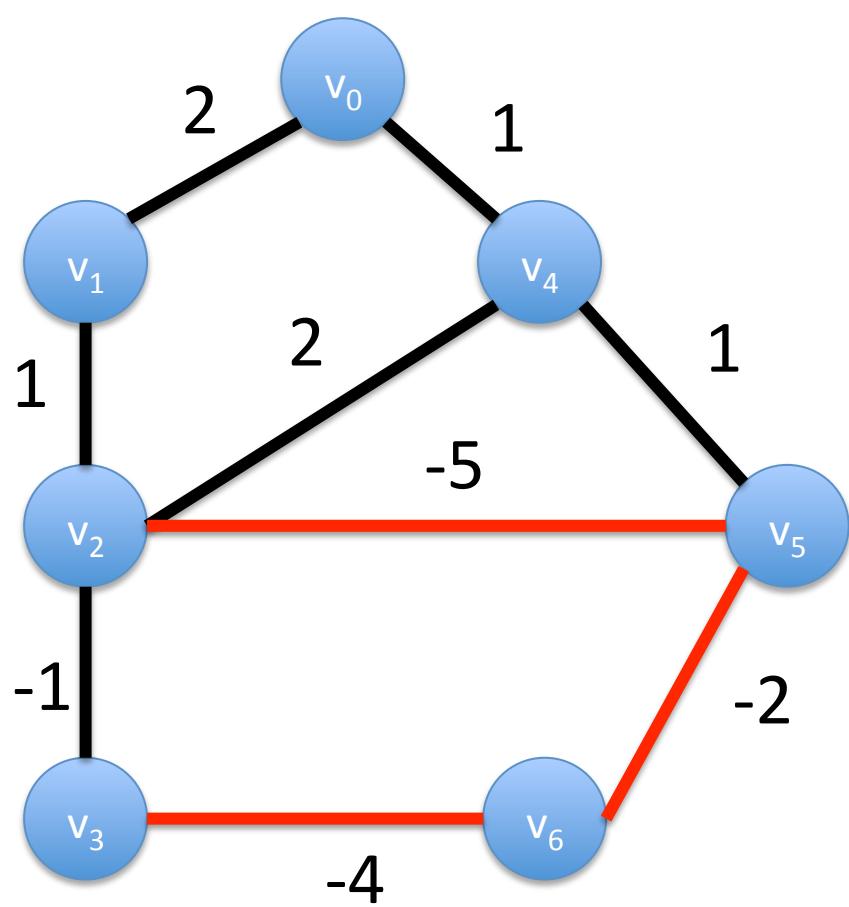


Kruskal's Method



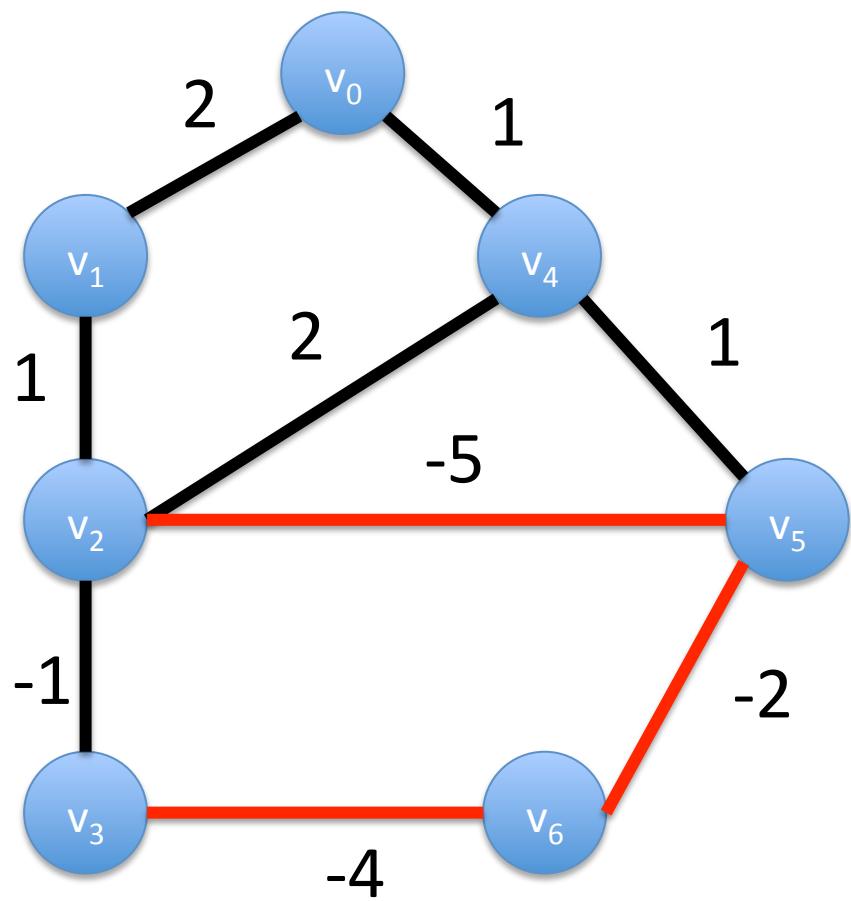
Select the edge with the minimum length

Kruskal's Method



Select the edge with the minimum length

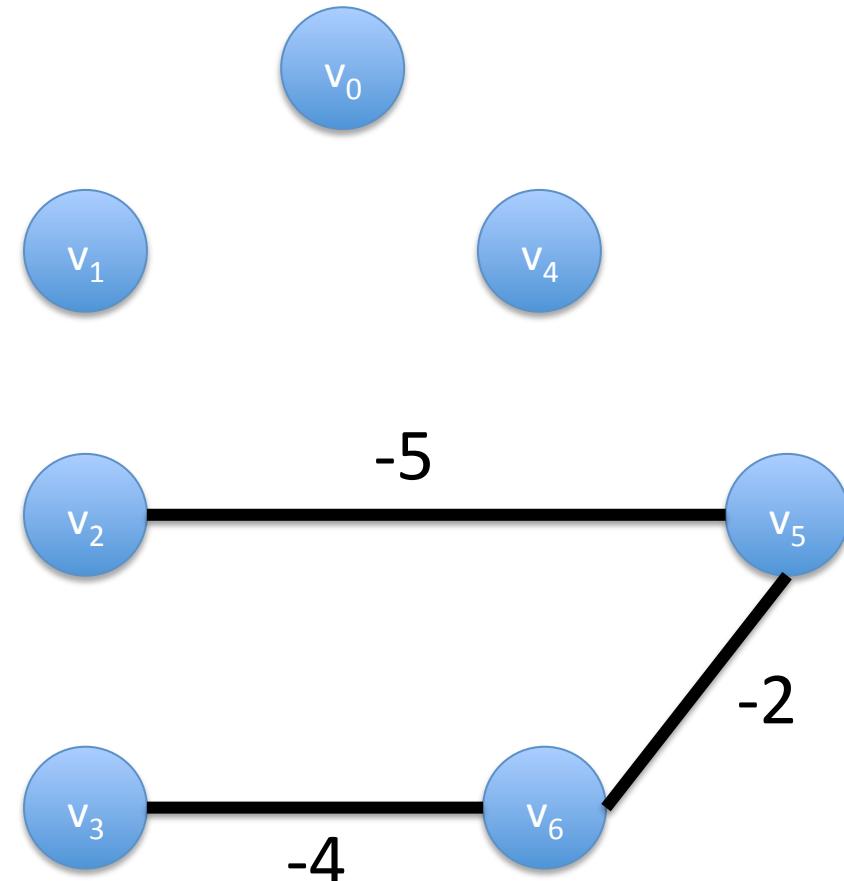
Kruskal's Method



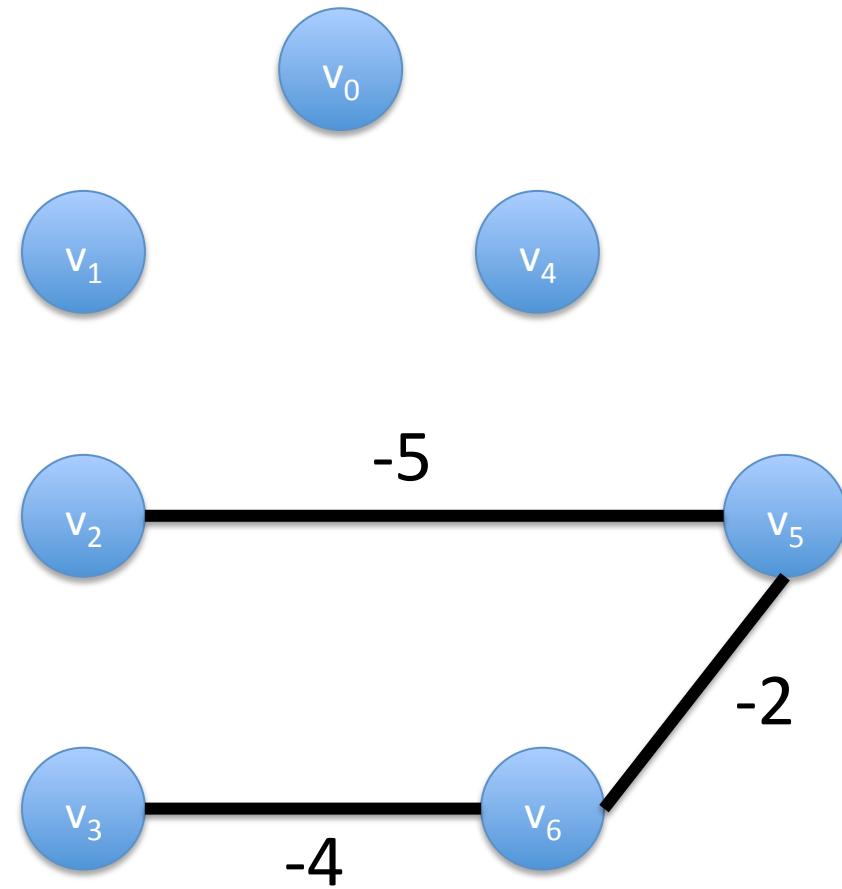
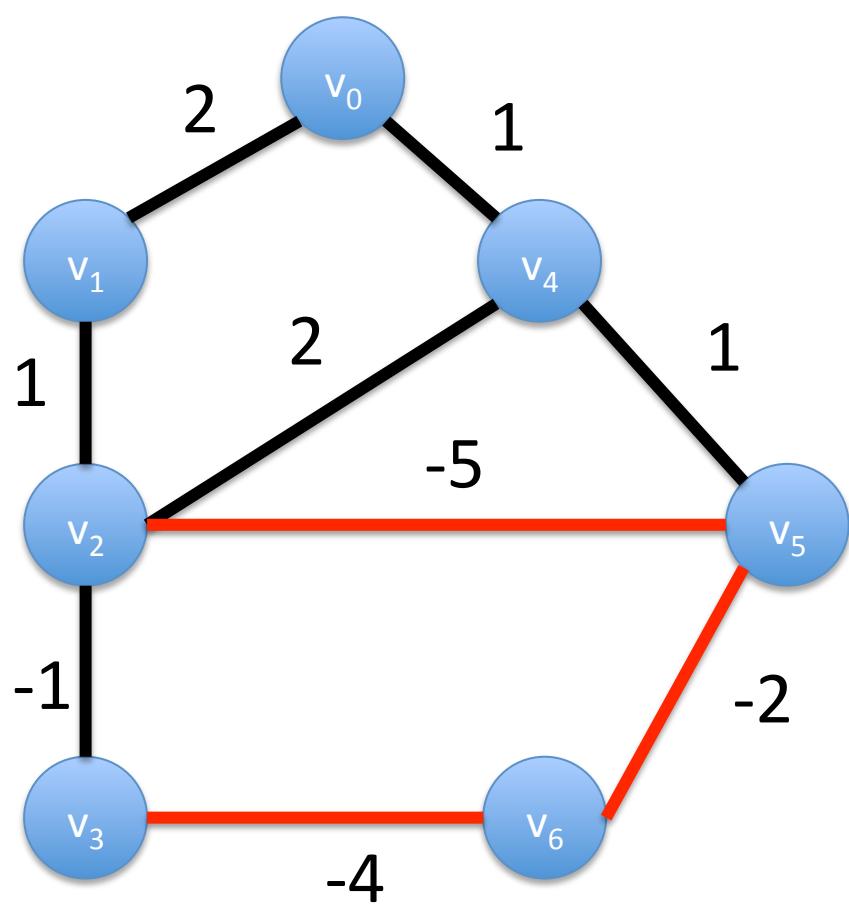
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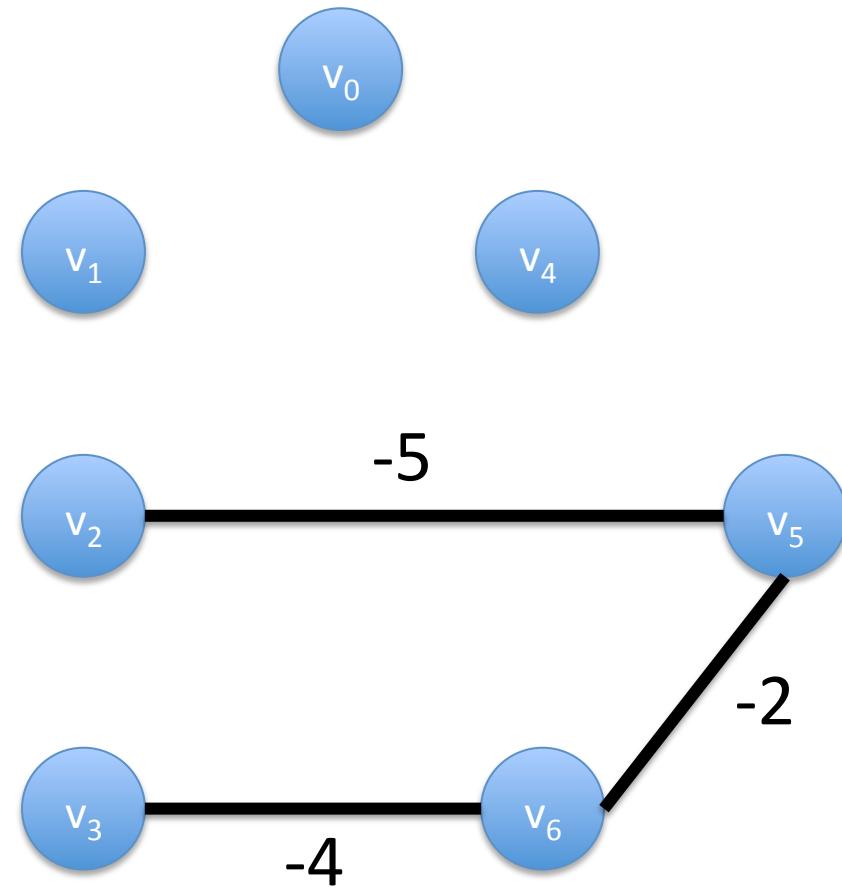
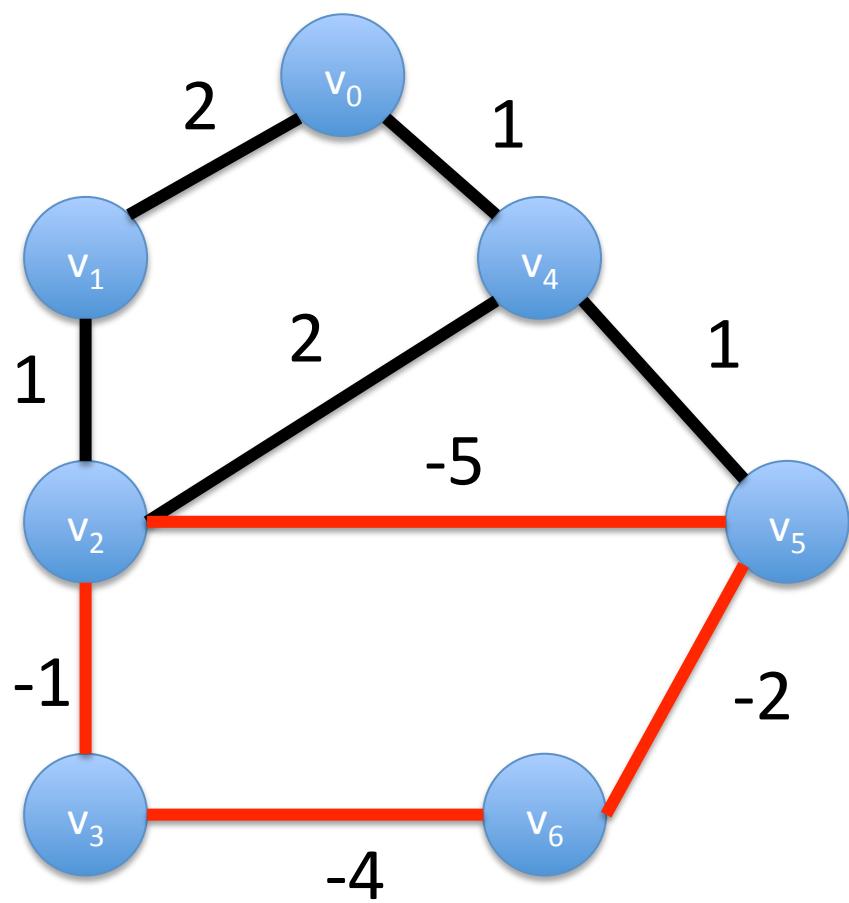


