

Discrete Optimization

MA2827

Fondements de l'optimisation discrète

<https://project.inria.fr/2015ma2827/>

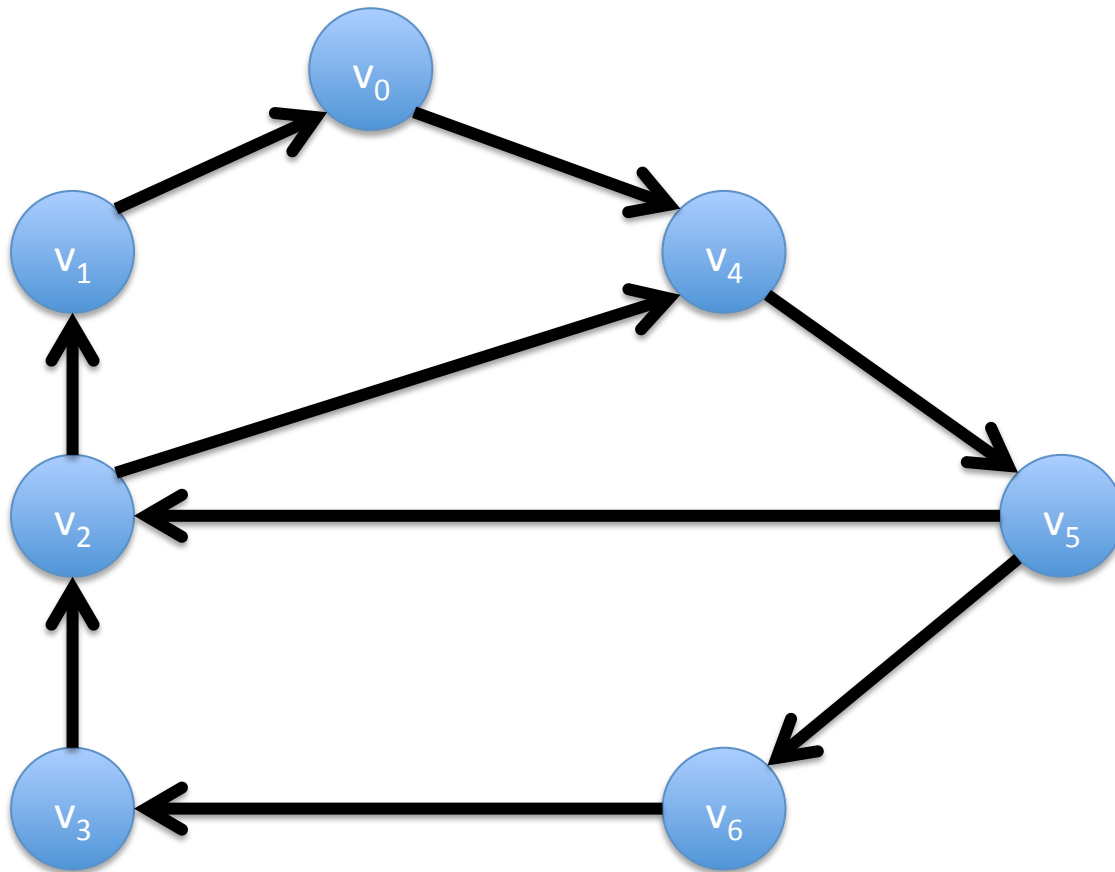
Slides courtesy of M. Pawan Kumar

Outline

- Preliminaries (Recap)
- Menger's Theorem for Disjoint Paths
- Path Packing

Directed Graphs (Digraphs)

