Discrete Optimization

MA2827 Fondements de l'optimisation discrète

Dynamic programming (Part 2)

https://project.inria.fr/2015ma2827/

Material based on the lectures of Erik Demaine at MIT and Pascal Van Hentenryck at Coursera

Outline

- Dynamic programming
 - Guitar fingering
- More dynamic programming
 - Tetris
 - Blackjack
- Quiz: bracket sequences

Dynamic programming

- DP ≈ "careful brute force"
- DP \approx recursion + memoization + guessing
- Divide the problem into subproblems that are connected to the original problem
- Graph of subproblems has to be acyclic (DAG)
- Time = #subproblems · time/subproblem

5 easy steps of DP

Analysis:

- 1. Define subproblems #subproblems
- 2. Guess part of solution #choices
- 3. Relate subproblems (recursion) time/subproblem
- 4. Recurse + memoize time
 OR build DP table bottom-up
 check subprobs be acyclic / topological order
- 5. Solve original problem

extra time

Task: find the best way to play a melody



Task: find the best way to play a melody

- Input: sequence of notes to play with right hand
- One note at a time!
- Which finger to use? 1, 2, ..., F = 5 for humans
- Measure d(f, p, g, q) of difficulty to go from note p with finger f to note q with finger g

Examples of rules:

crossing fingers: 1 < f < g and p > q => uncomfortable

stretching: p << q

=> uncomfortable

legato (smooth): ∞ if f = g

Task: find the best way to play a melody Goal: minimize overall difficulty

```
Subproblems:
min. difficulty for suffix note[ i : ]
#subproblems = O( n ) where n = #notes
```

Guesses: finger f for the first note[i] #choices = F

```
Recurrence:
DP[i] = min{ DP[i+1]+d(note[i], f, note[i+1], next finger ) }
```

Task: find the best way to play a melody Goal: minimize overall difficulty

Subproblems: min. difficulty for suffix note[i :] #subproblems = O(n) where n = #notes

Guesses: finger f for the first note[i] #choices = F

Recurrence: DP[i] = min{ DP[i+1] + d(note[i], f, note[i+1], next finger) } Not enough information!

Task: find the best way to play a melody Goal: minimize overall difficulty

Subproblems: min. difficulty for suffix note[i :] when finger f is on note[i] #subproblems = O(n F)

Guesses: finger f for the next note, note[i + 1] #choices = F

Recurrence: DP[i, f] = min{ DP[i + 1, g] + d(note[i], f, note[i+1], g) | all g } Base-case: DP[n, f] = 0 time/subproblem = O(F)

Task: find the best way to play a melody



Final problem:
 find minimal DP[0, f] for f = 1, ..., F
 guessing the first finger

Task: win in the game of Tetris!



Task: win in the game of Tetris!

- Input: a sequence of n Tetris pieces and an empty board of small width w
- Choose orientation and position for each piece
- Must drop piece till it hits something
- Full rows do not clear
- Goal: survive i.e., stay within height h

Task: stay within height h

```
Subproblem:
survival? in suffix [ i : ]
given a particular column profile
#subproblems = O( n (h+1)<sup>w</sup> )
```



Guesses: where to drop piece i? #choices = O(w)

```
Recurrence:
DP[i, p] = max { DP[i + 1, q] | q is a valid move from p }
Base-case: DP[n+1, p] = true for all profiles p
time/subproblem = O(w)
```

Task: stay within height h



Final problem: DP[0, empty]

Blackjack

Task: beat the blackjack (twenty-one)!



Blackjack

Task: beat the blackjack!

Rules of Blackjack (simplified):

- The player and the dealer are initially given 2 cards each
- Each card gives points:
 - Cards 2-10 are valued at the face value of the card
 - Face cards (King, Queen, Jack) are valued at 10
 - The Ace card can be valued either at 11 or 1
- The goal of the player is to get more points than the dealer, but less than 21, if more than 21 than he looses (busts)
- Player can take any number of cards (hits)
- After that the dealer hits deterministically: until \geq 17 points

Perfect-information Blackjack

Task: beat the blackjack with a marked deck!

- Input: a deck of cards c₀, ..., c_{n-1}
- Player vs. dealer one-on-one
- Goal: maximize winning for a fixed bet \$1
- Might benefit from loosing to get a better deck

Quiz (as homework)

- Write the DP for perfect-information blackjack
- Derive the number of subproblems for the tetris problem