Kruskal's Method



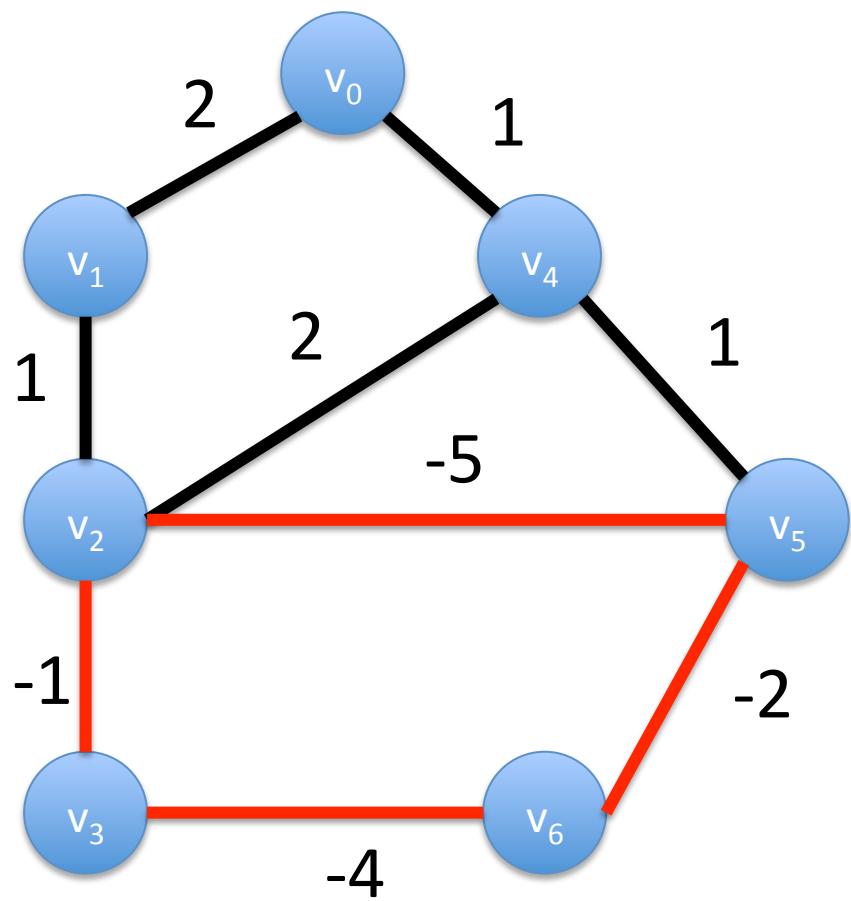
Select the edge with the minimum length

Kruskal's Method

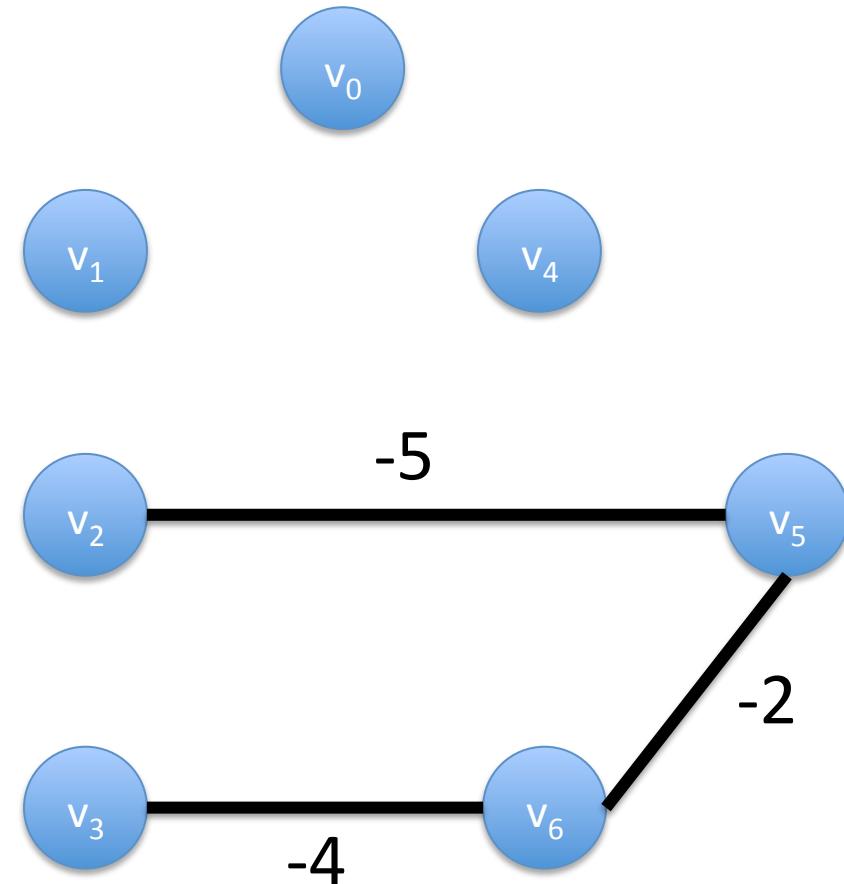


Select the edge with the minimum length

Kruskal's Method



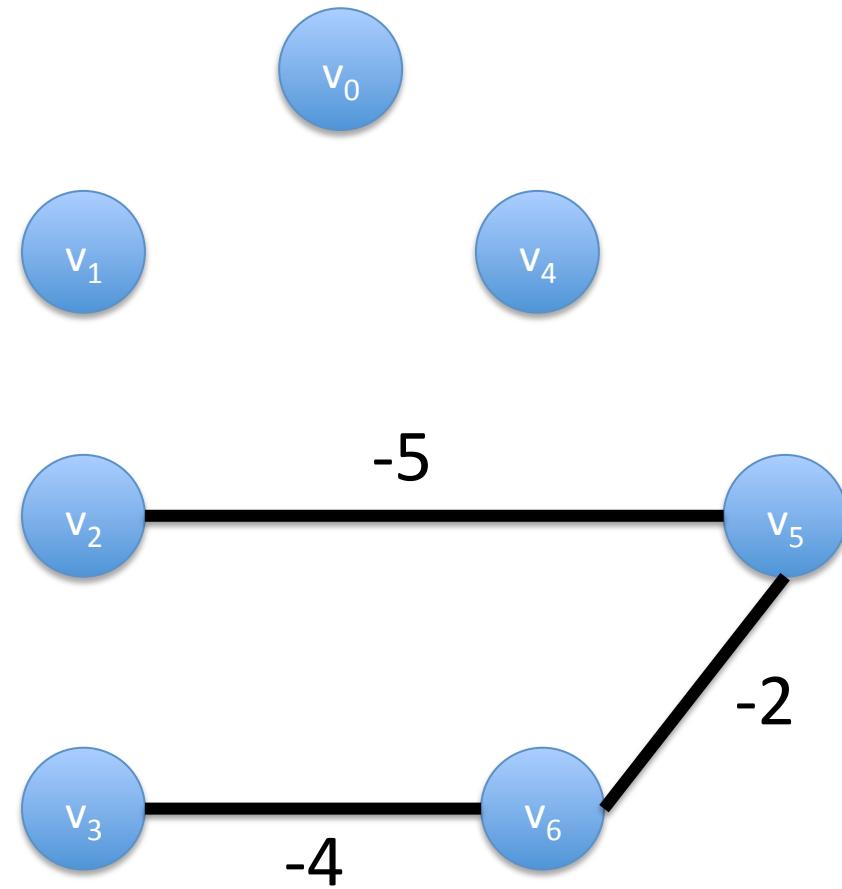
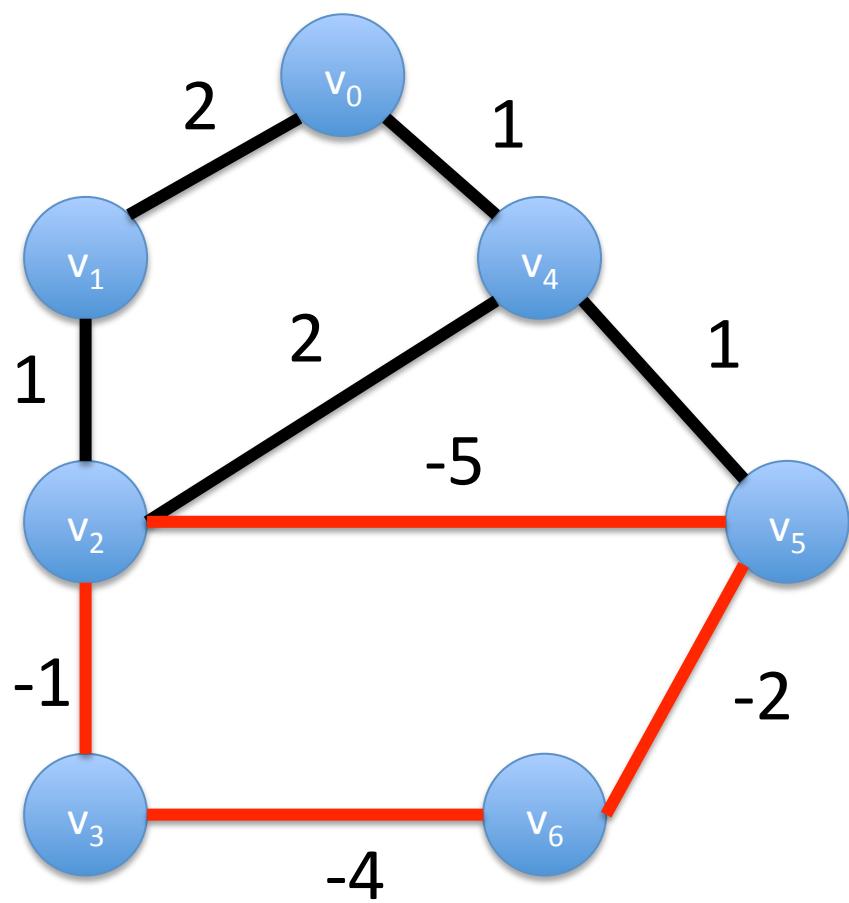
NO



Discard the edge

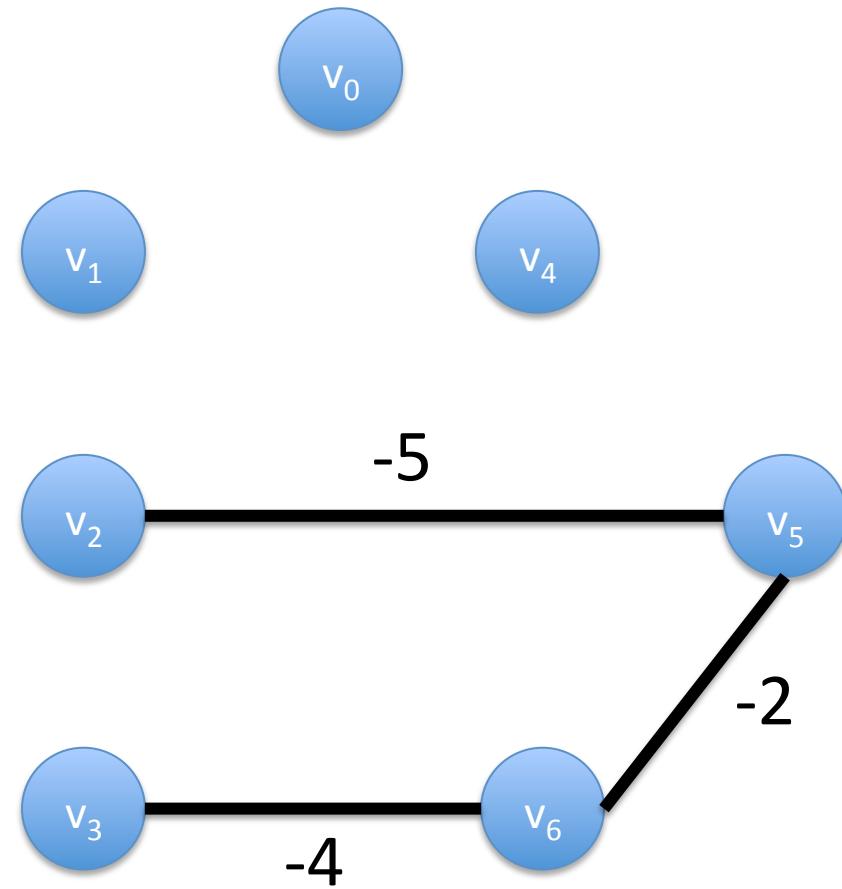
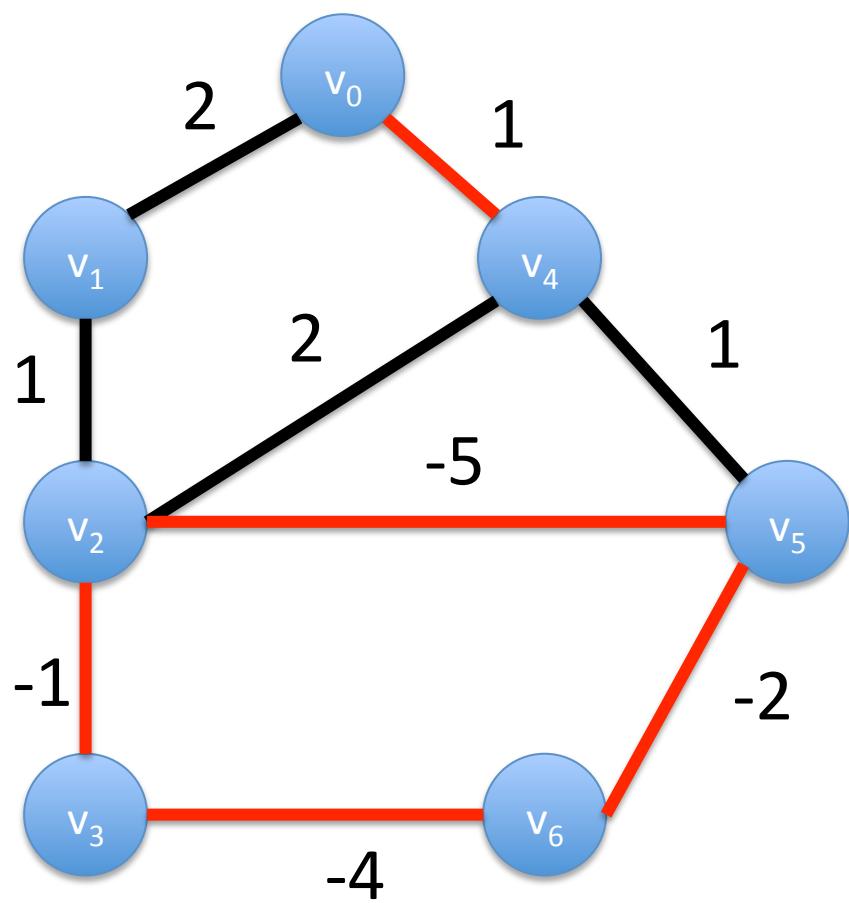
Does this edge connect two different trees?