$$D = (V, A)$$

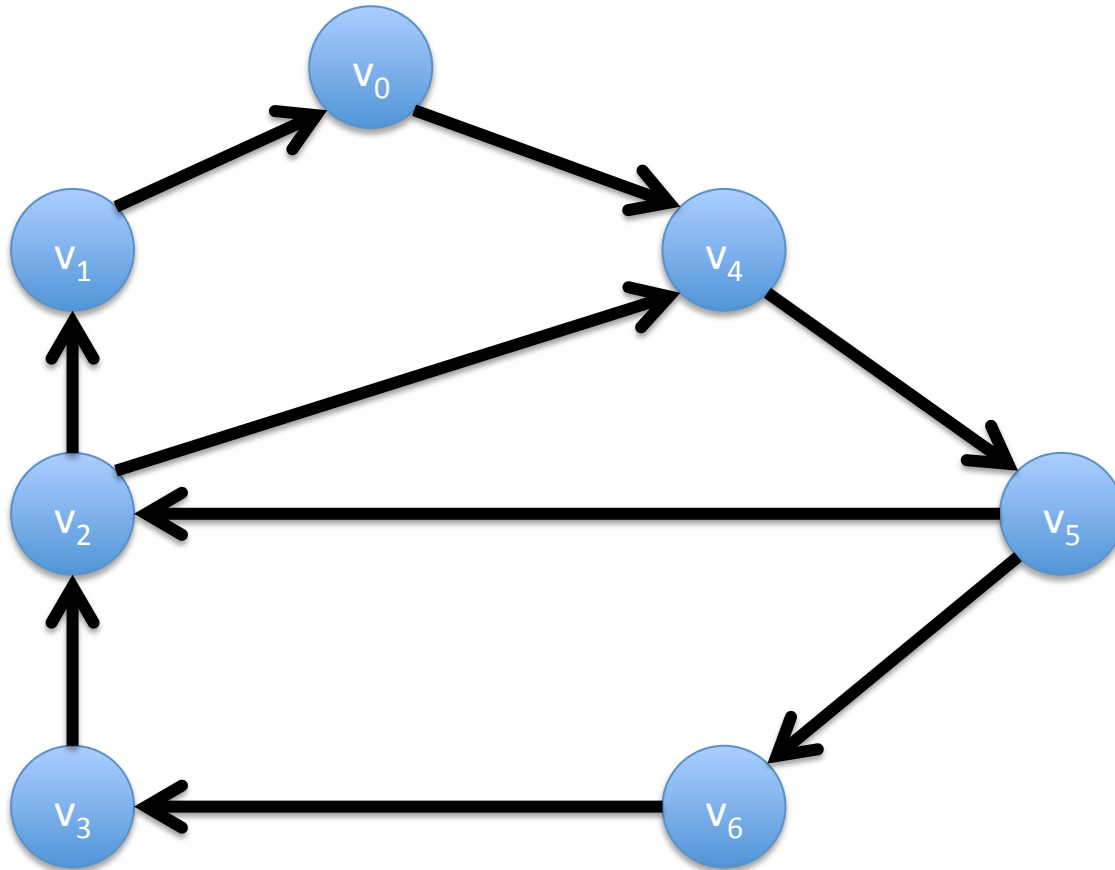


'n' vertices or nodes V

'm' arcs A : ordered pairs from V

Indegree of a Vertex

$$D = (V, A)$$

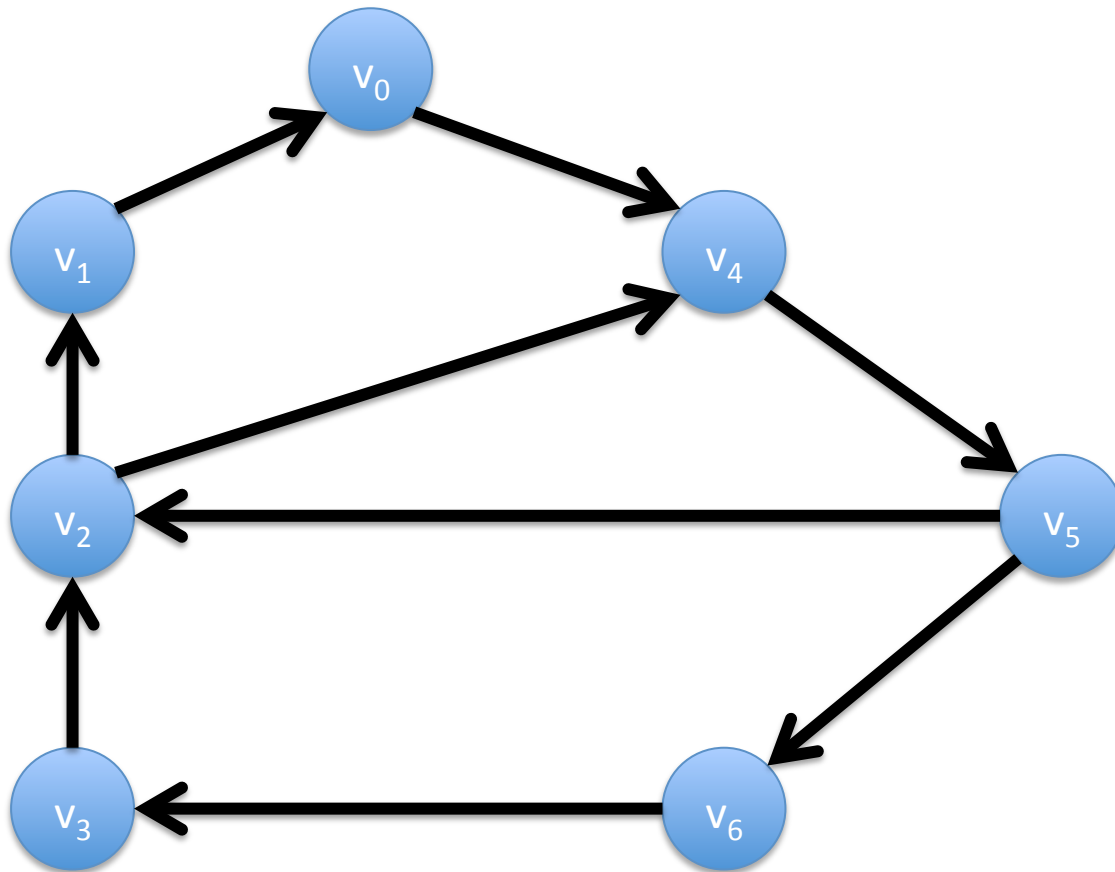


Number of arcs entering the vertex.

$$\text{indeg}(v_0) = 1, \text{indeg}(v_1) = 1, \text{indeg}(v_4) = 2, \dots$$

Indegree of a Subset of Vertices

$$D = (V, A)$$

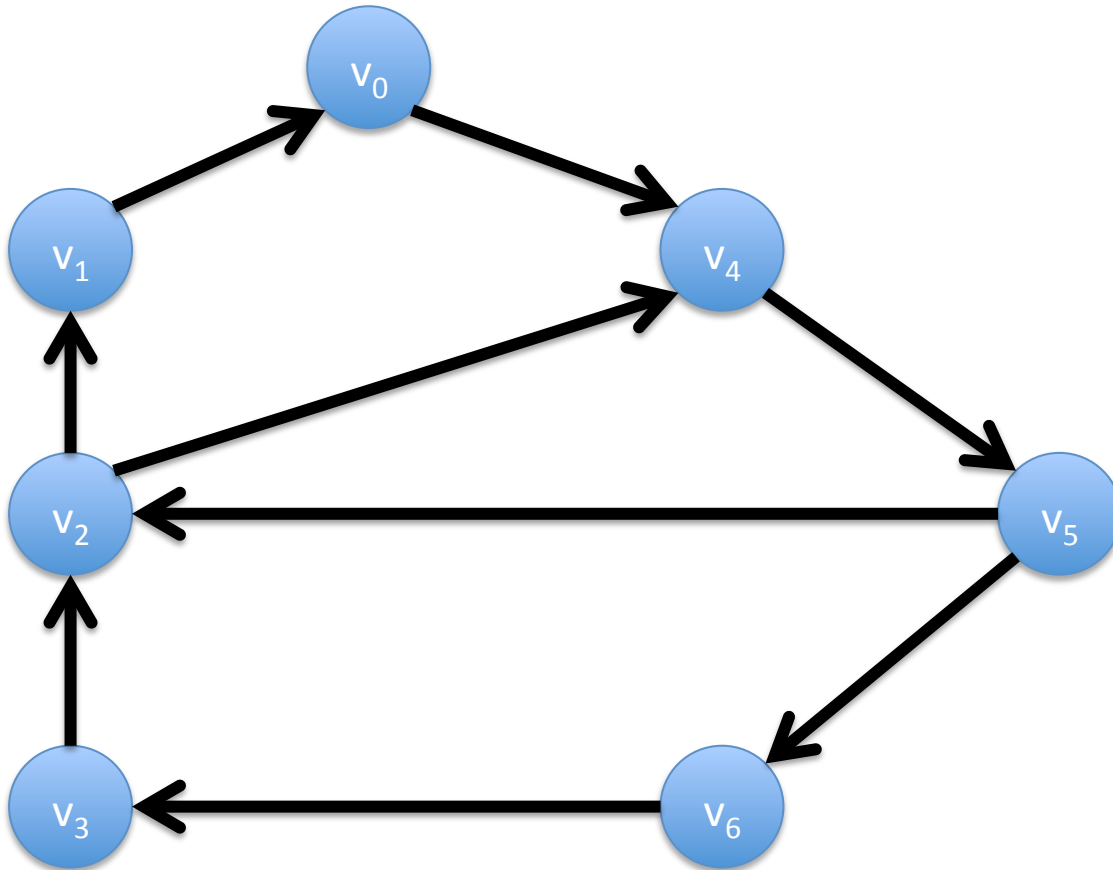


Number of arcs entering the subset.

$$\text{indeg}(\{v_0, v_1\}) = 1, \text{indeg}(\{v_1, v_4\}) = 3, \dots$$

Outdegree of a Vertex

$$D = (V, A)$$

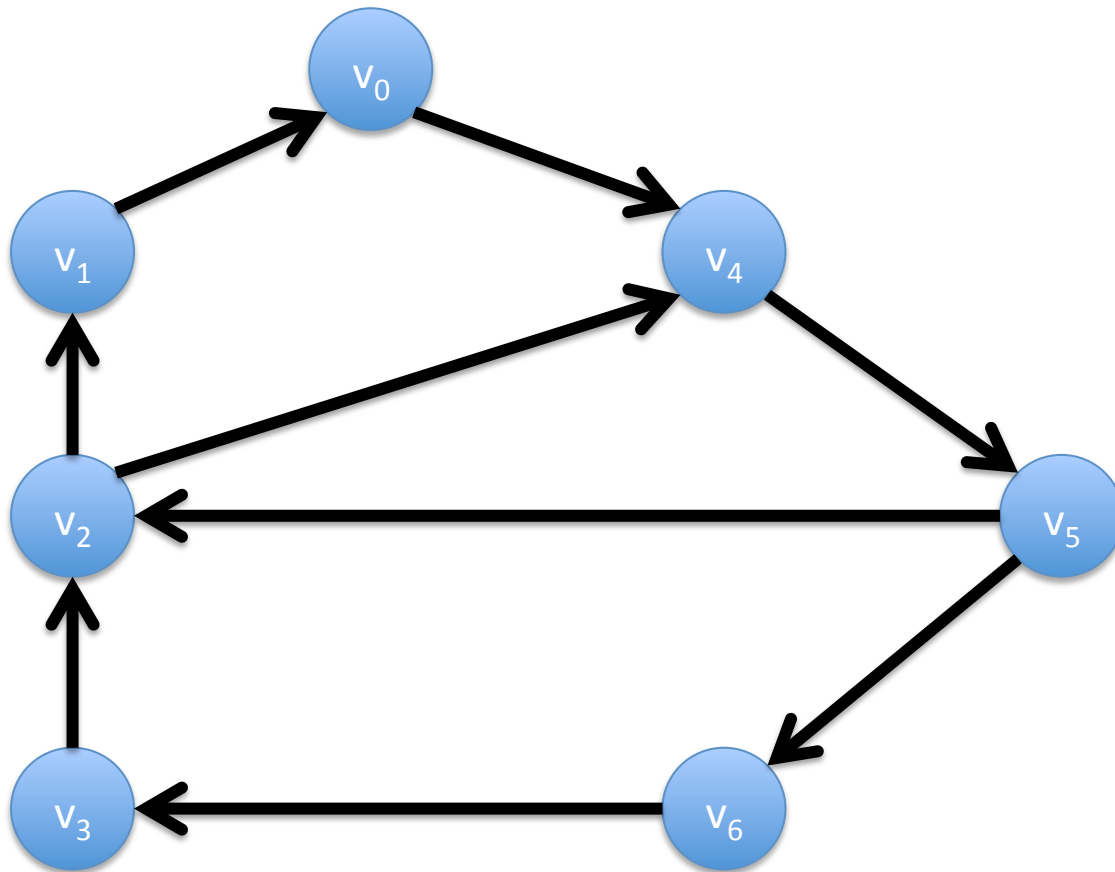


Number of arcs leaving the vertex.

$$\text{outdeg}(v_0) = 1, \text{outdeg}(v_1) = 1, \text{outdeg}(v_2) = 2, \dots$$

Outdegree of a Subset of Vertices

$$D = (V, A)$$

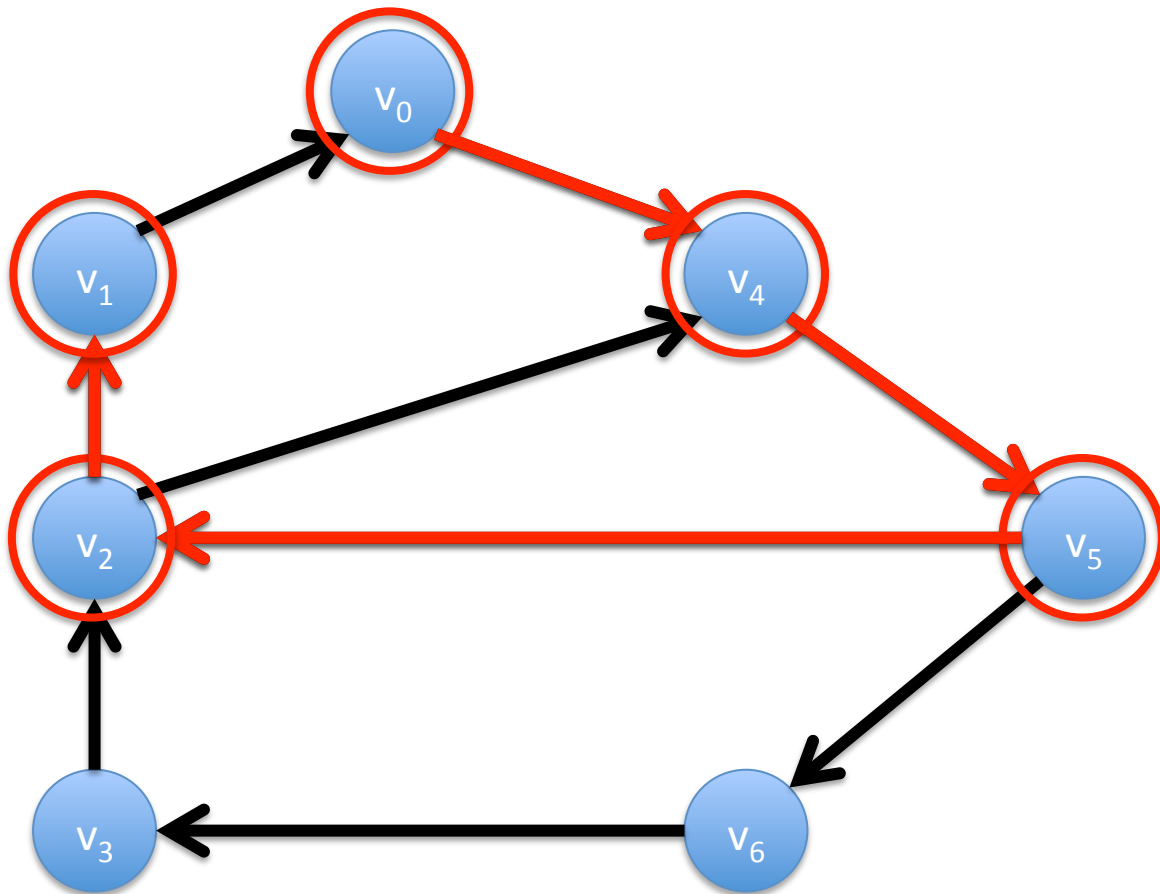


Number of arcs leaving the subset.