Kruskal's Method



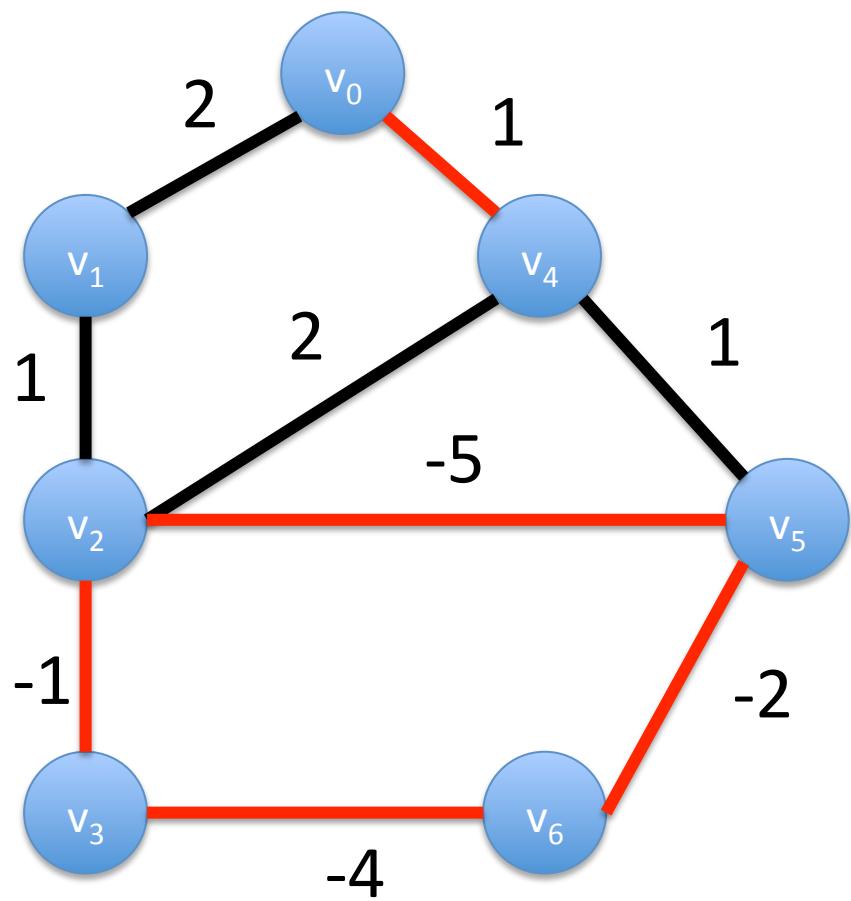
Select the edge with the minimum length

Kruskal's Method

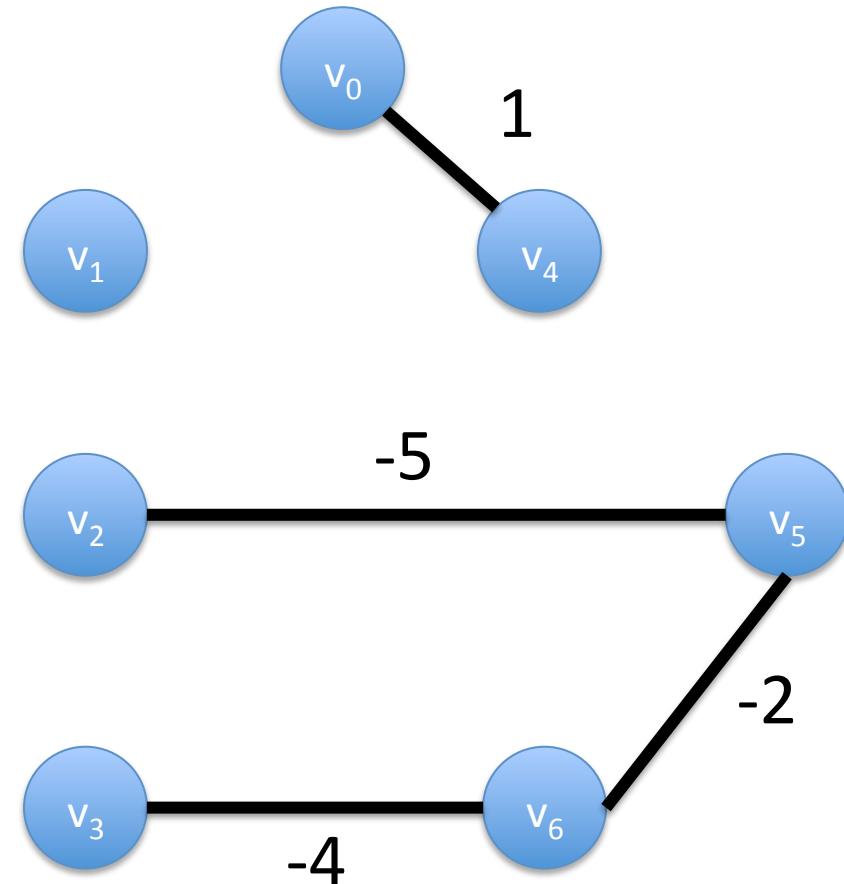


Select the edge with the minimum length

Kruskal's Method



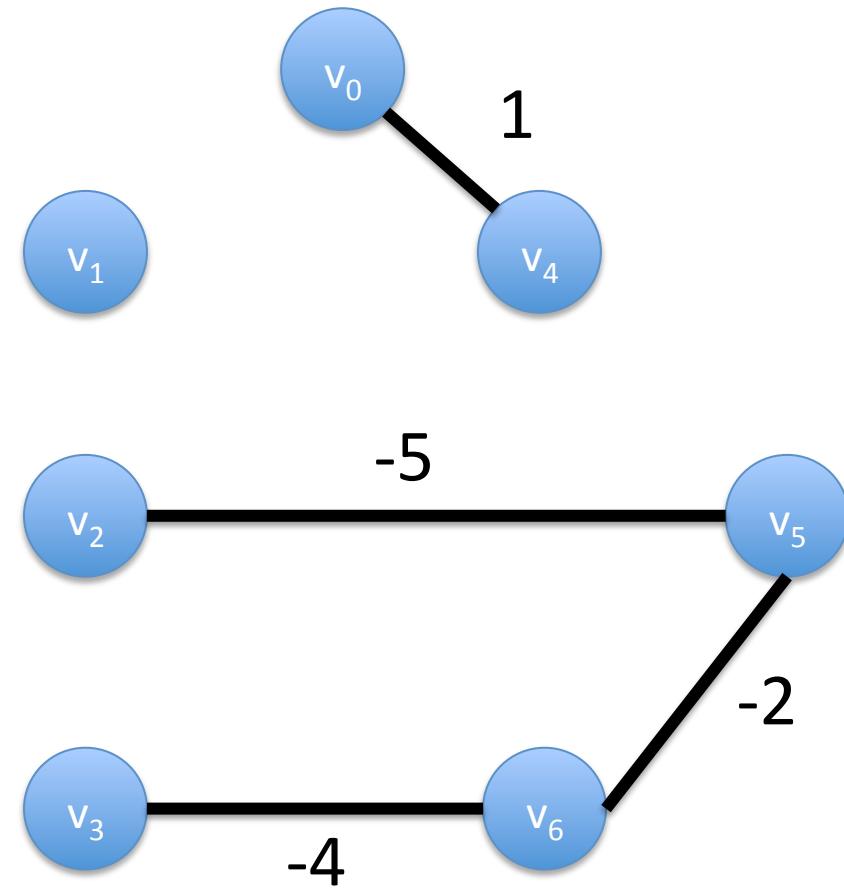
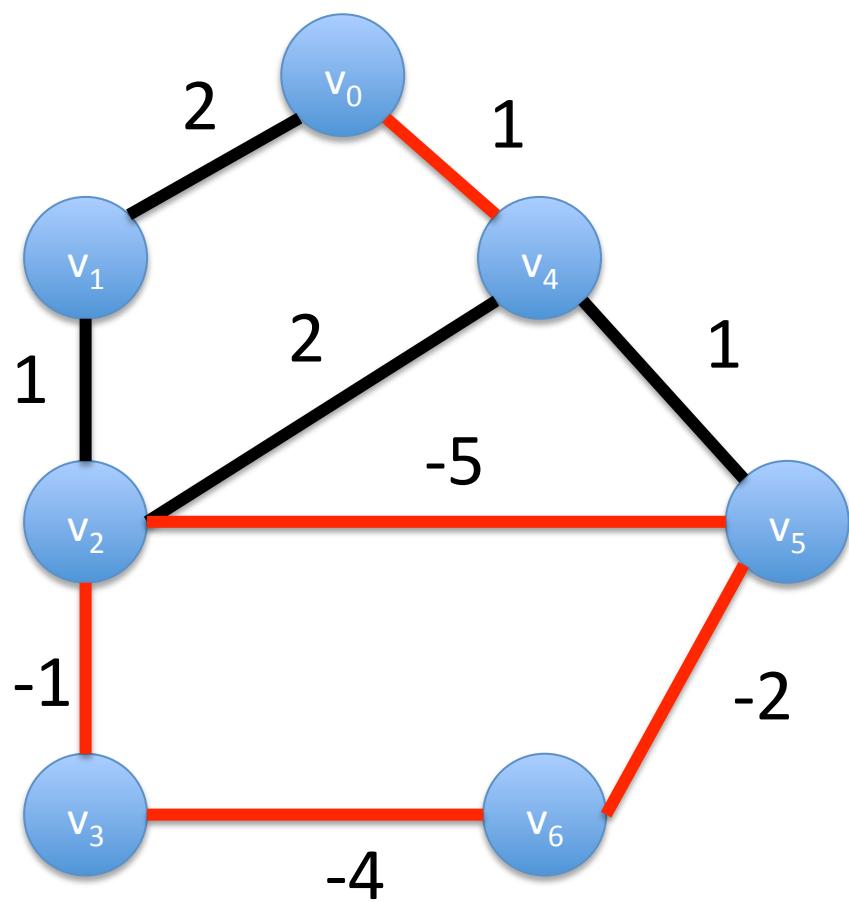
YES



Add the edge to the forest

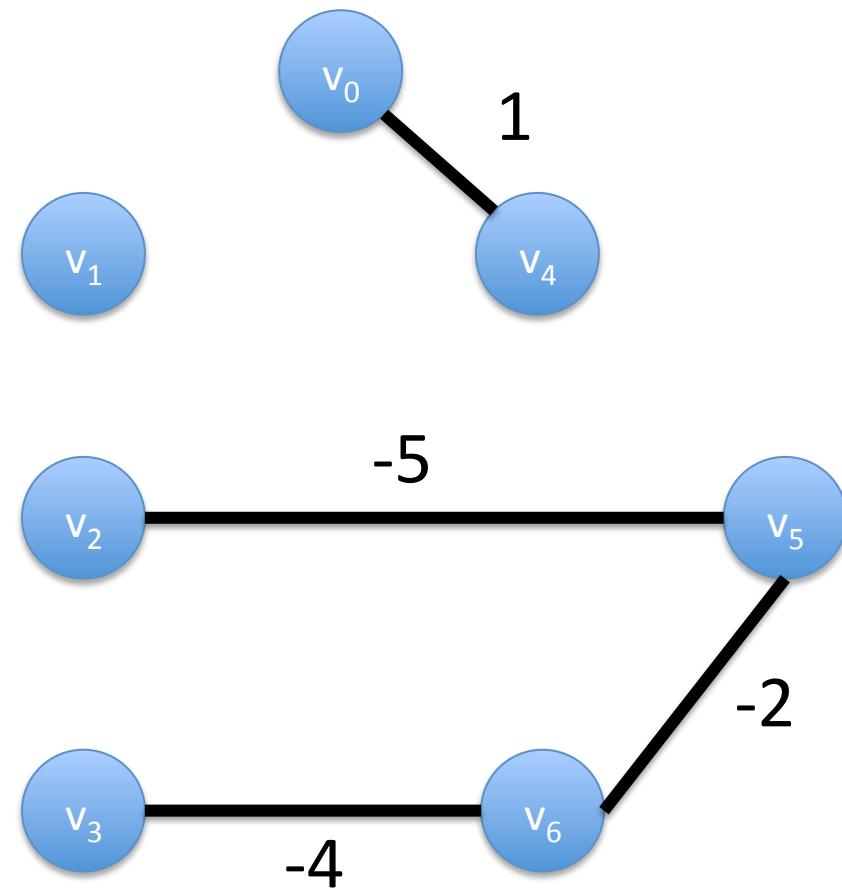
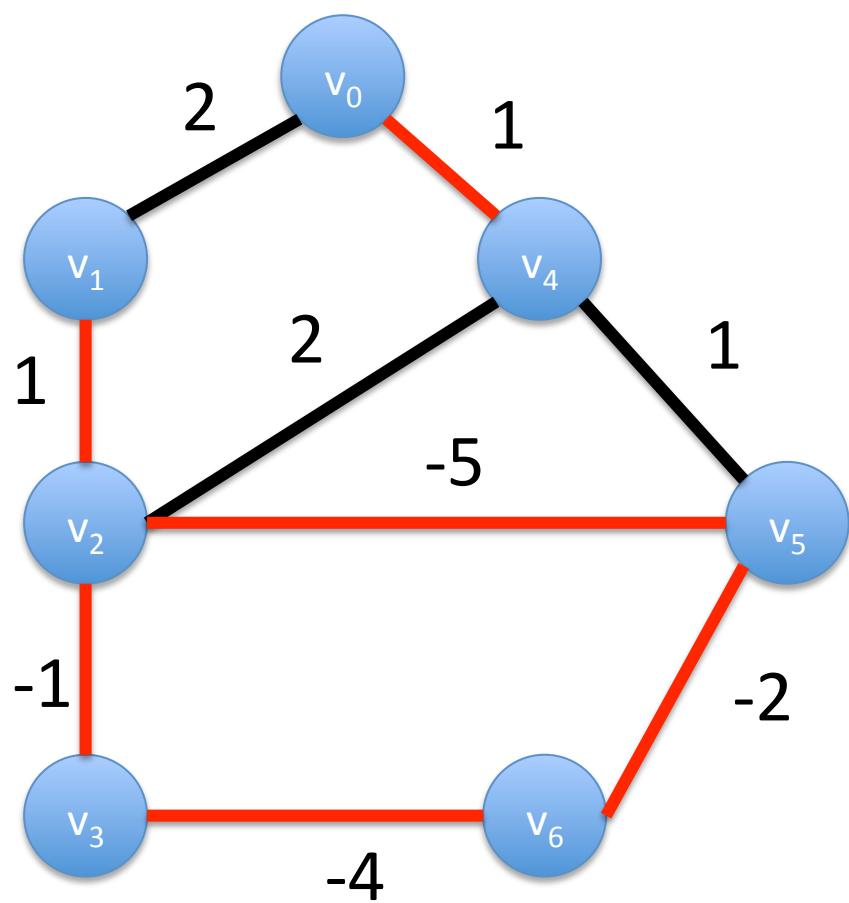
Does this edge connect two different trees?

Kruskal's Method



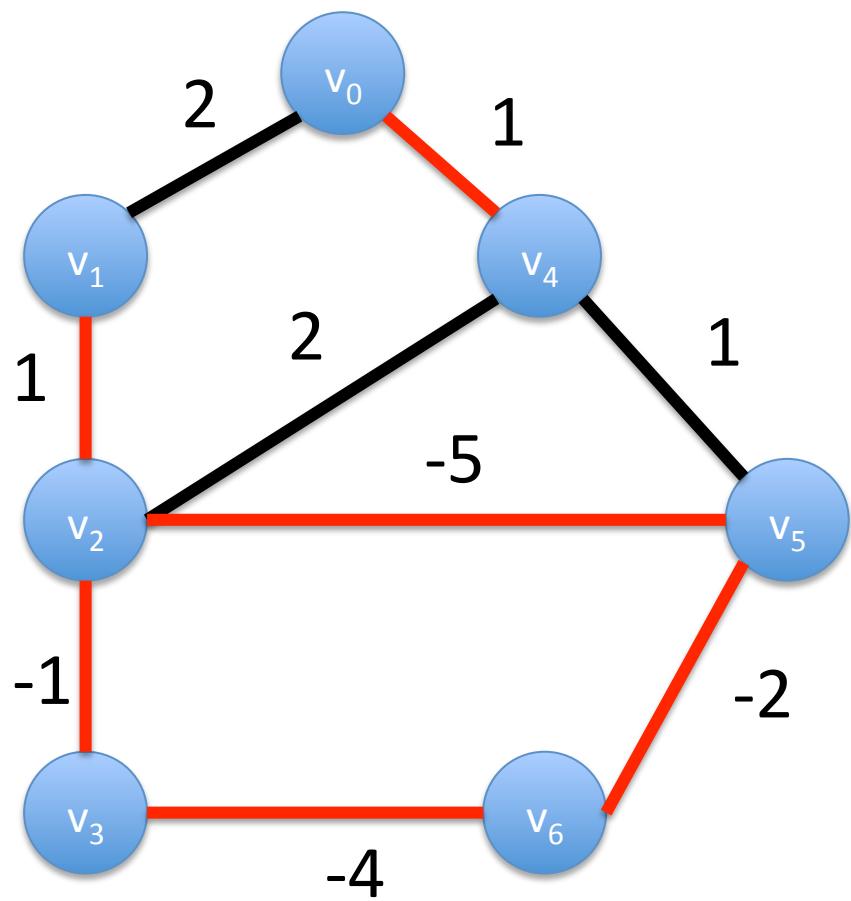
Select the edge with the minimum length

Kruskal's Method

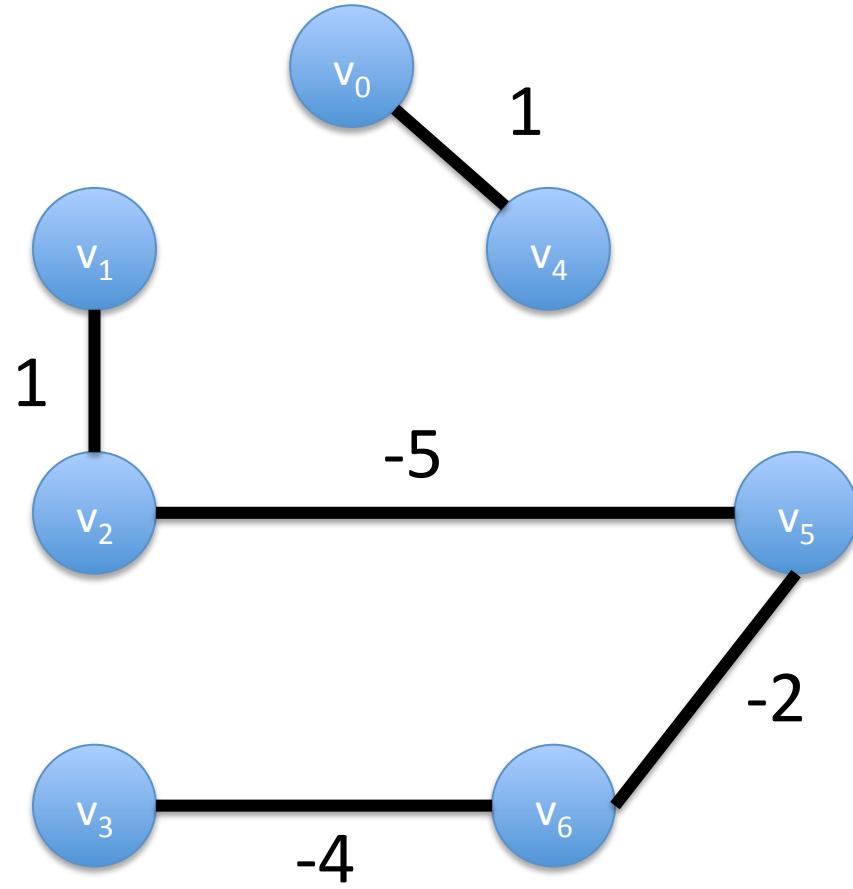


Select the edge with the minimum length

Kruskal's Method



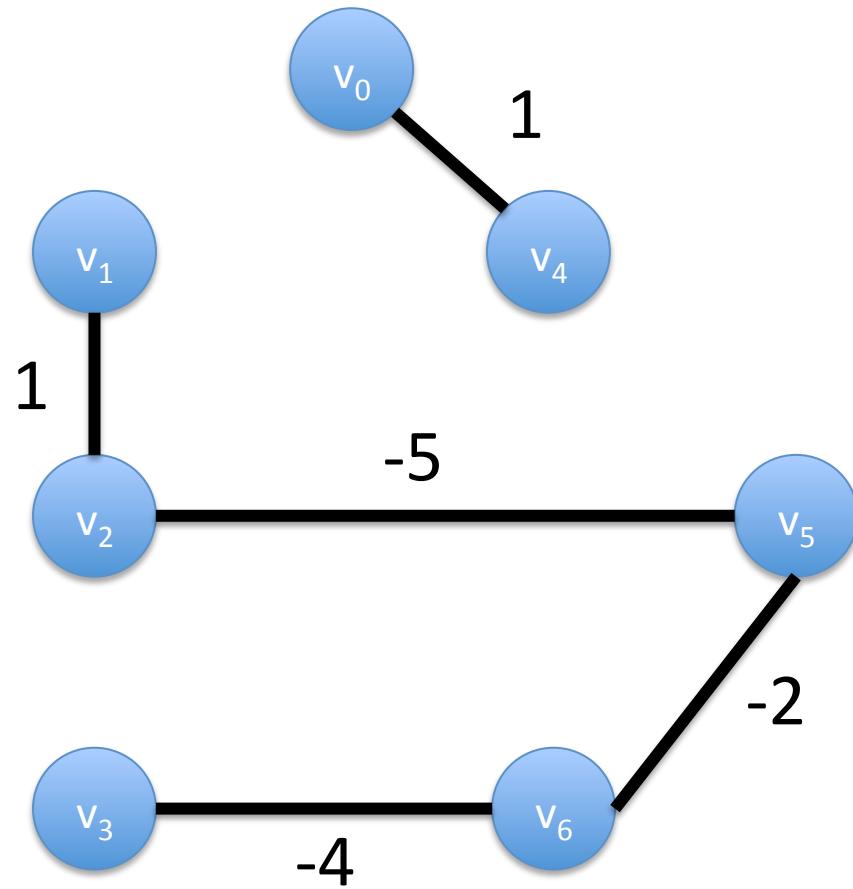
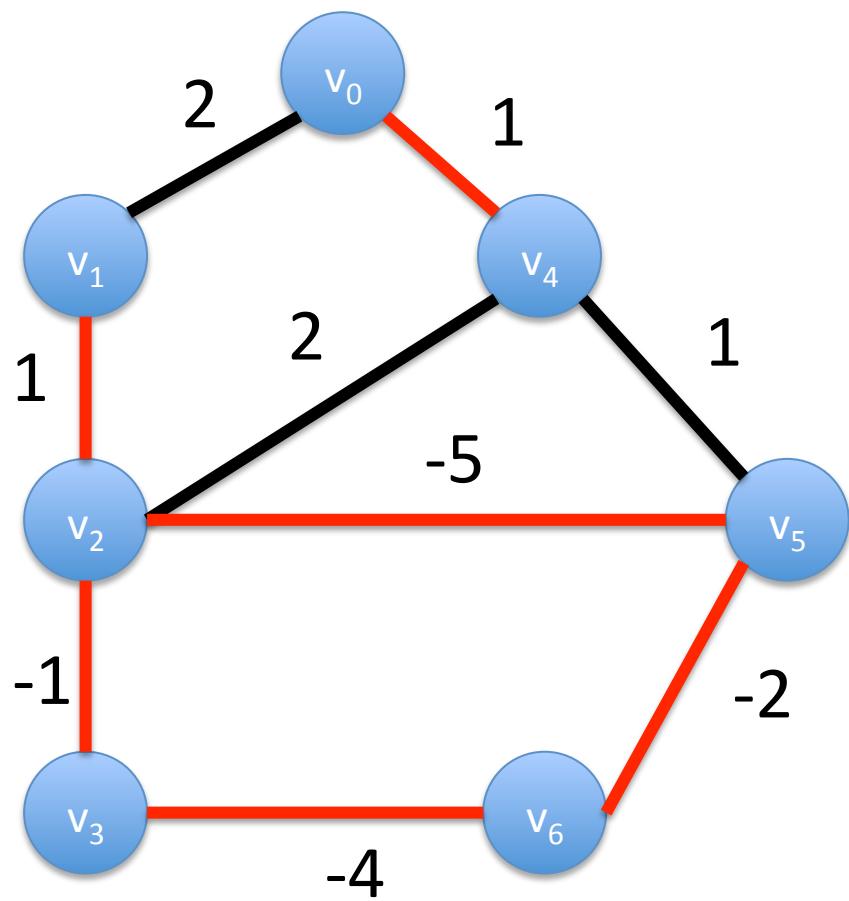
YES