$$\text{outdeg}(\{v_0, v_1\}) = 1, \text{outdeg}(\{v_1, v_4\}) = 2, \dots$$

s-t Path

$$D = (V, A)$$

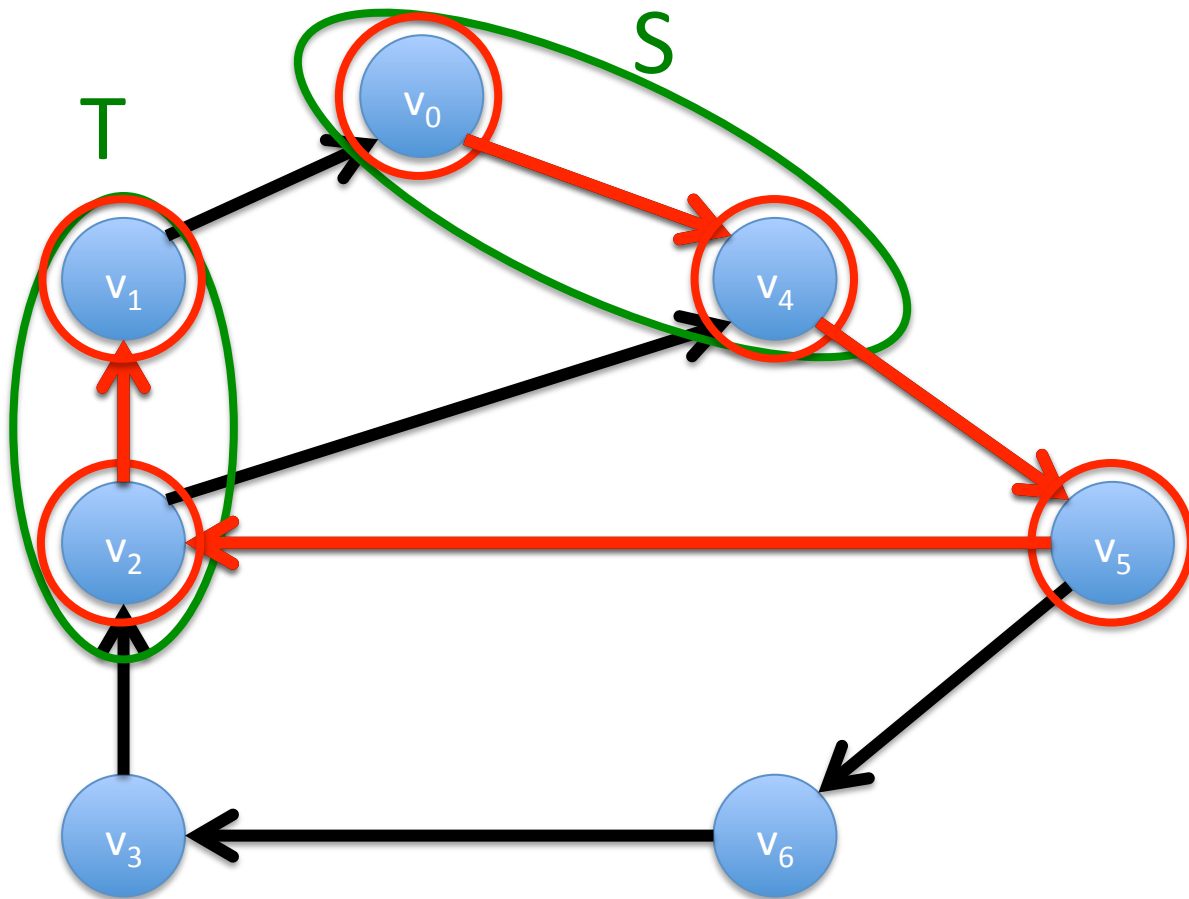


Sequence $P = (s=v_0, a_1, v_1, \dots, a_k, t=v_k)$, $a_i = (v_{i-1}, v_i)$

Vertices $s=v_0, v_1, \dots, t=v_k$ are distinct

S-T Path

$$D = (V, A)$$

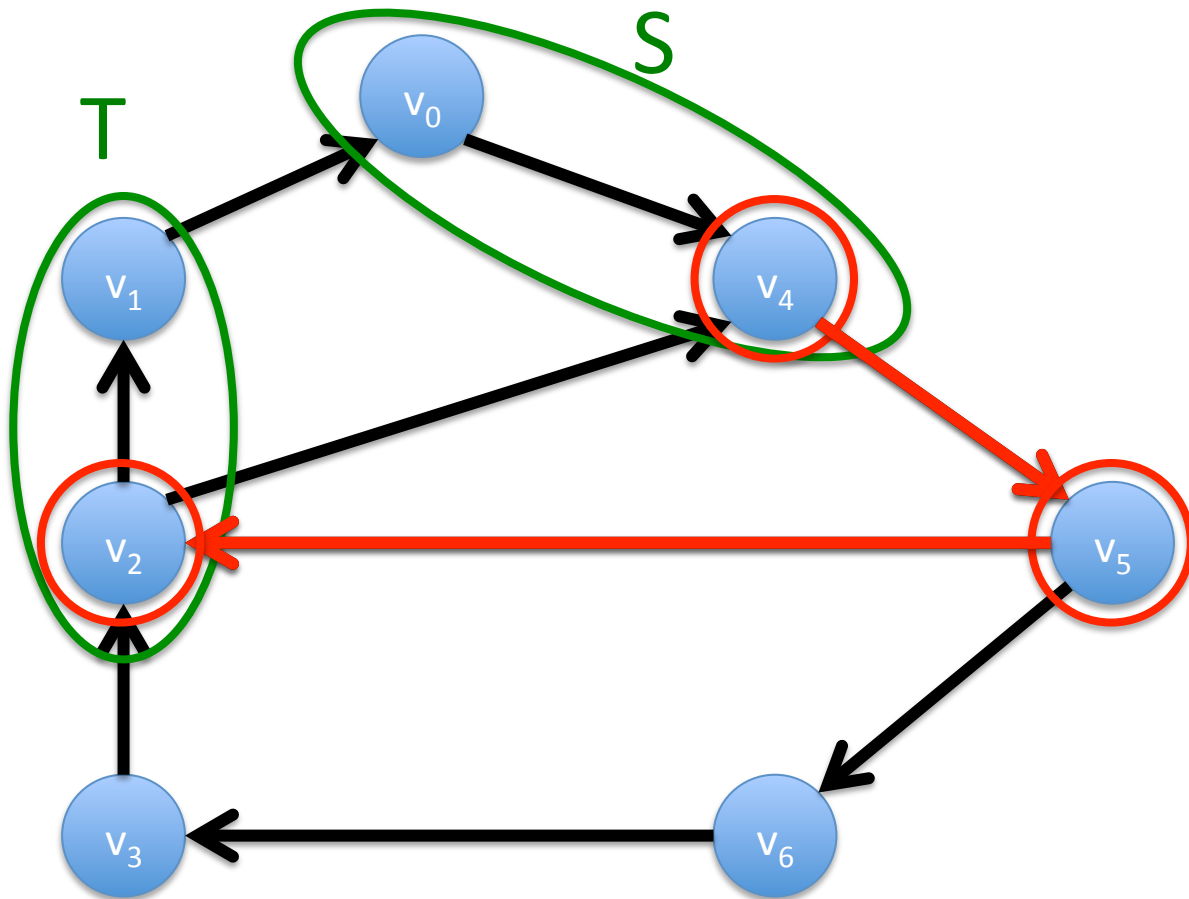