Add the edge to the forest

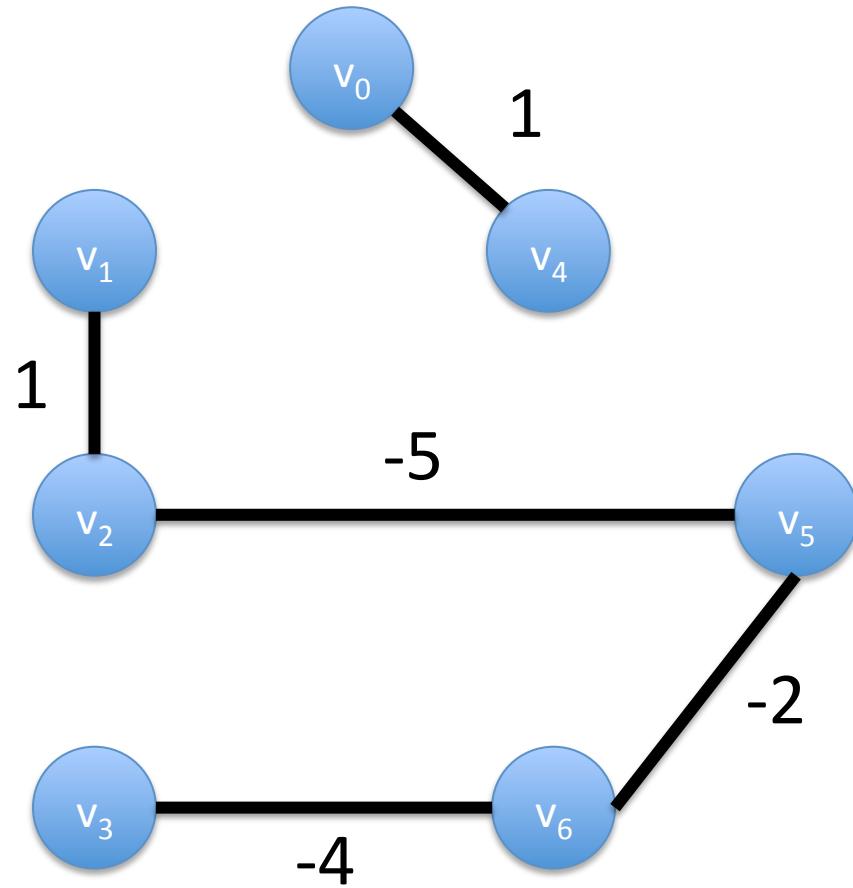
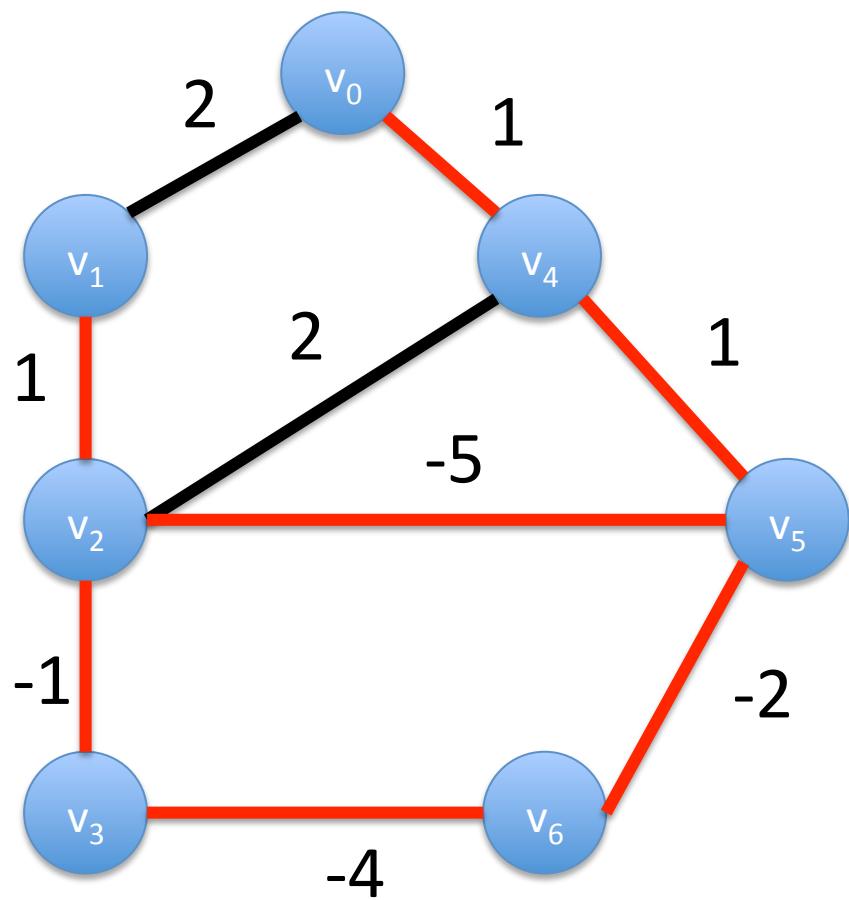
Does this edge connect two different trees?

Kruskal's Method



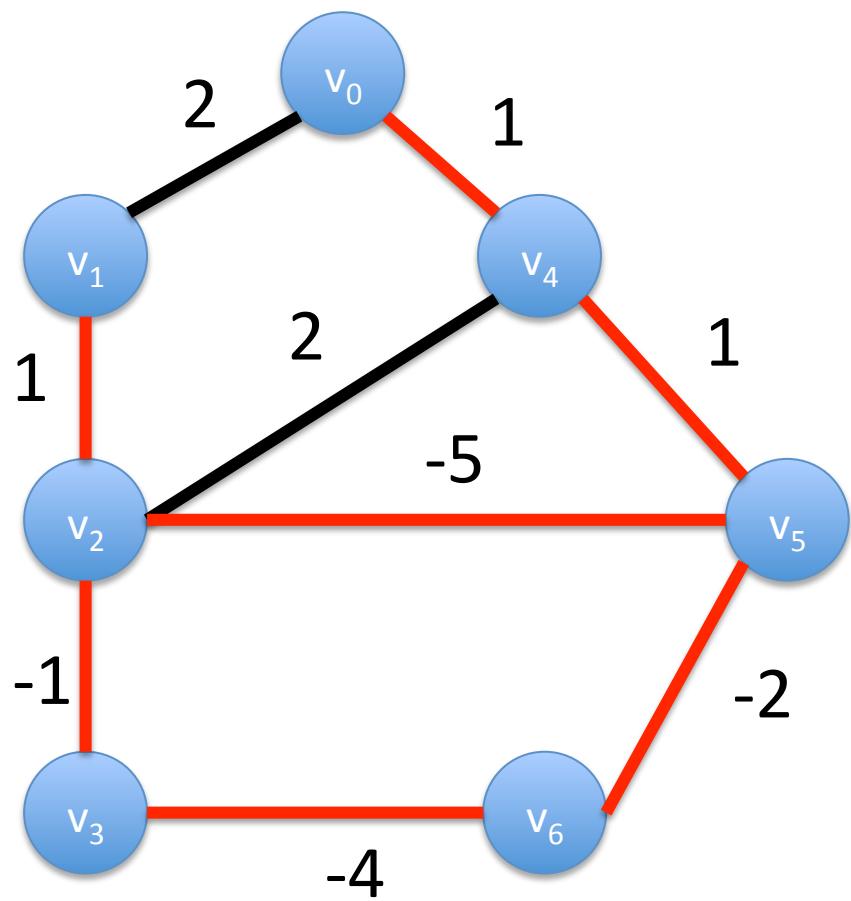
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Kruskal's Method

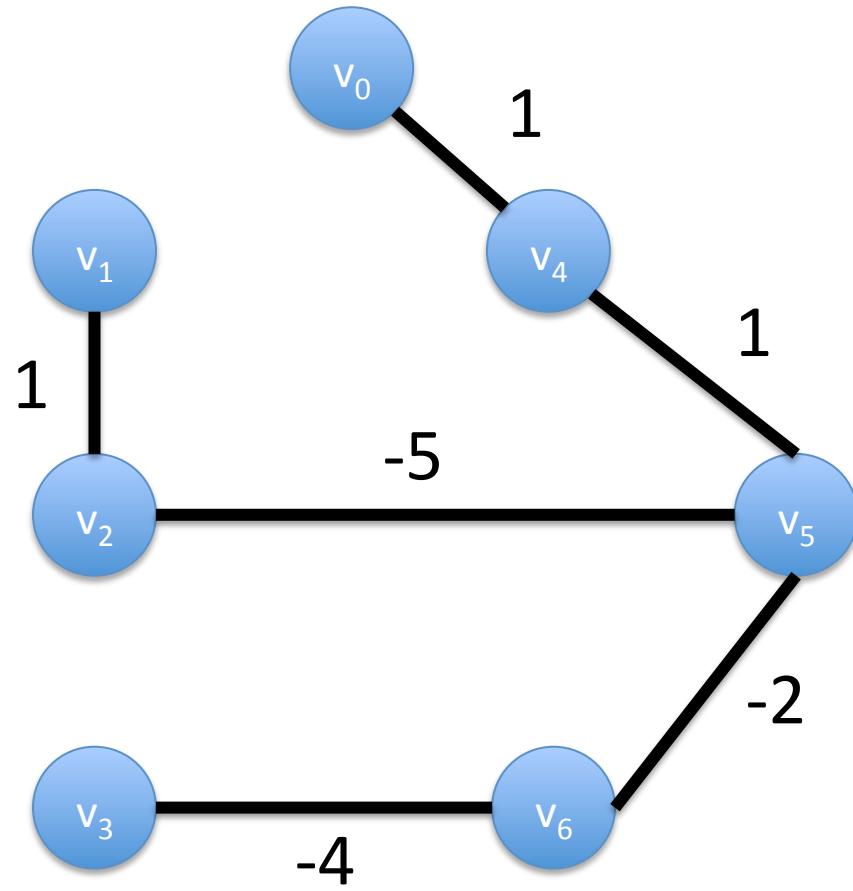


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Kruskal's Method



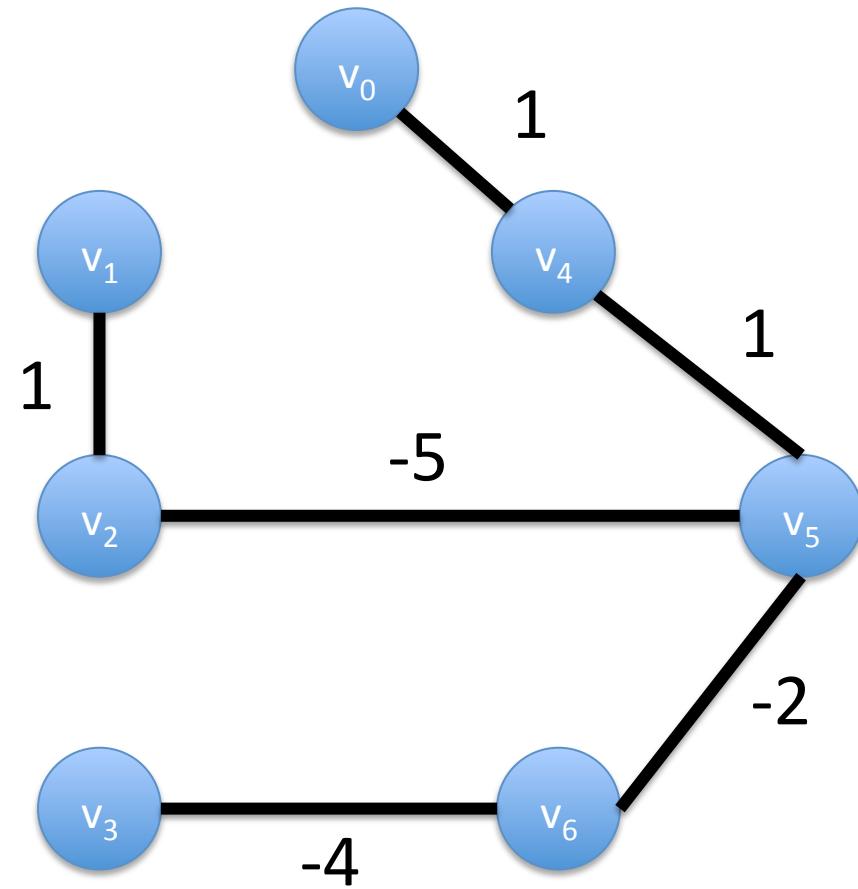
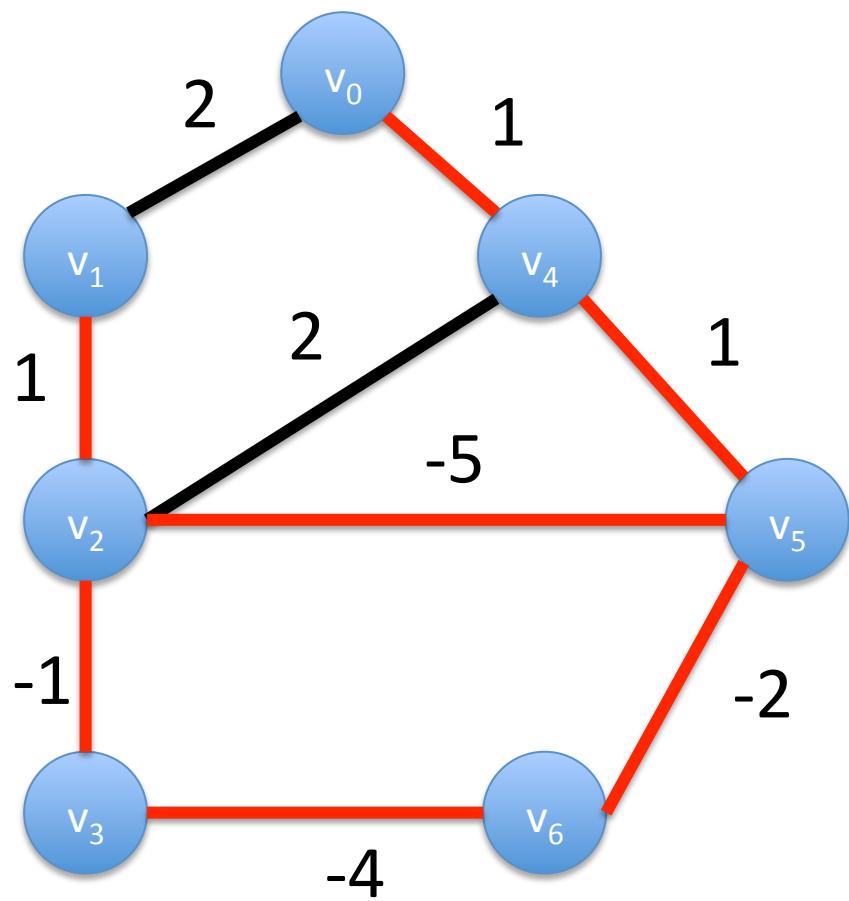
YES



Add the edge to the forest

Does this edge connect two different trees?

Kruskal's Method



Spanning Tree !!! Algorithm terminates.

Summary

Given $G = (V, E)$. Define $T = (V, \{\})$ and $S = E$.