S and T are subsets of V

Any st -path where $s \in S$ and $t \in T$

S-T Path

$$D = (V, A)$$



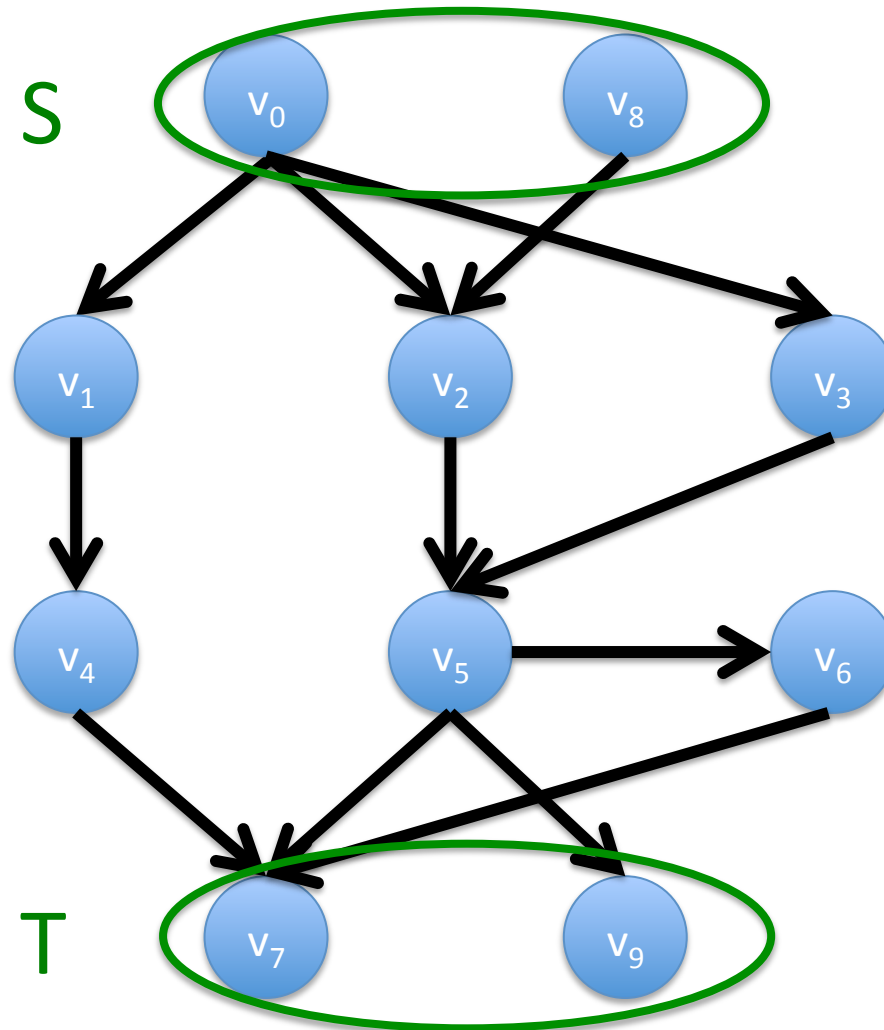
S and T are subsets of V

Any st-path where $s \in S$ and $t \in T$

Outline

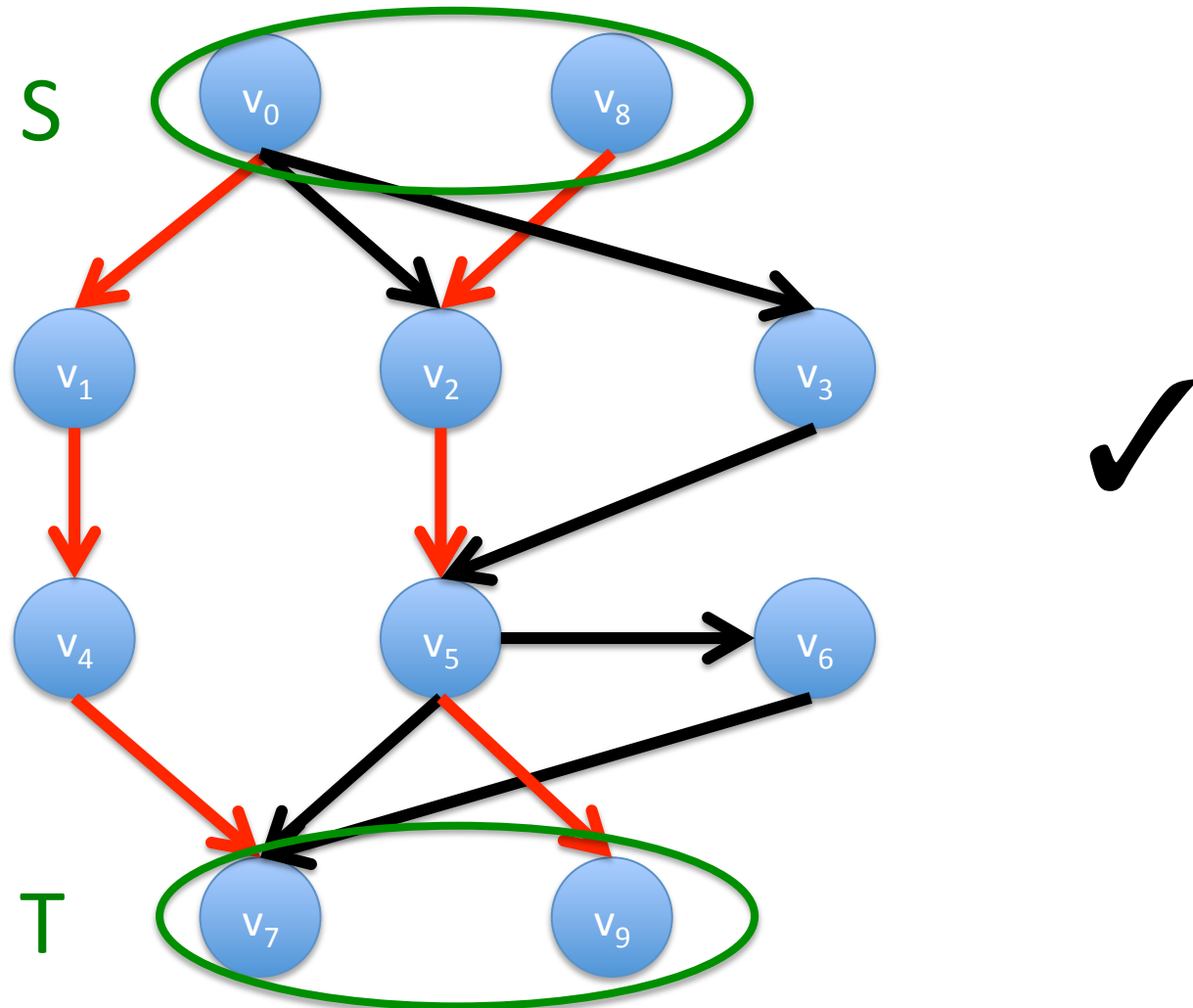
- Preliminaries
- **Menger's Theorem for Disjoint Paths**
- Path Packing