While T is not a spanning tree

Select the minimum length edge e from S

If e connects two different trees

Add e to T

End

Remove e from S

End

Time Complexity

$O(m \log(m))$ where $m = |E|$

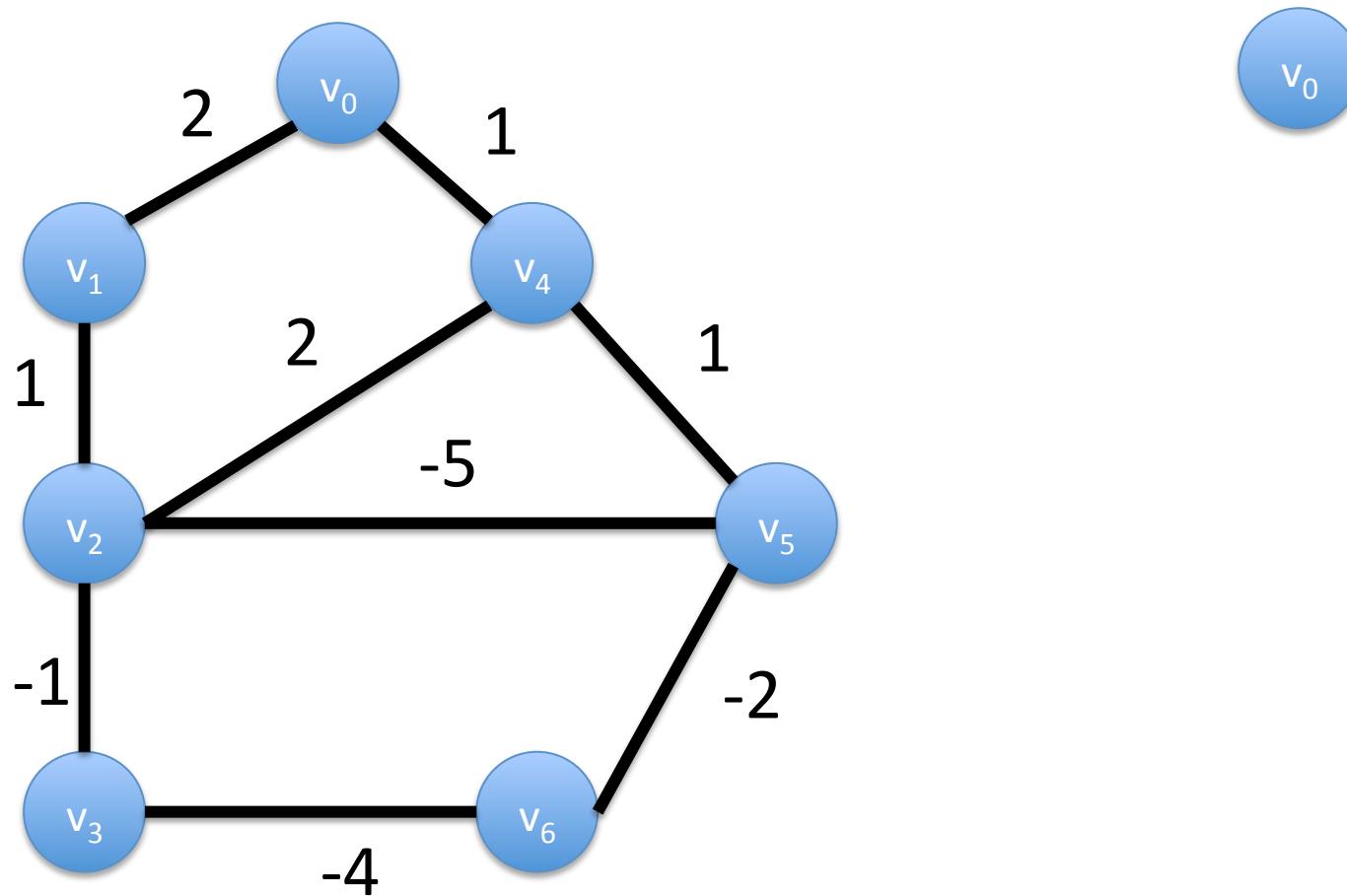
Proof of Kruskal's Algorithm

On the board

Outline

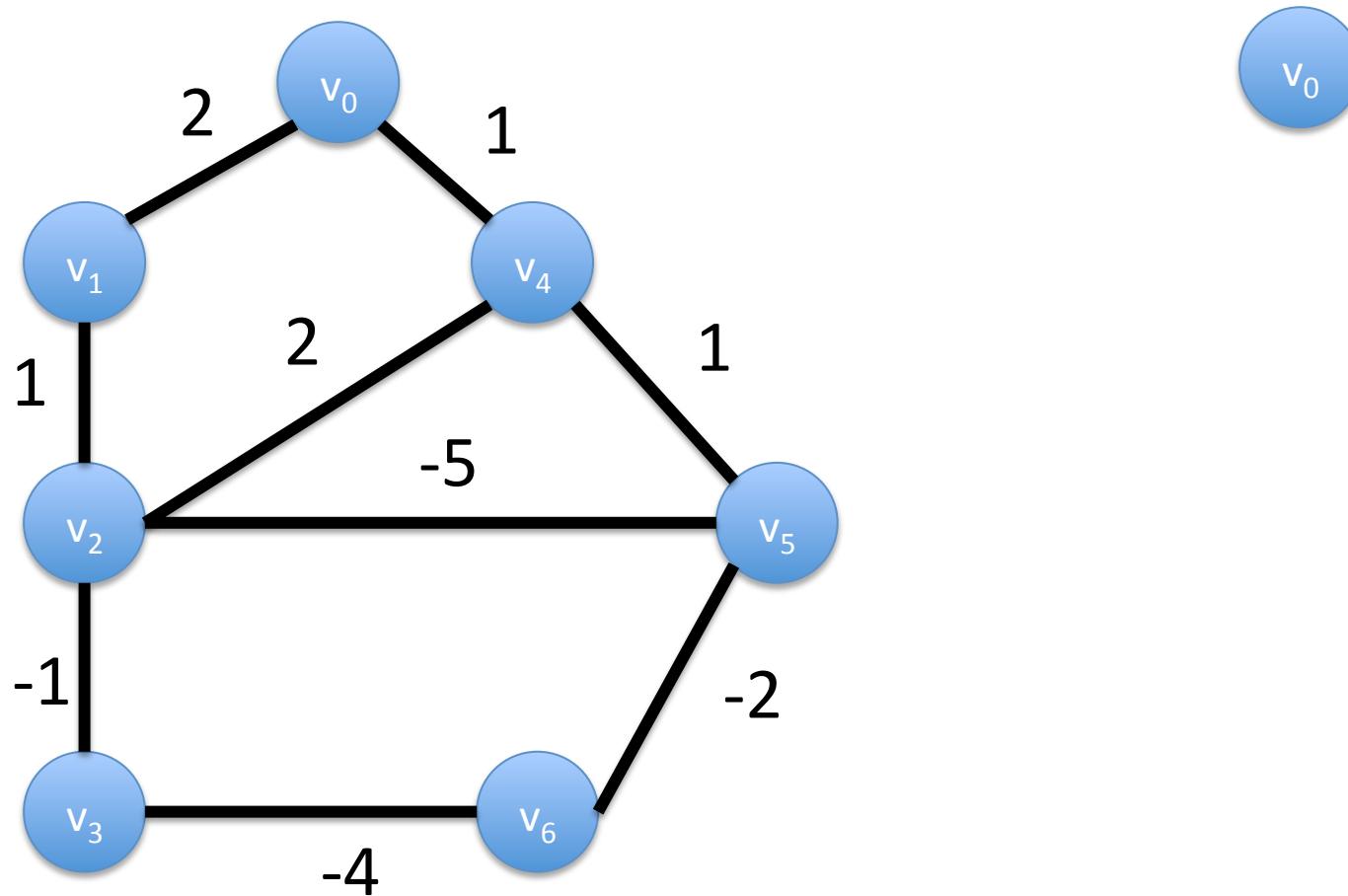
- Chow-Liu Tree
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Prim's Method



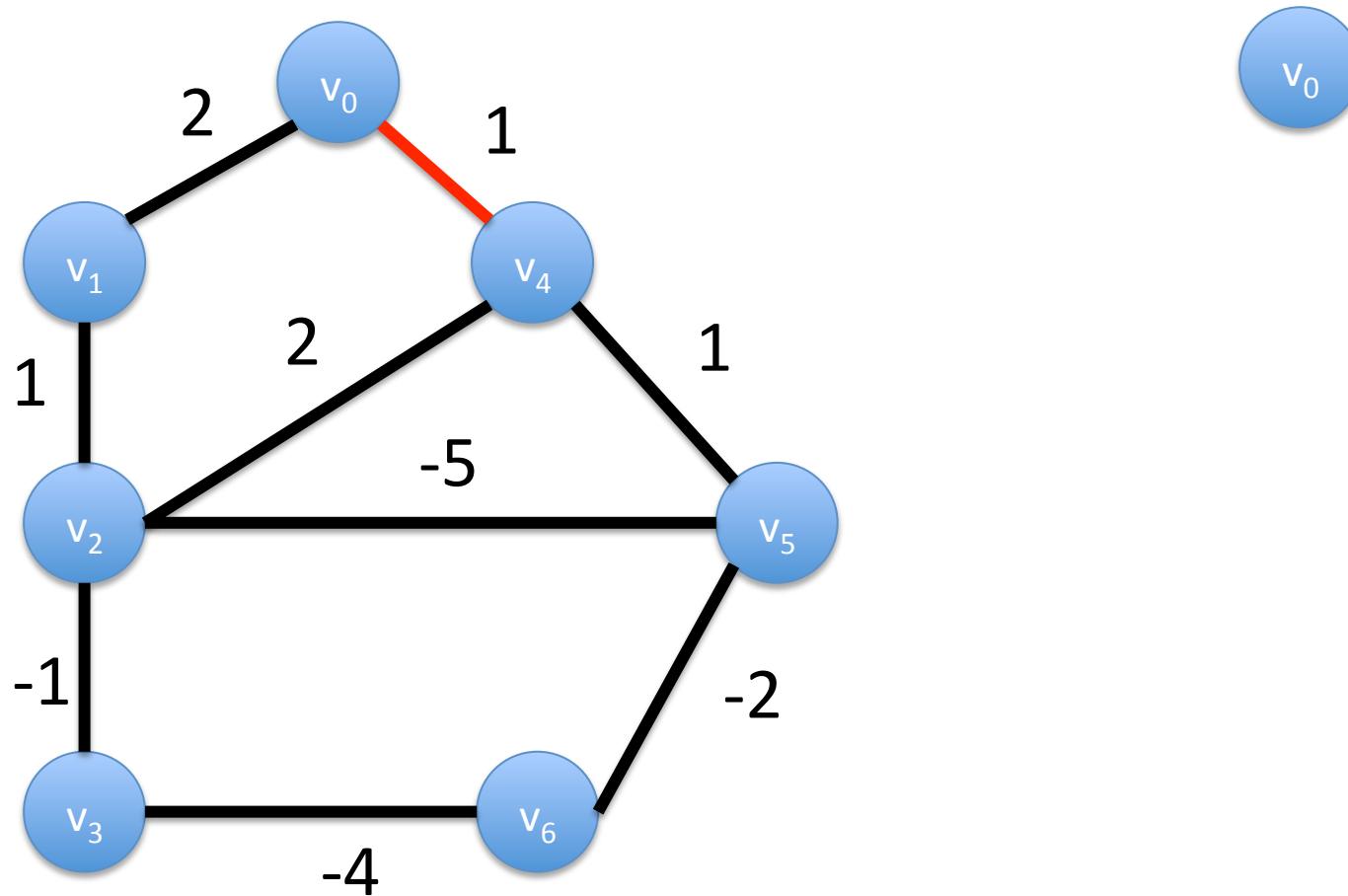
Initialize $T = (V_T, \{\})$ where $V_T = \{v_0\}$

Prim's Method



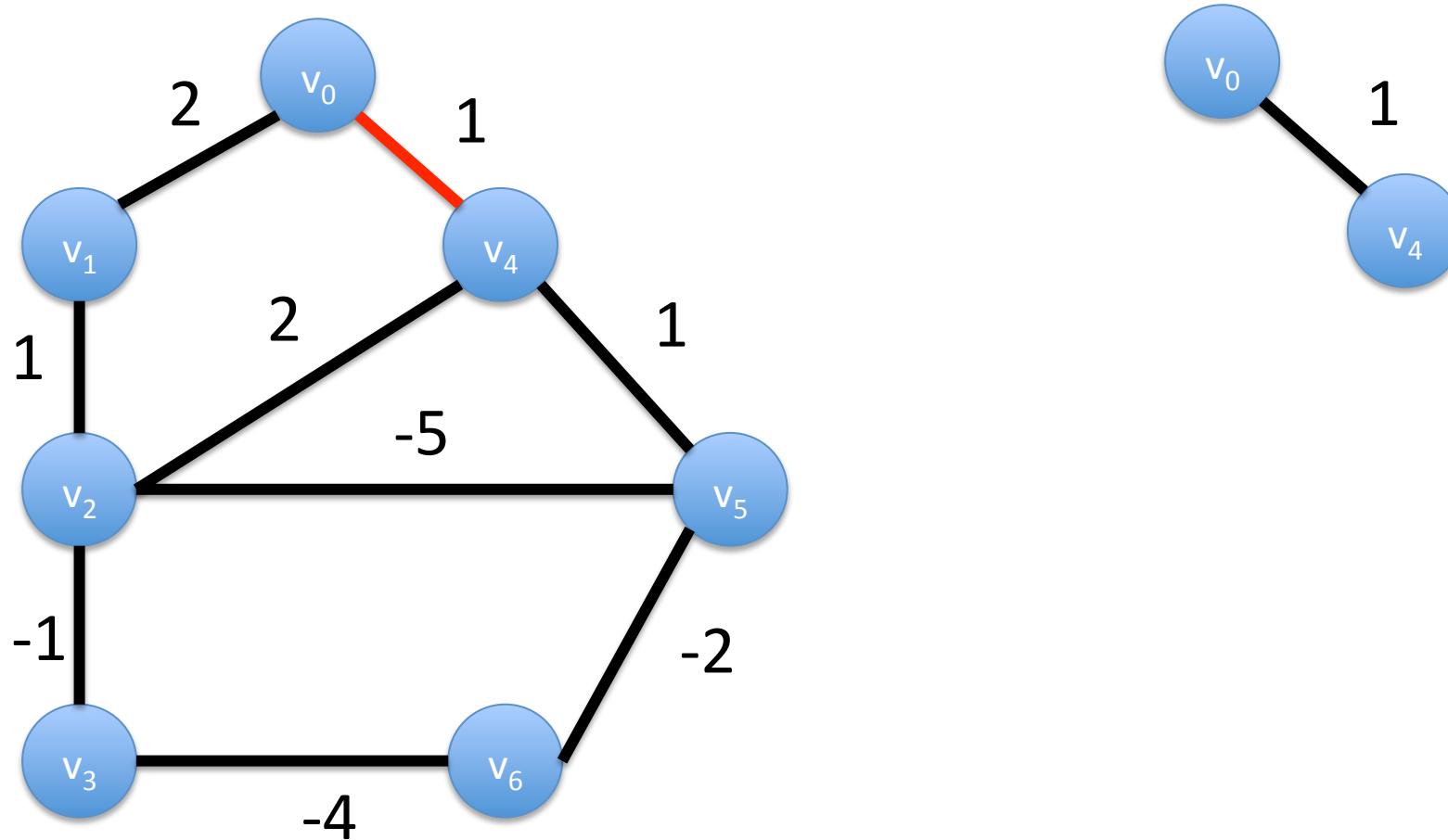
Choose $(u^*, v^*) = \min I(u, v), u \in V_T, v \in V \setminus V_T$

Prim's Method



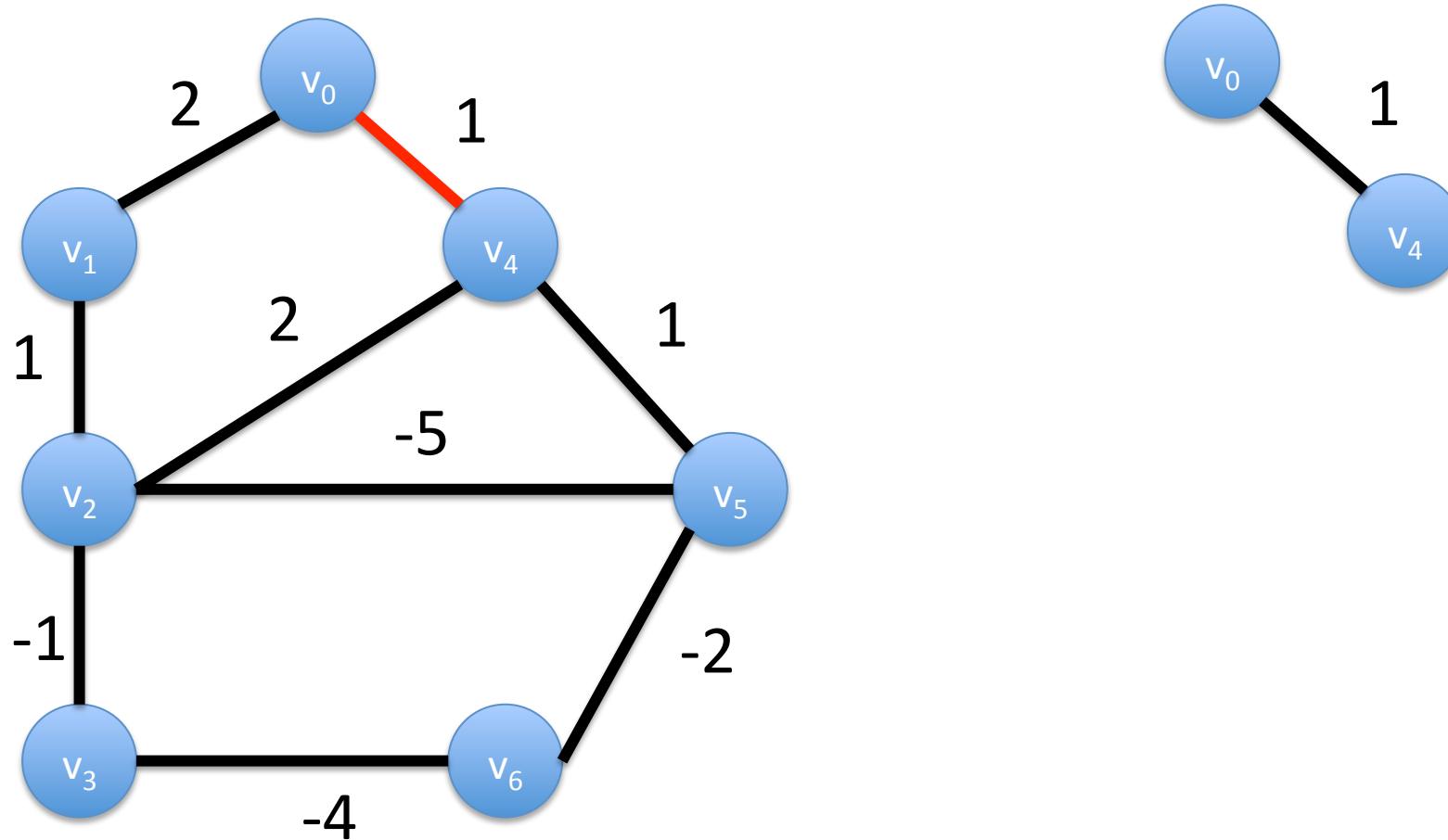
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Prim's Method



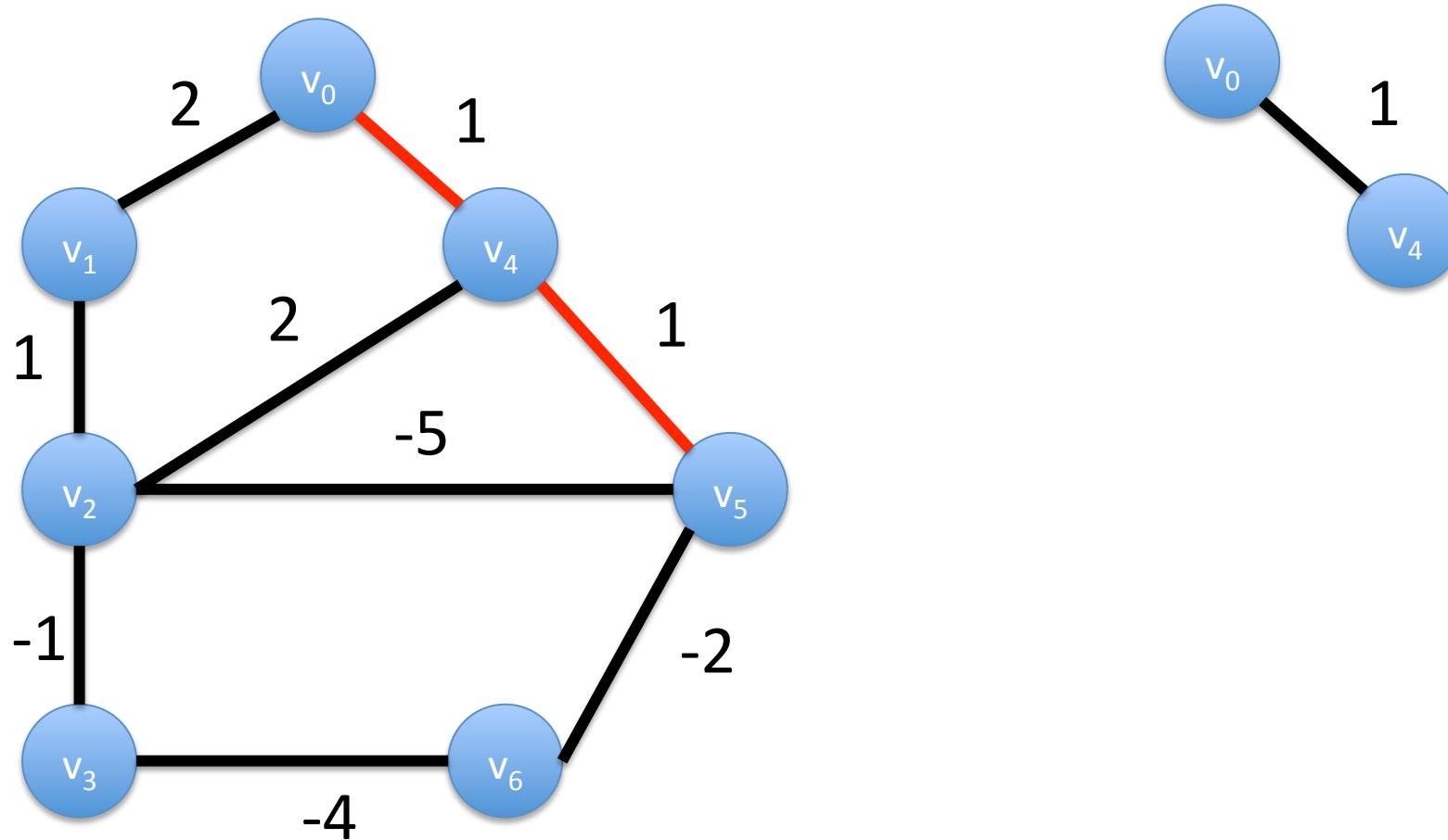
Add v^* to V_T . Add (u^*, v^*) to E_T .

Prim's Method



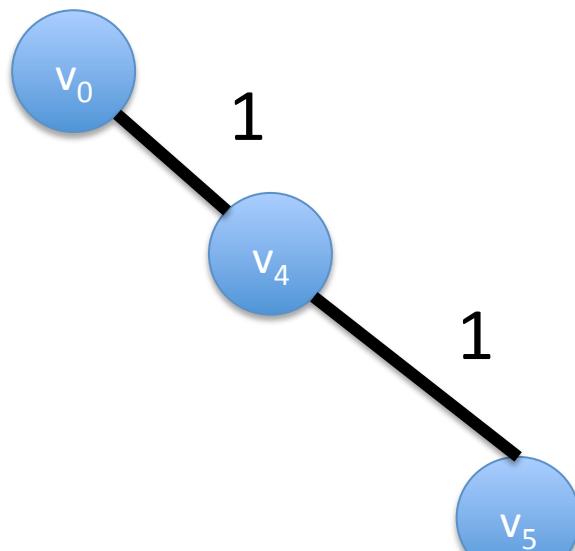
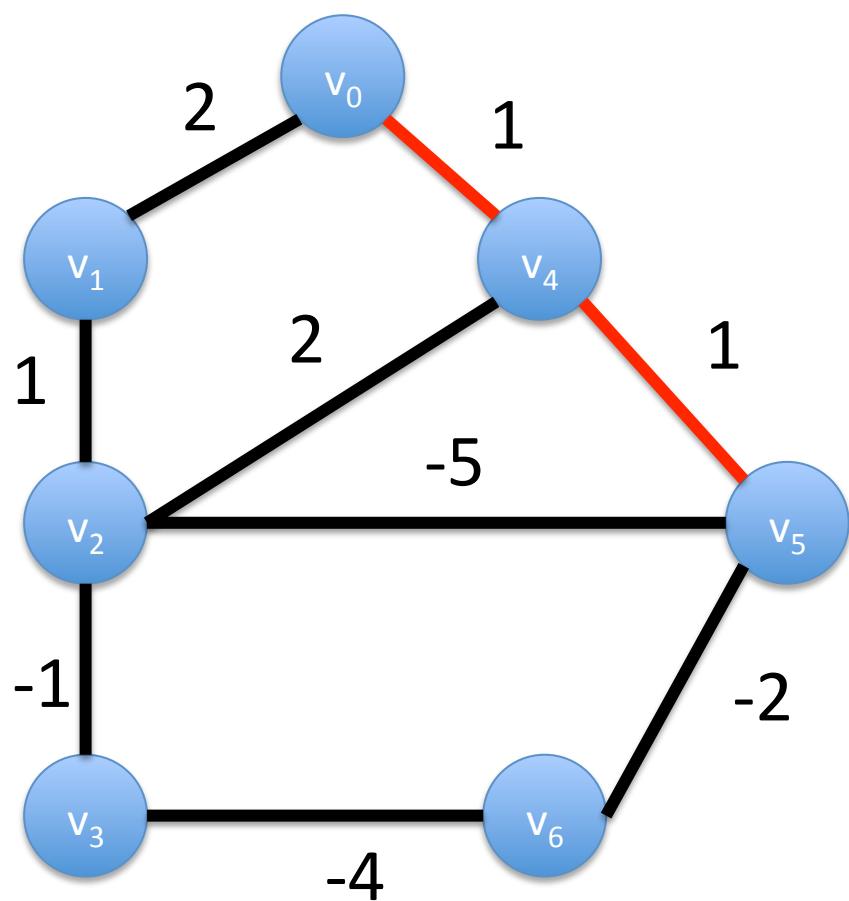
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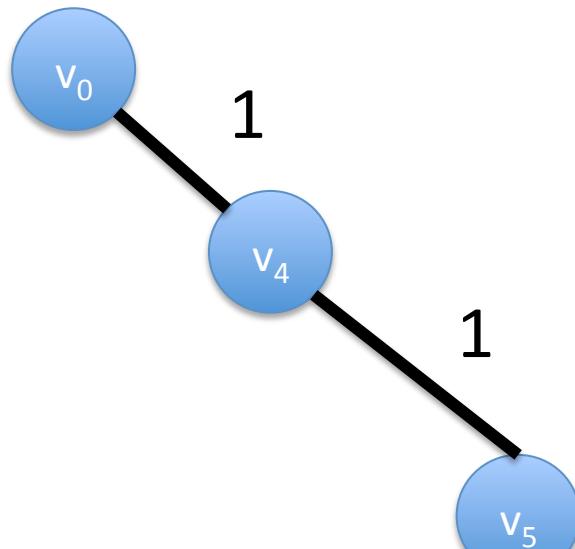
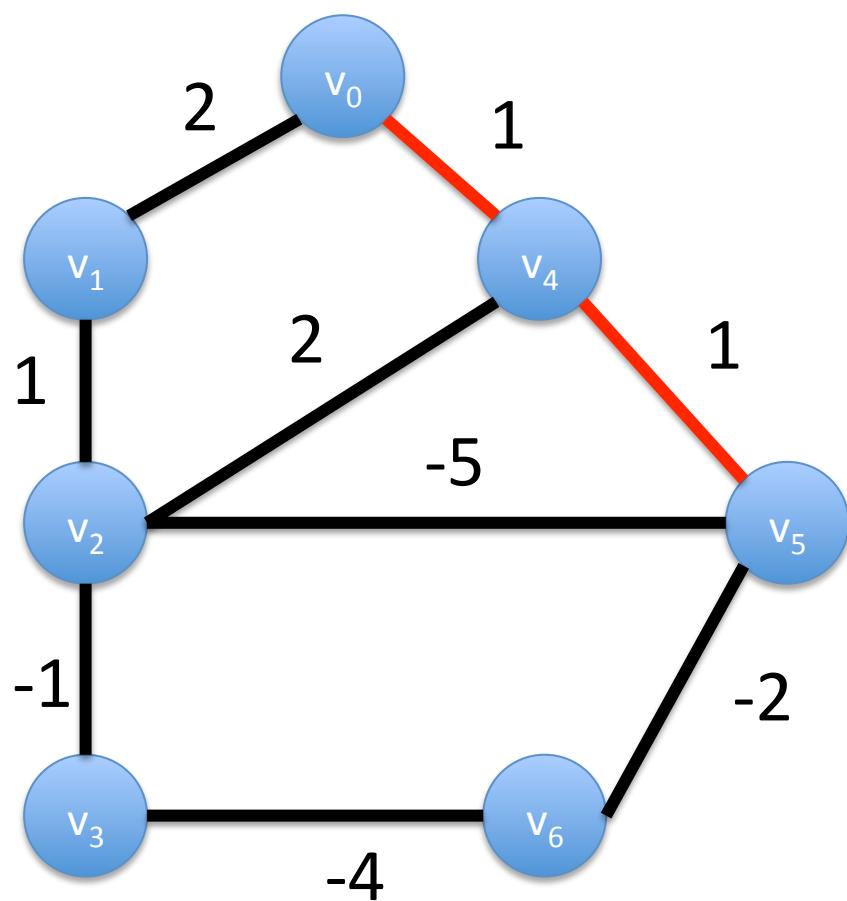
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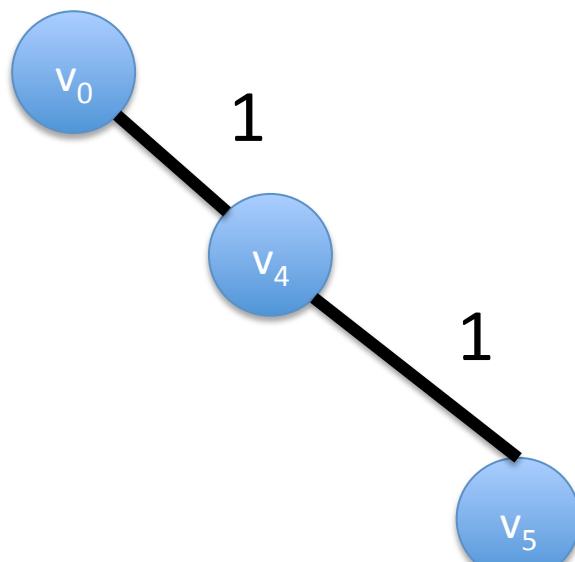
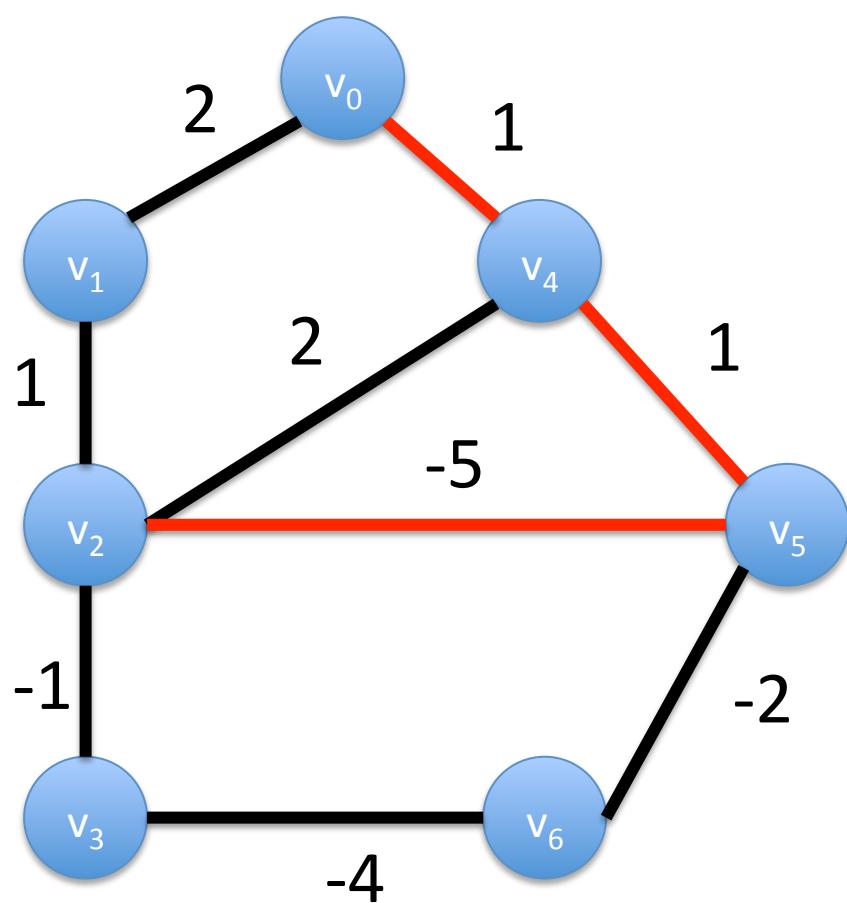
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Prim's Method



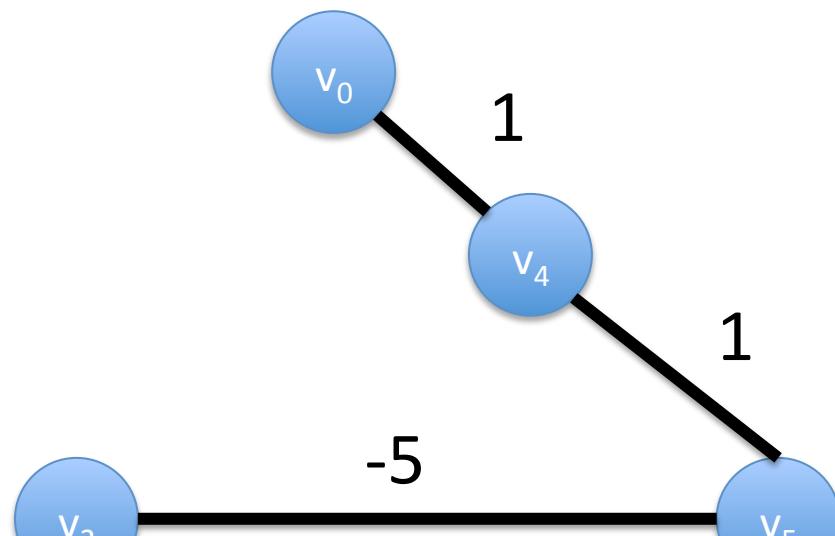
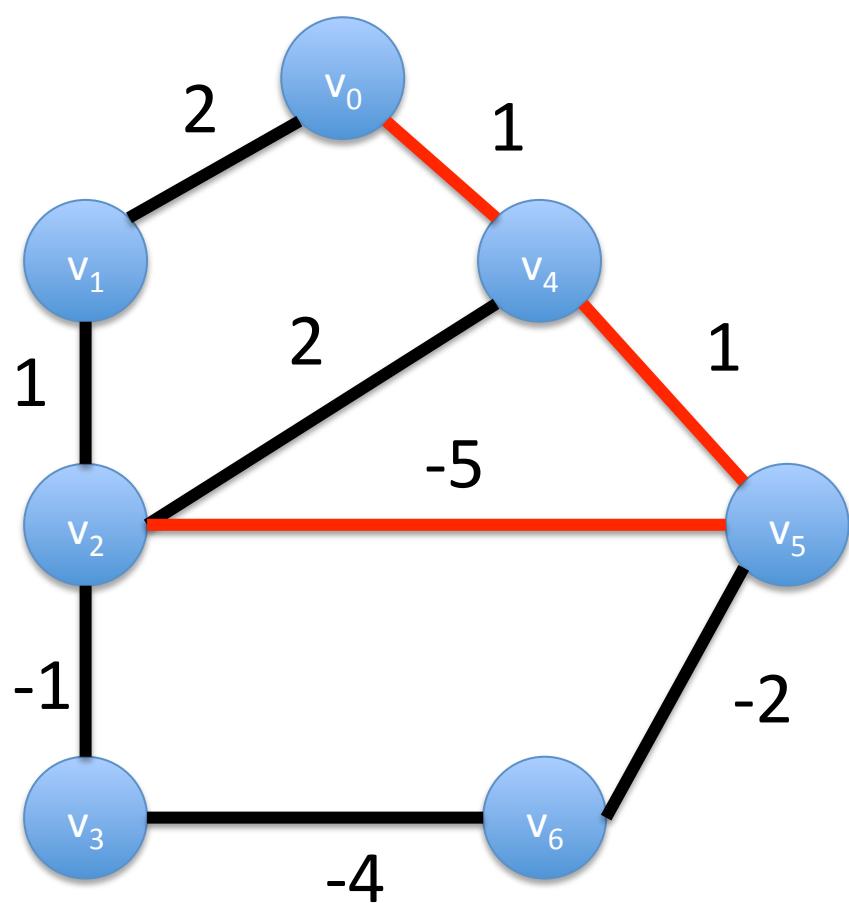
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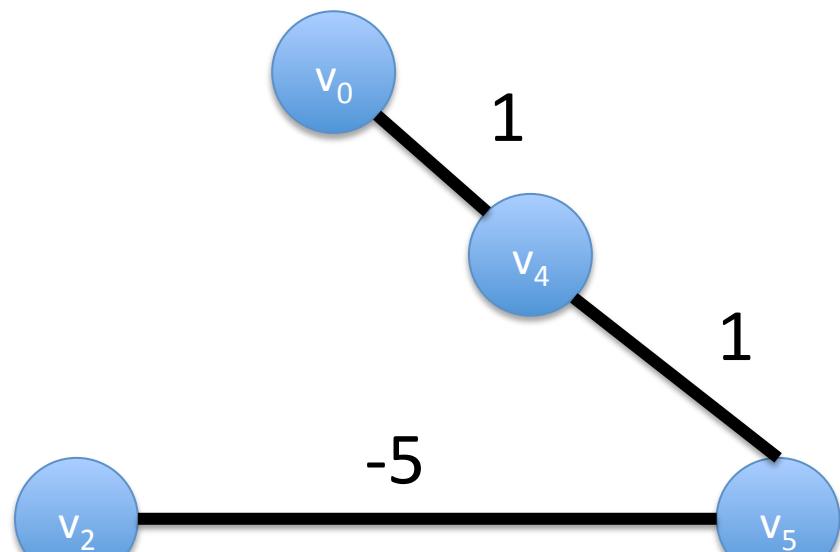
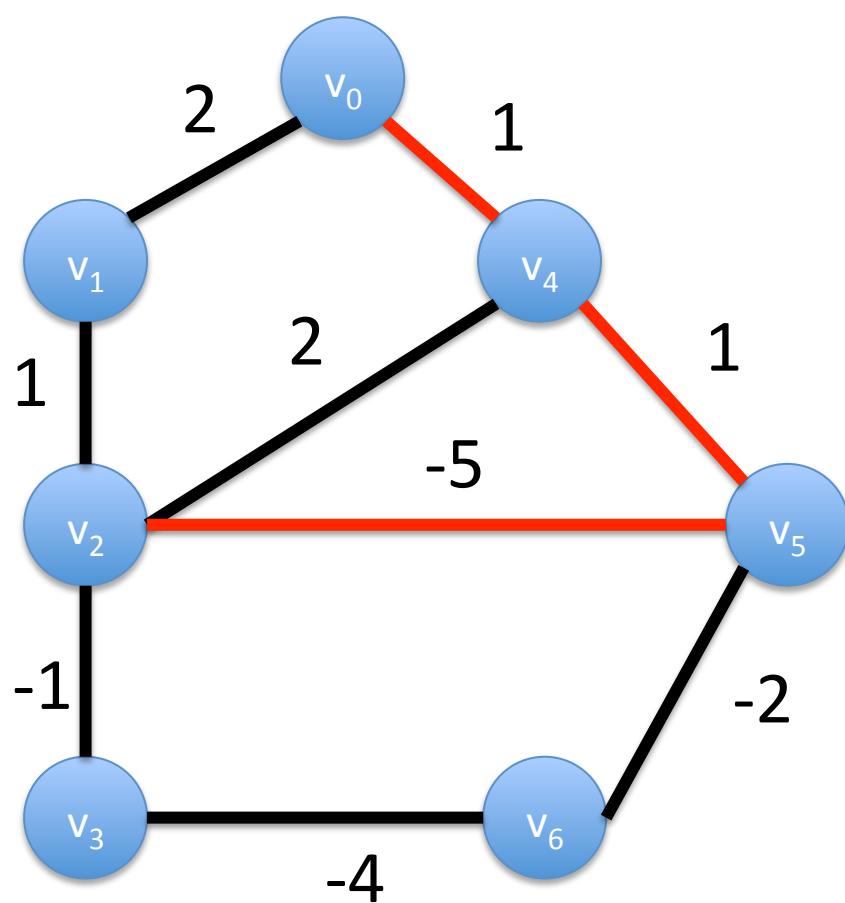
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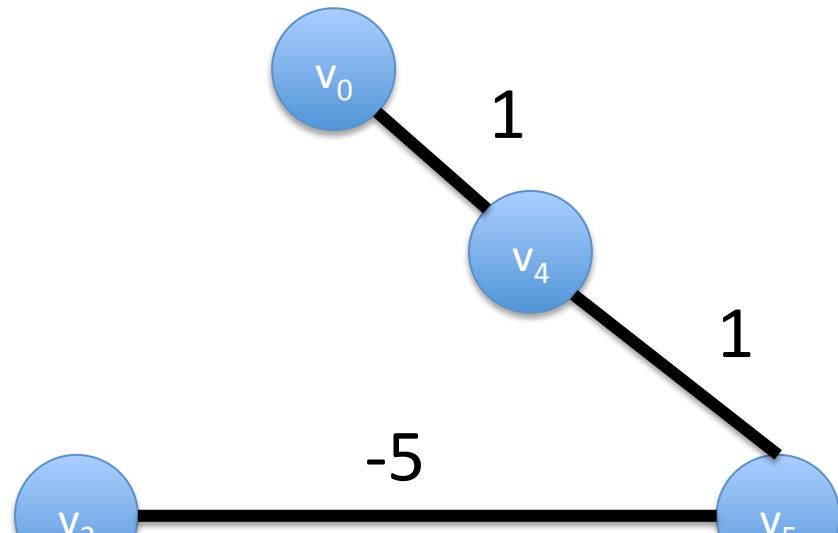
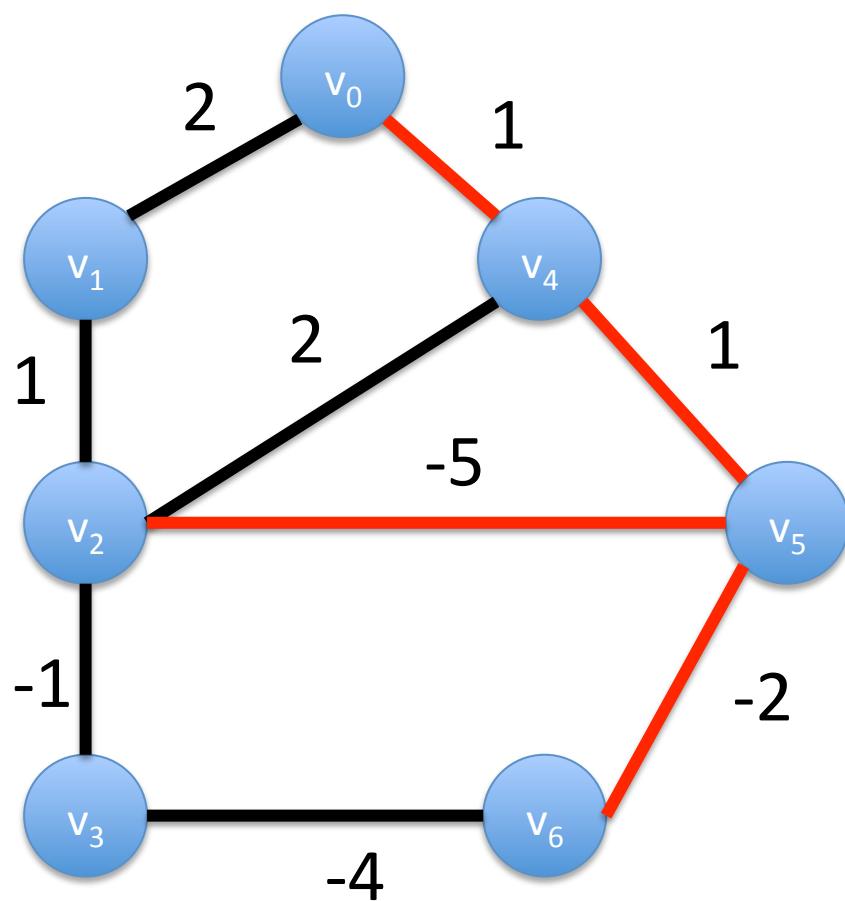
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Prim's Method