Vertex Disjoint S-T Paths



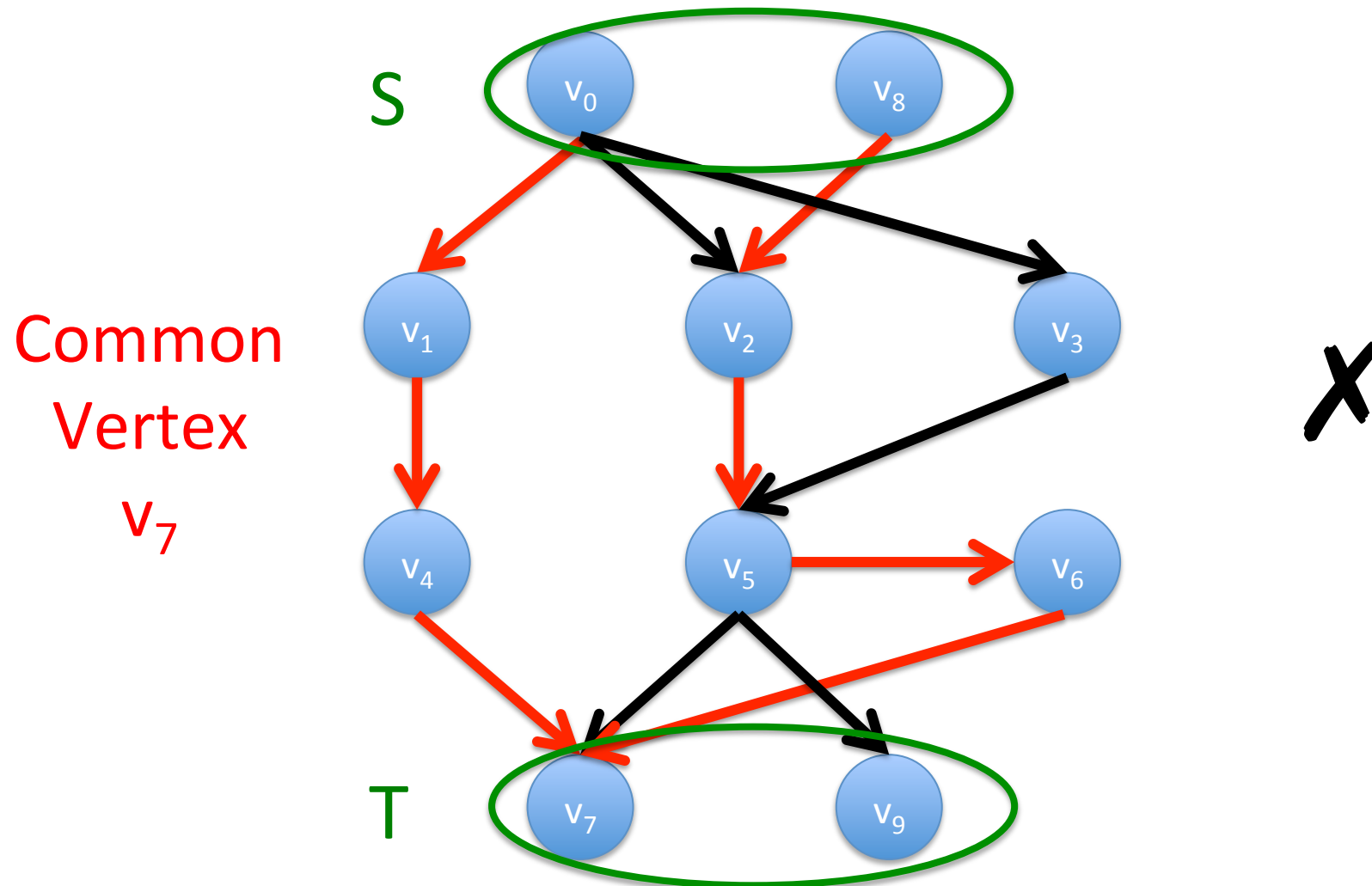
Set of S-T Paths with no common vertex

Vertex Disjoint S-T Paths



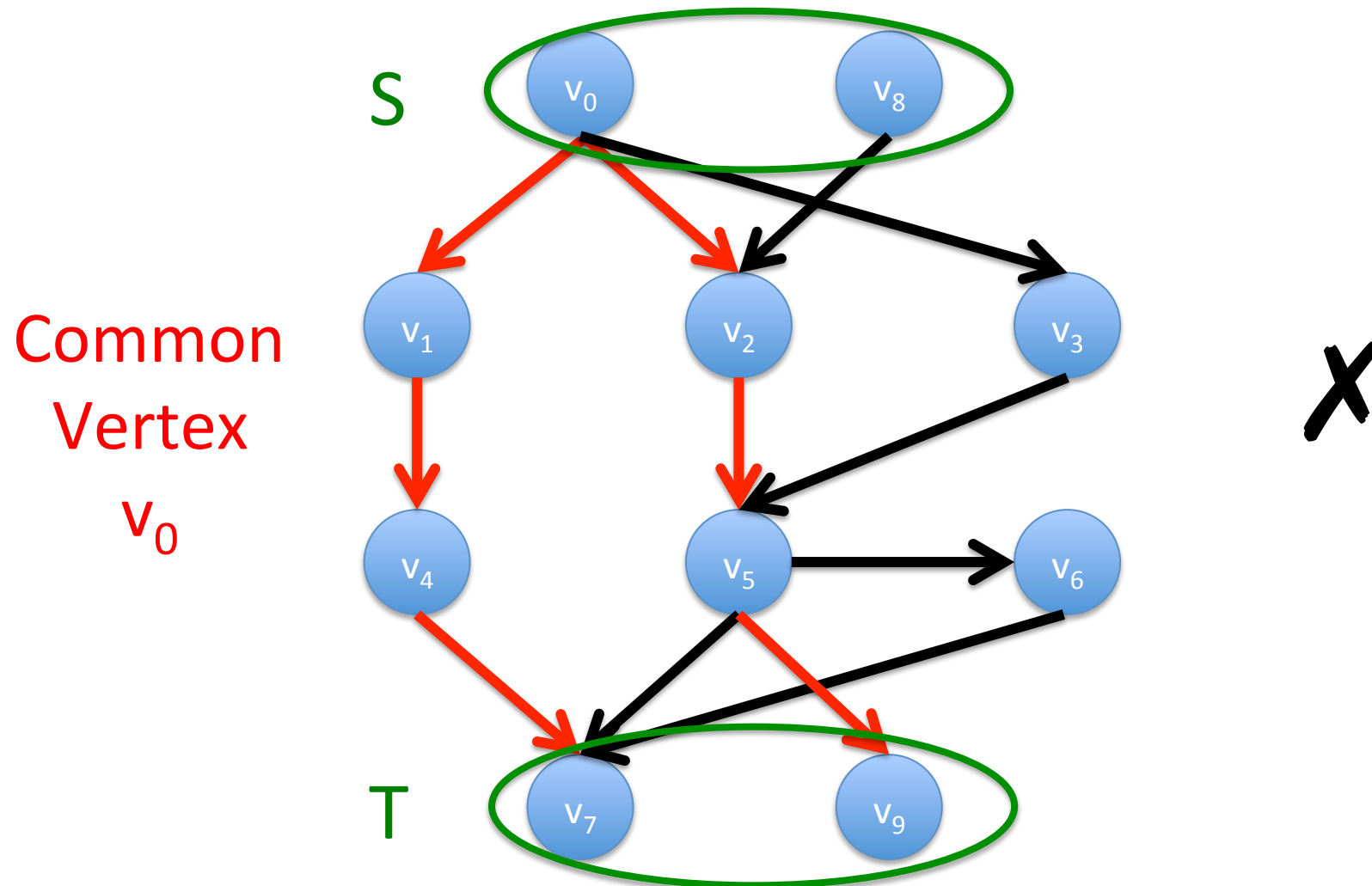
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Vertex Disjoint S-T Paths



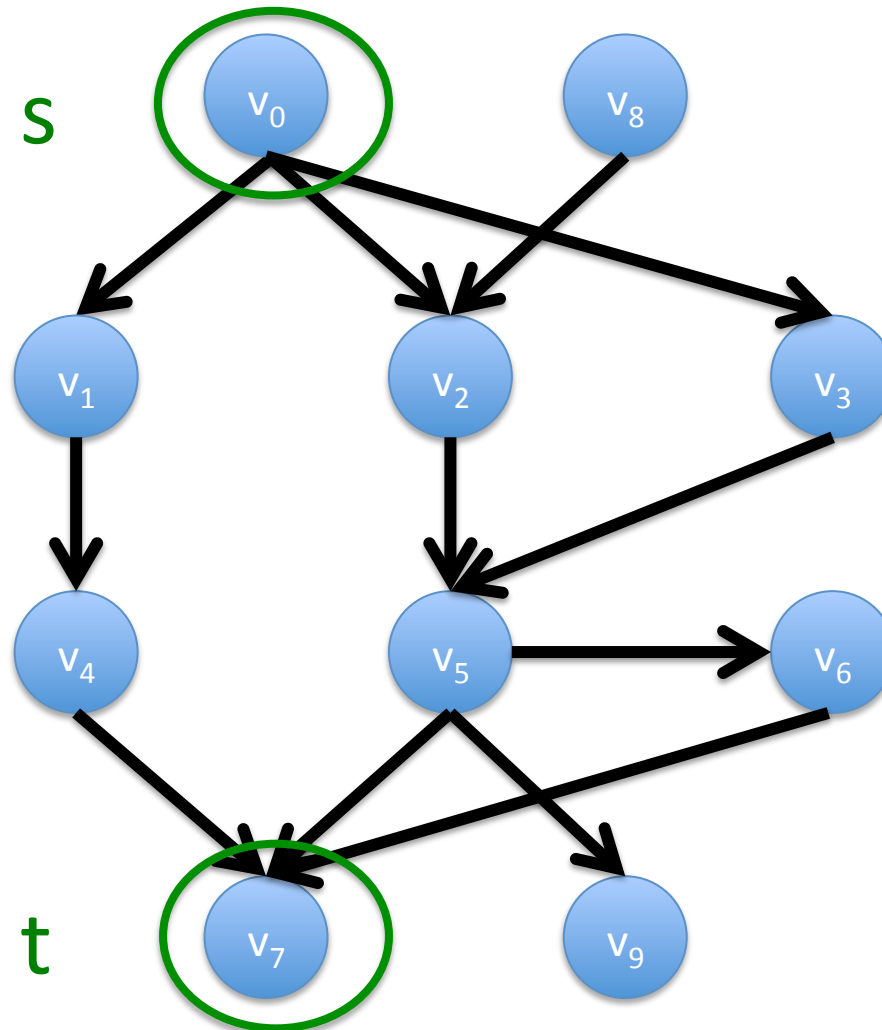
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Vertex Disjoint S-T Paths