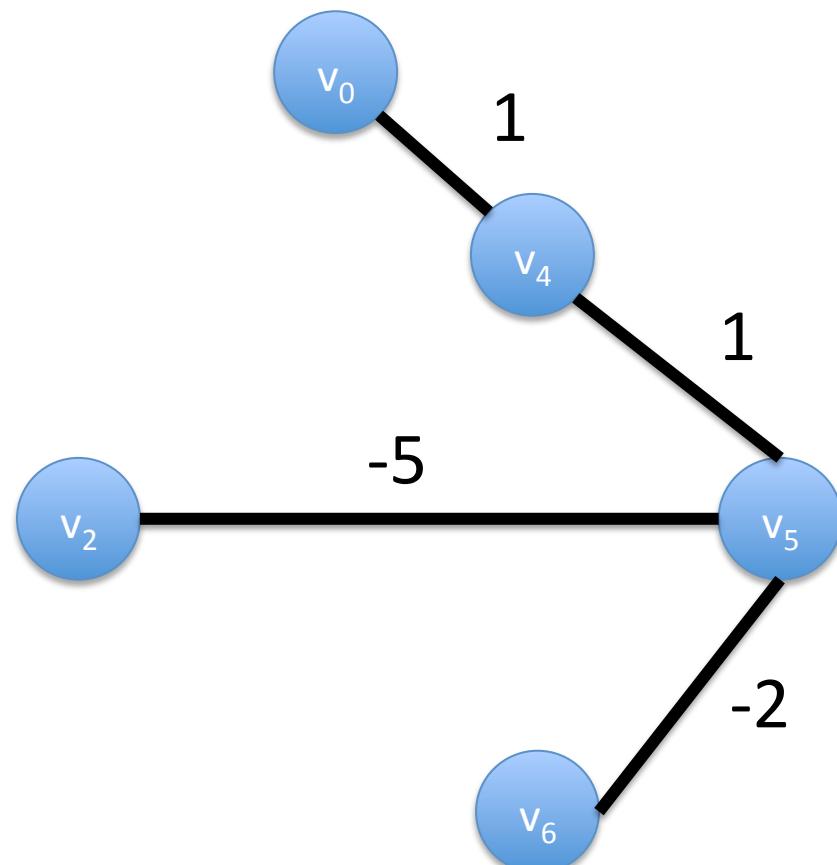
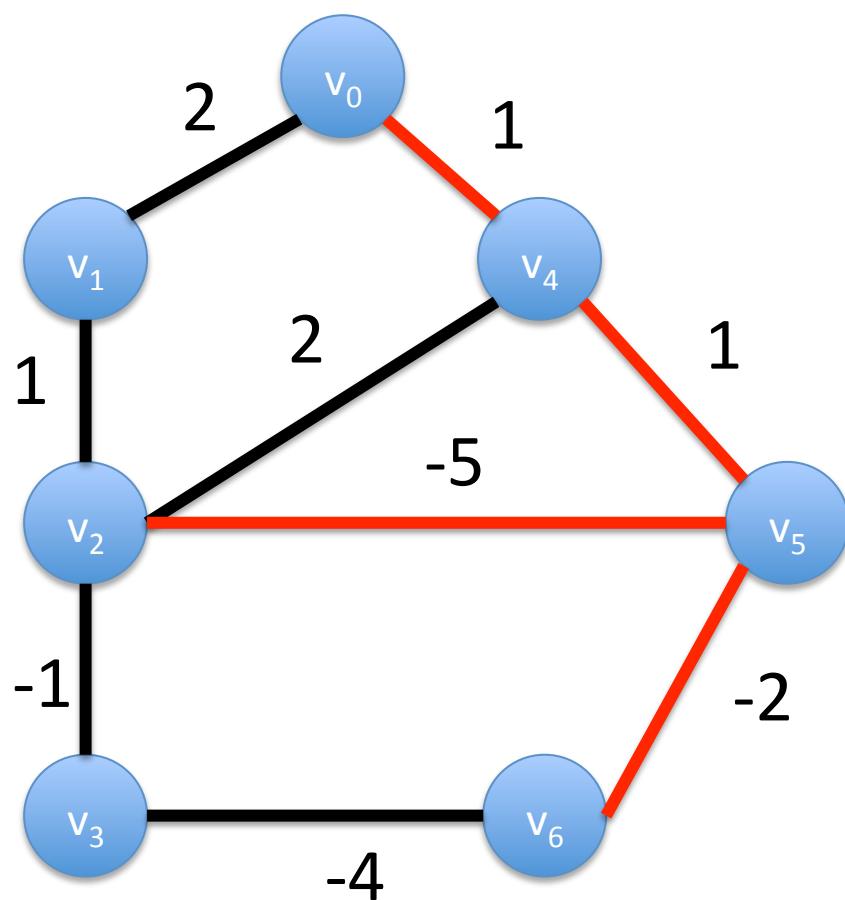
Choose $(u^*, v^*) = \min I(u, v), u \in V_T, v \in V \setminus V_T$

Prim's Method



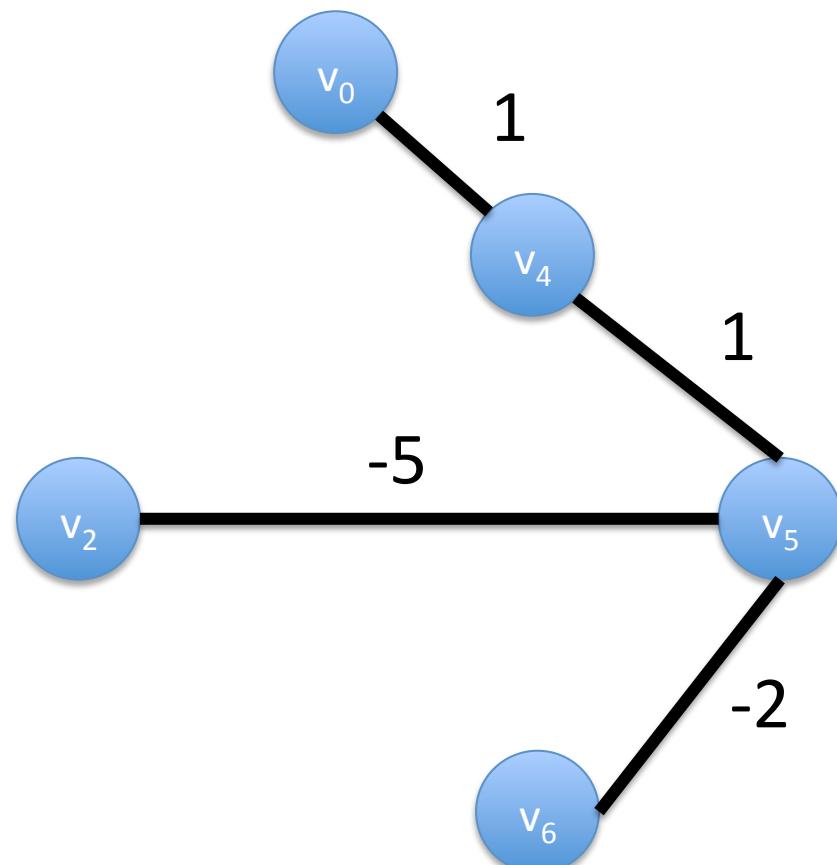
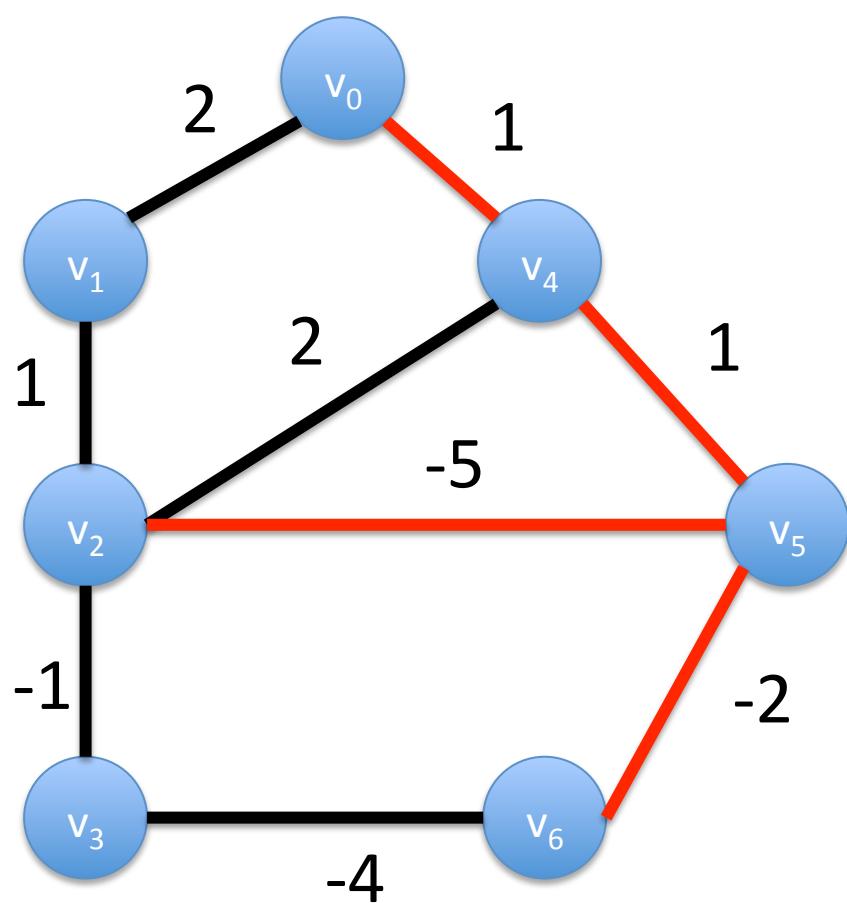
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Prim's Method



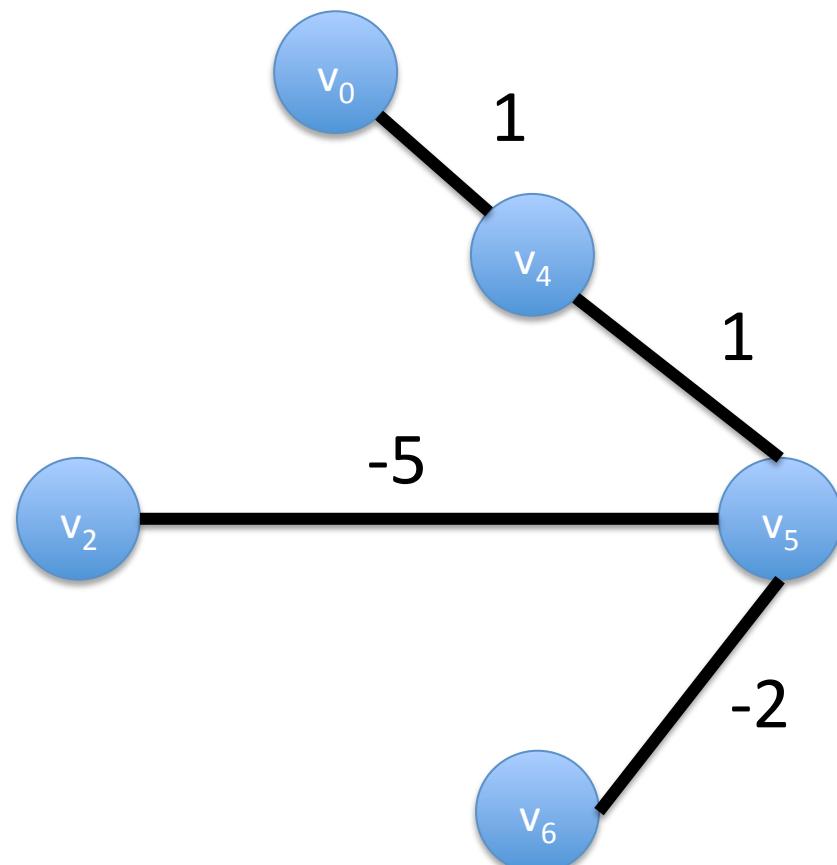
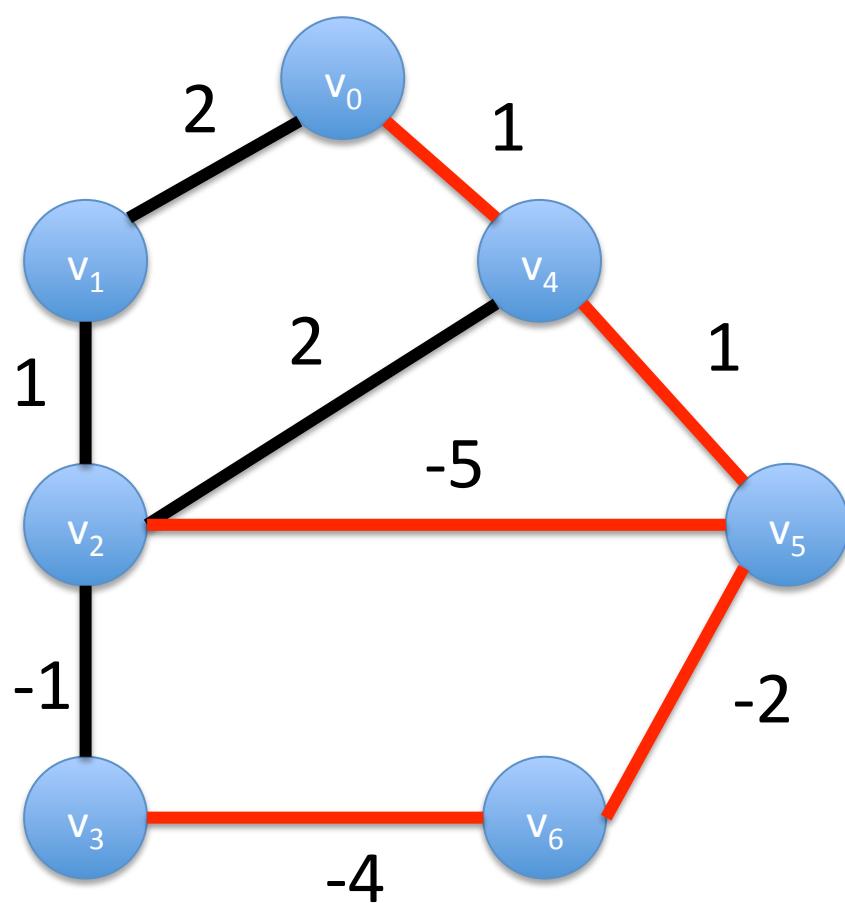
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Prim's Method



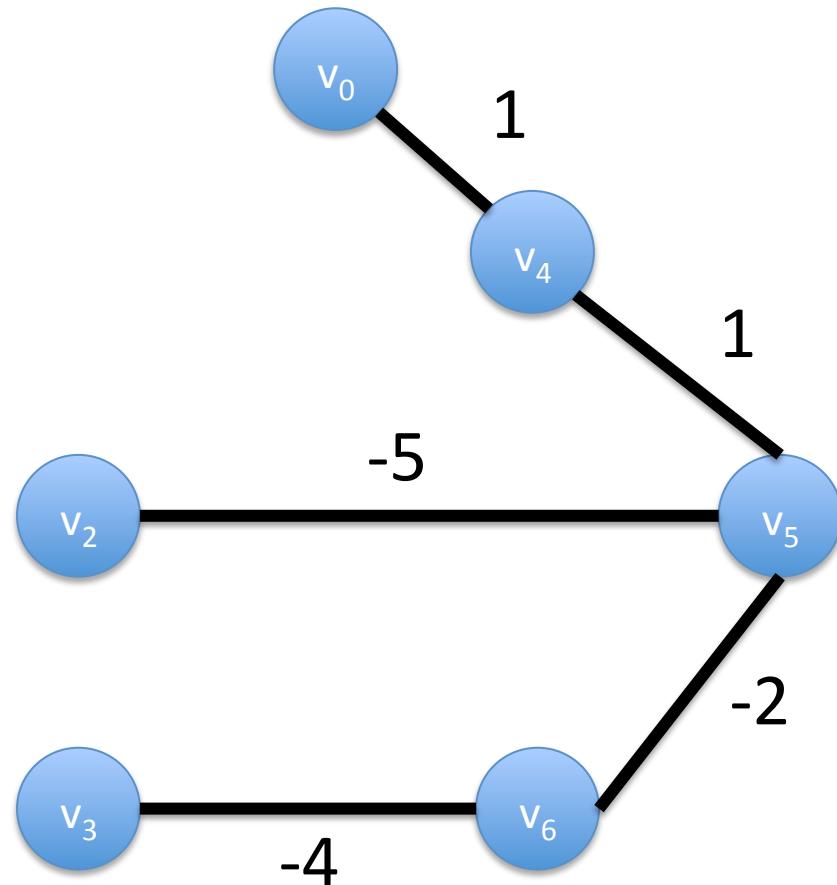
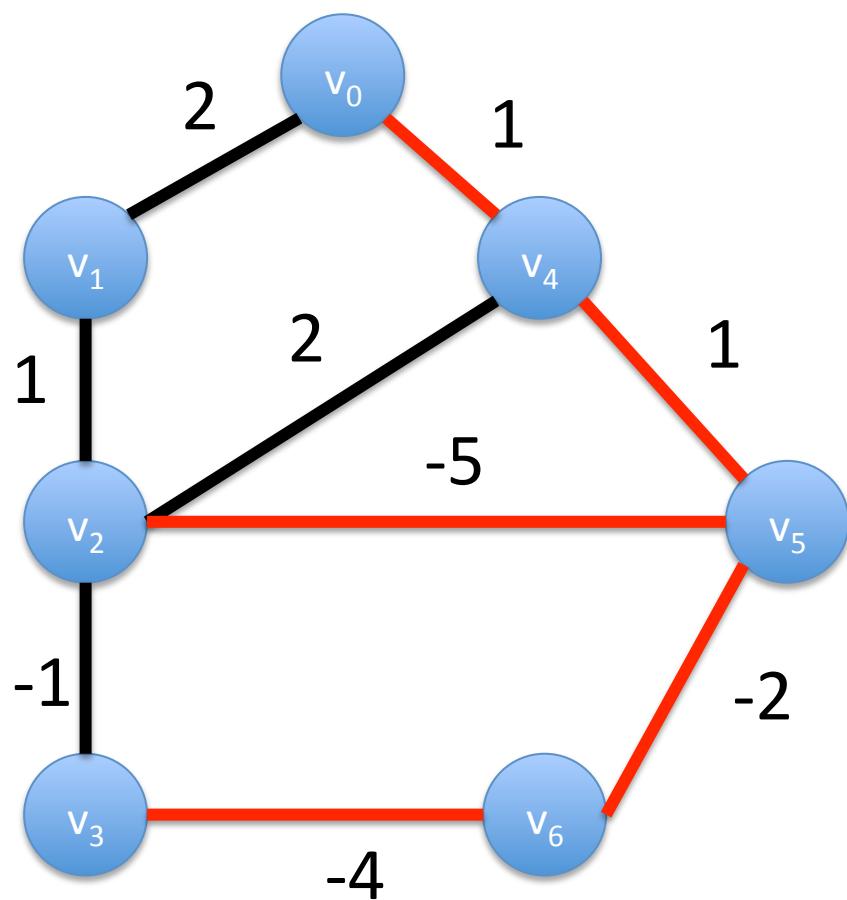
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Prim's Method



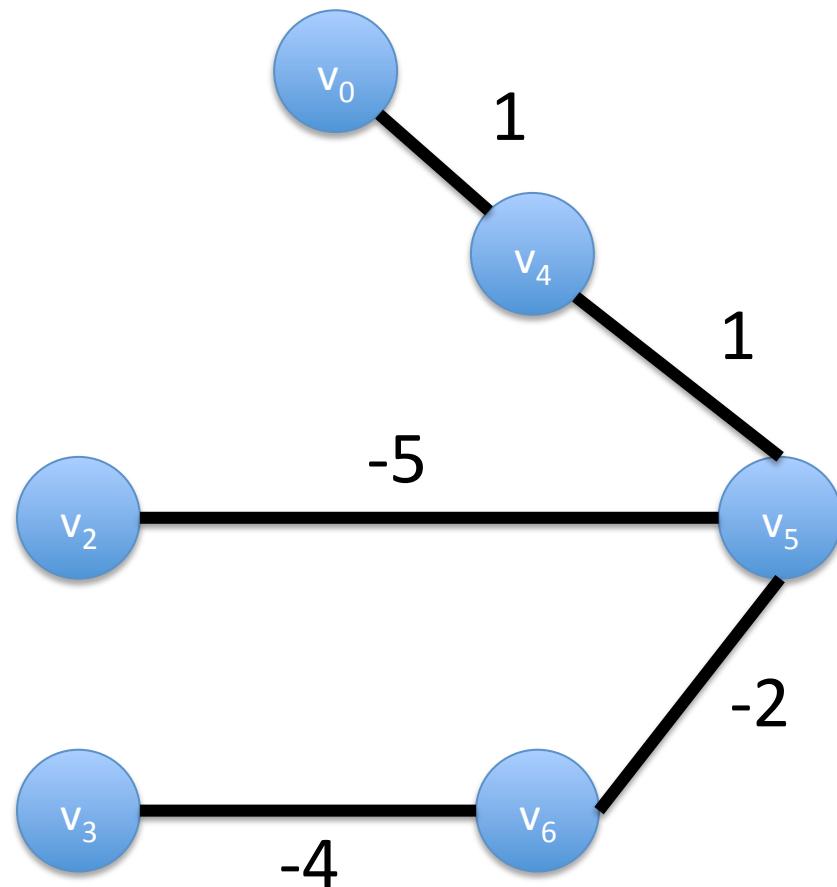
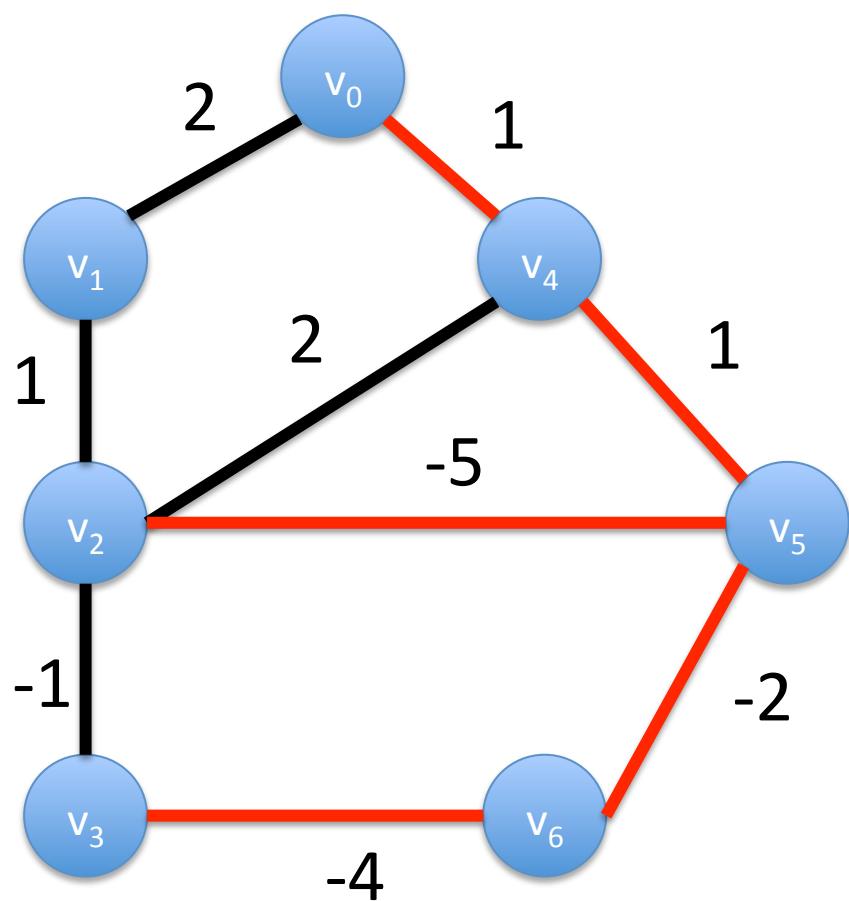
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Prim's Method



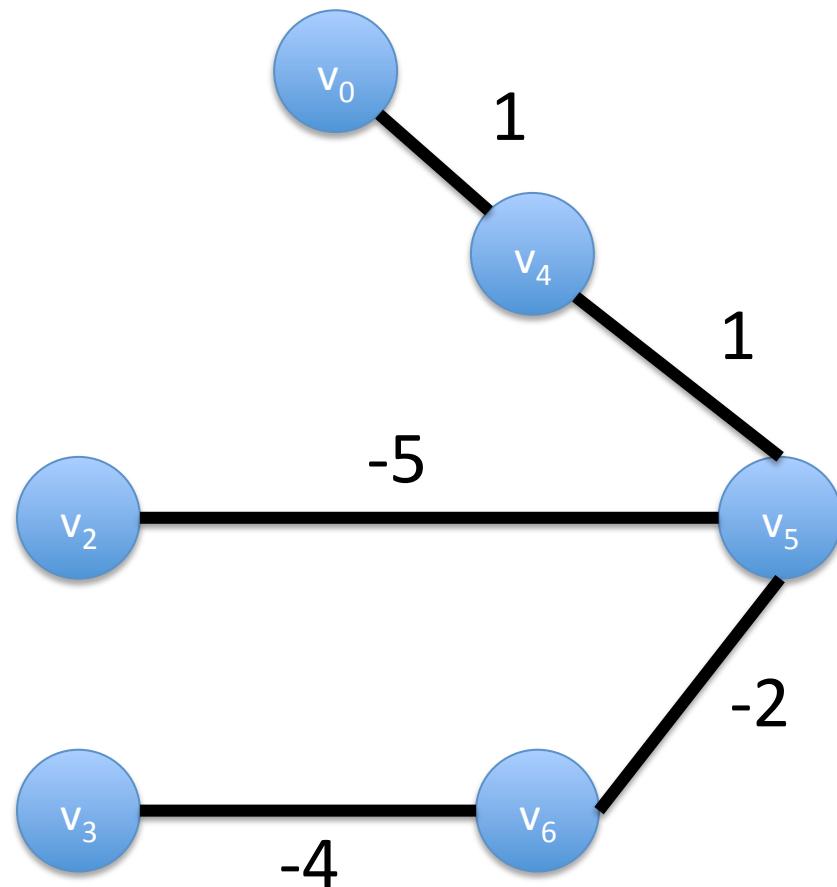
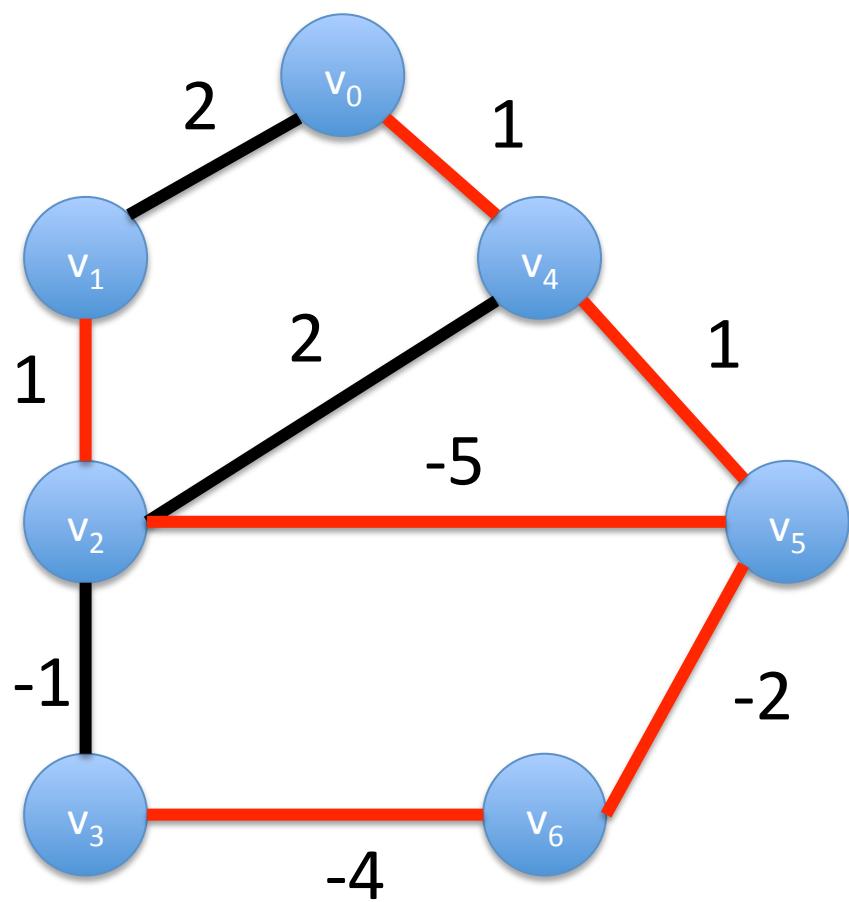
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Prim's Method



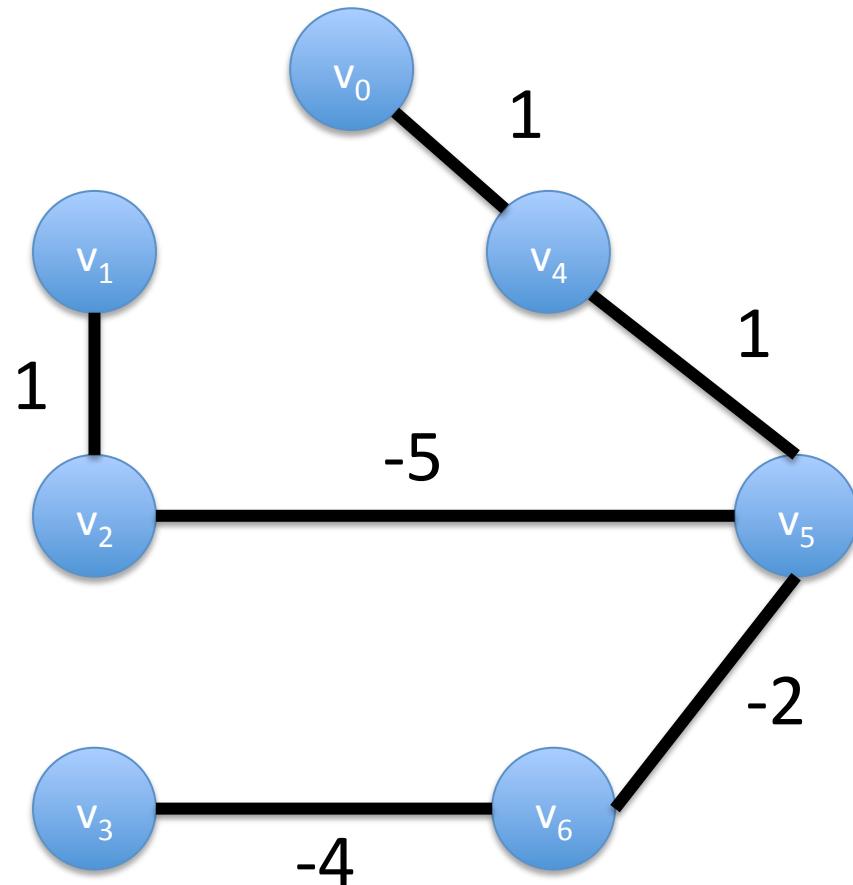
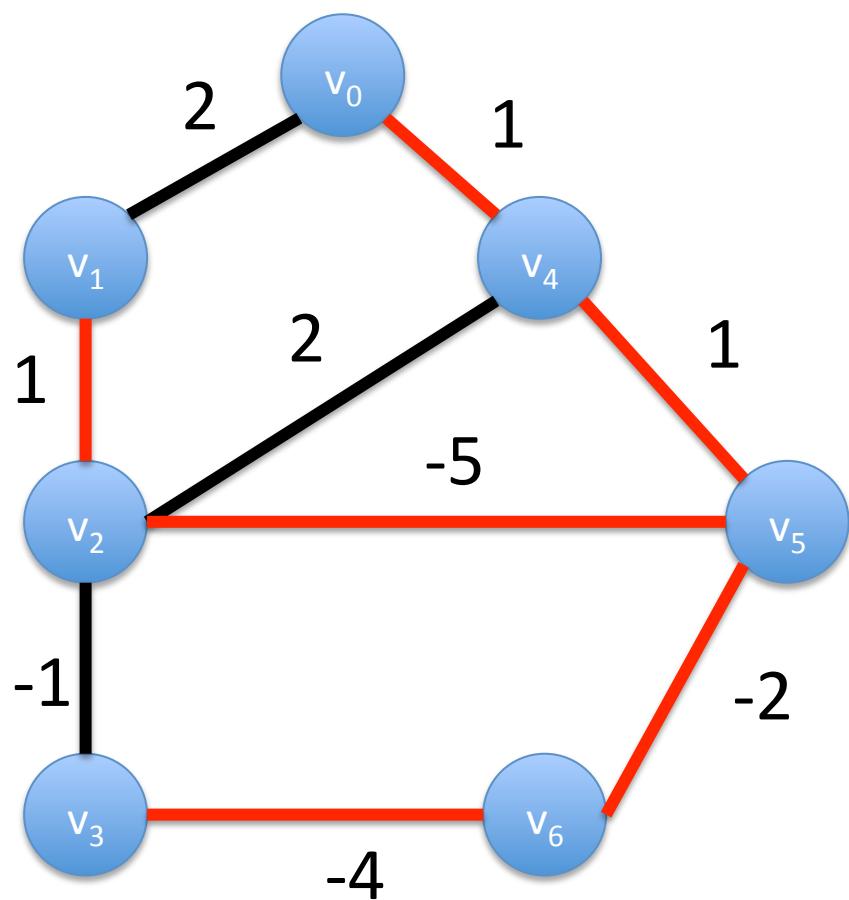
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Prim's Method



Choose $(u^*, v^*) = \min I(u, v), u \in V_T, v \in V \setminus V_T$

Prim's Method



Add v^* to V_T . Add (u^*, v^*) to E_T .

Summary

Given $G = (V, E)$. Define $T = (V_T, \{\}) = (\{v_0\}, \{\})$.

While V_T is not equal to V

Select $(u, v) = \min I(u, v)$ where

(1) u belongs to V_T

(2) v does not belong to V_T

Add v to V_T . Add (u, v) to E_T

End

Time Complexity

$O(n^2)$ where $n = |V|$

Proof of Prim's Algorithm

Homework!