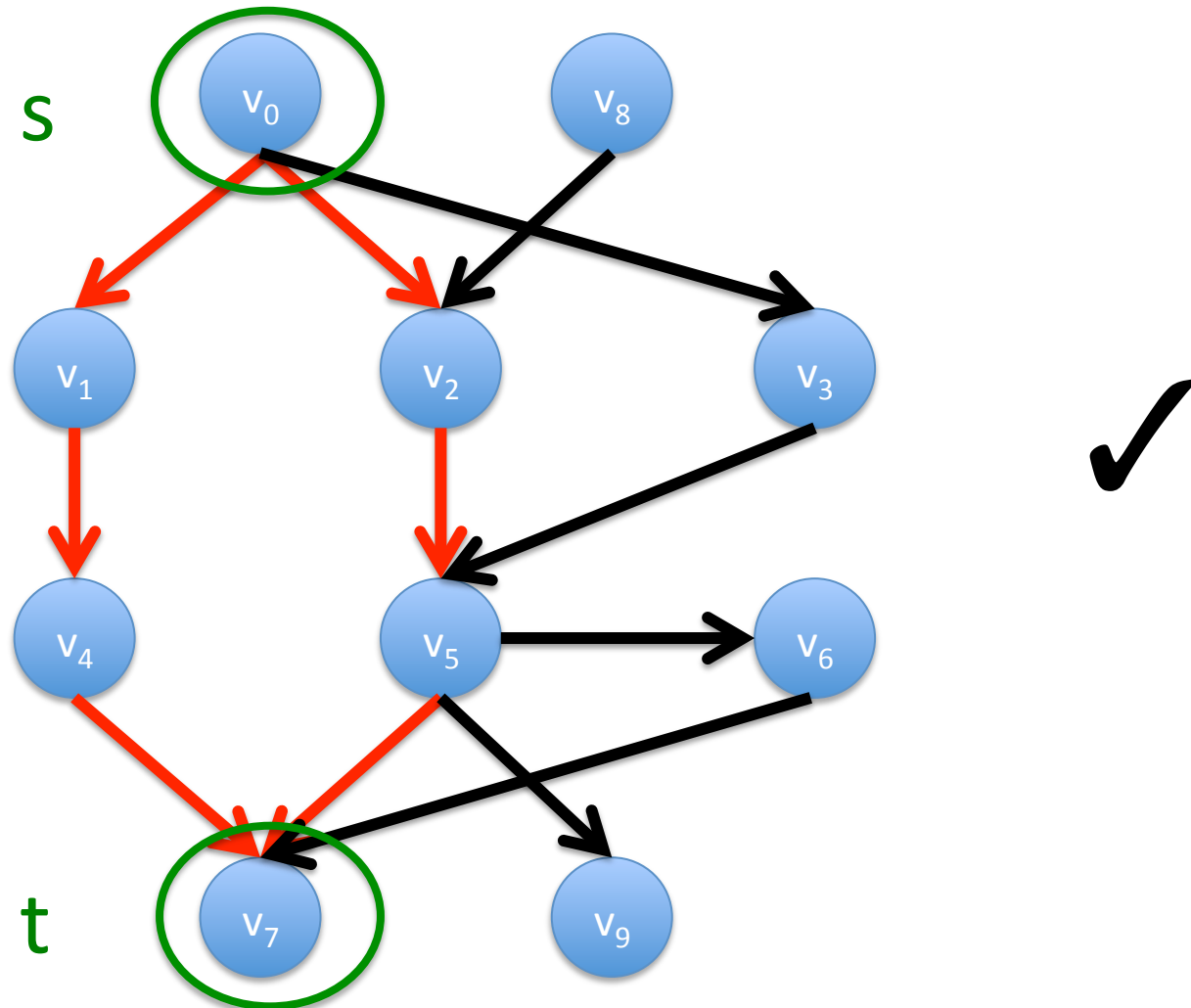
Set of S-T Paths with no common vertex

Internally Vertex Disjoint s-t Paths



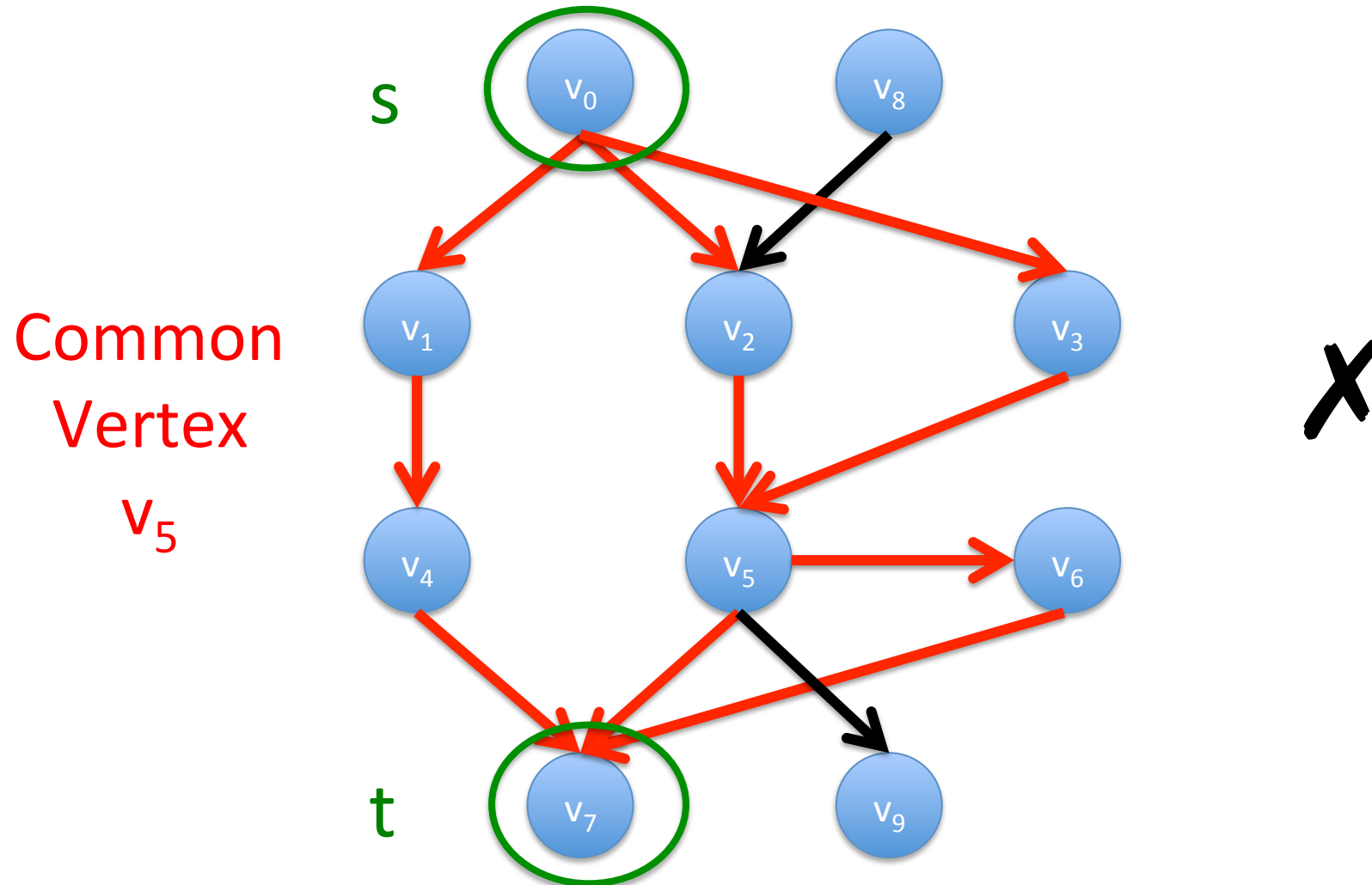
Set of s-t Paths with no common internal vertex

Internally Vertex Disjoint s-t Paths



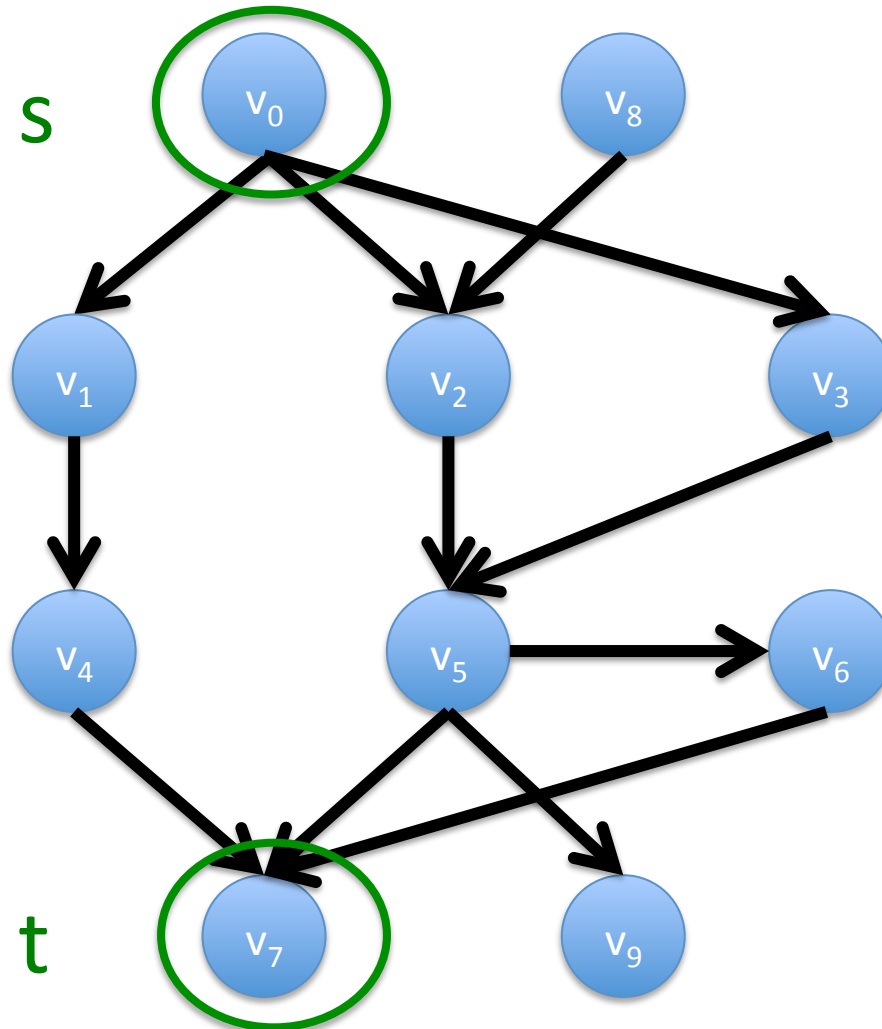
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Internally Vertex Disjoint s-t Paths



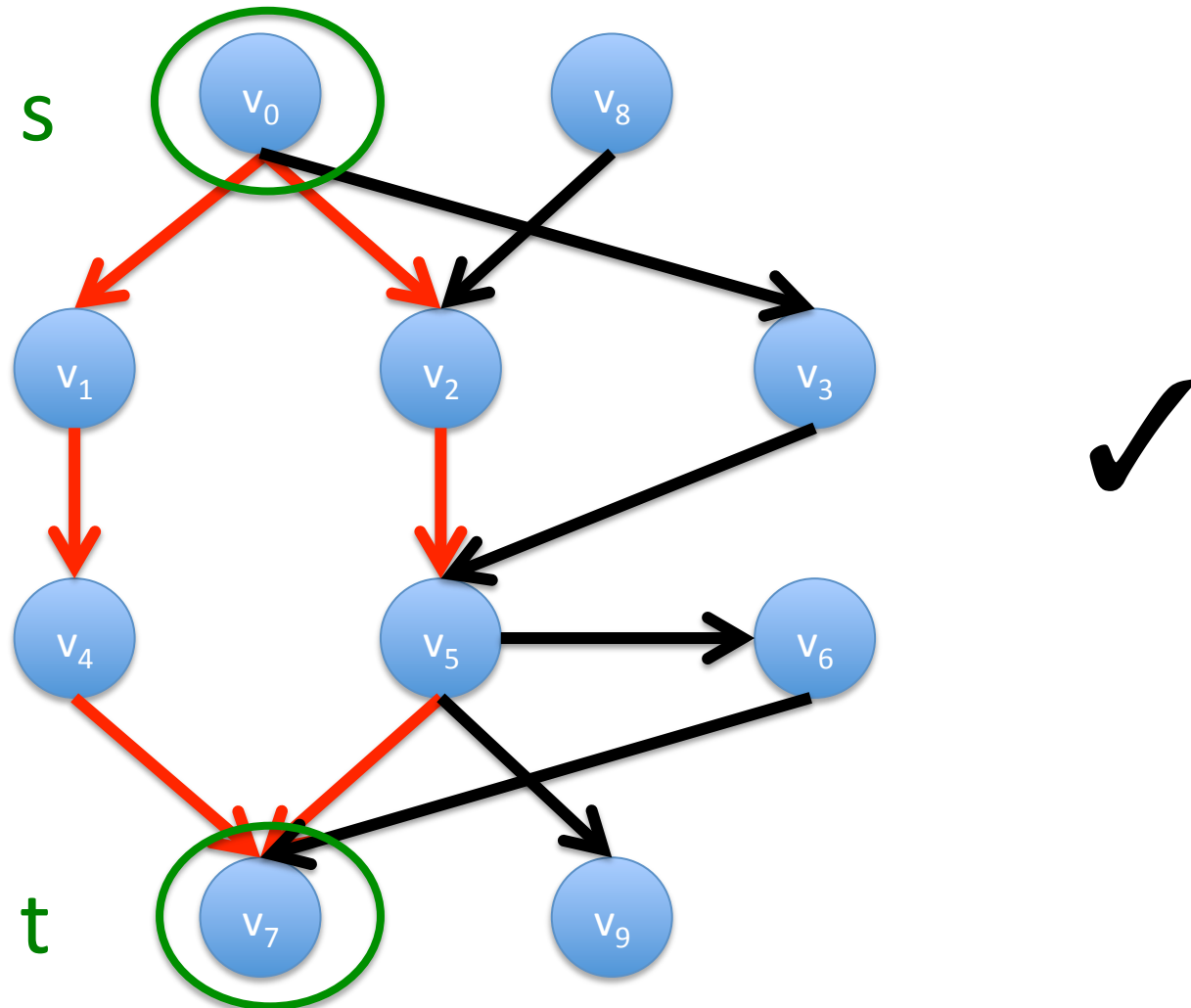
Set of s-t Paths with no common internal vertex

Arc Disjoint s-t Paths



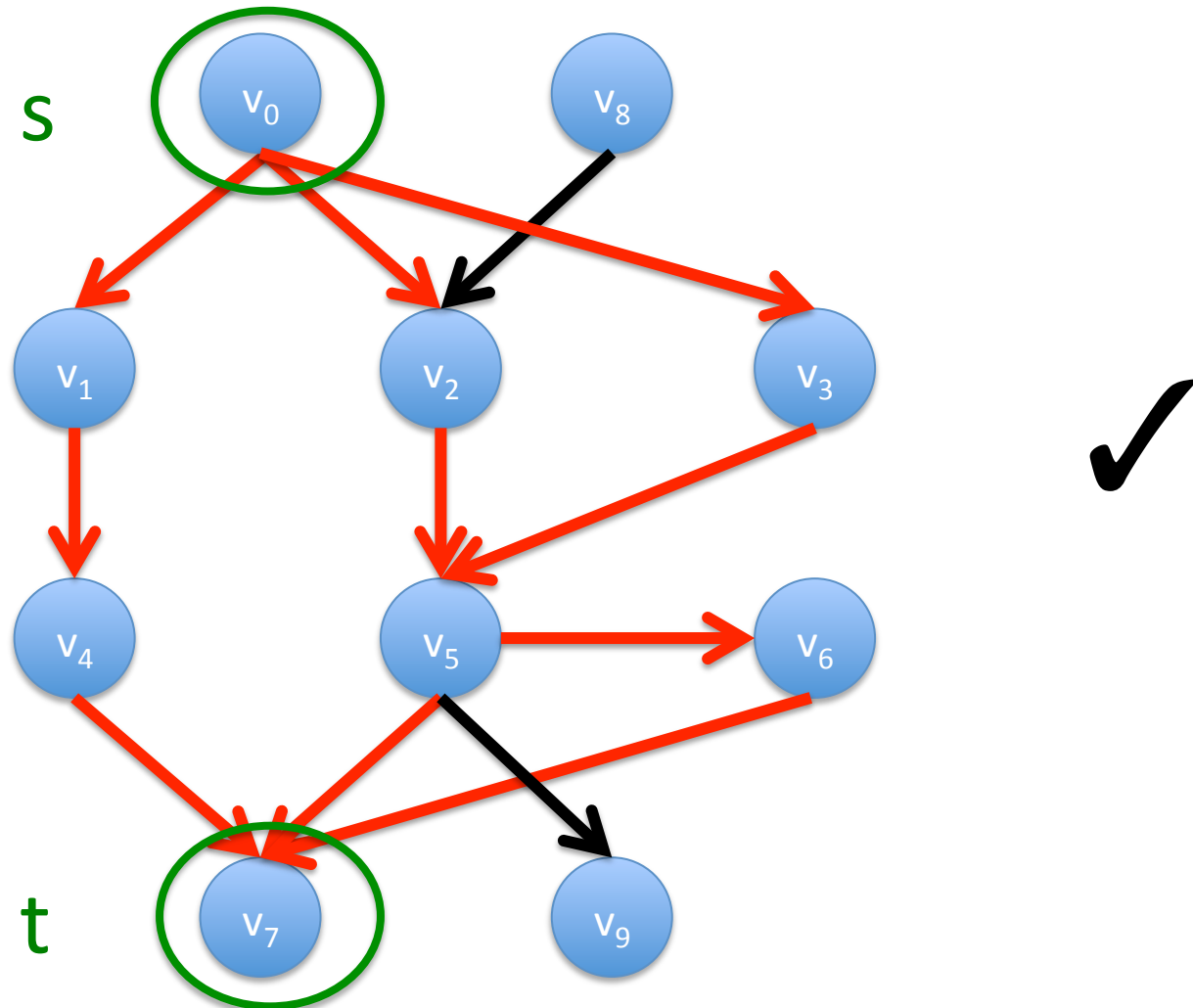
Set of s-t Paths with no common arcs

Arc Disjoint s-t Paths



Set of s-t Paths with no common arcs

Arc Disjoint s-t Paths

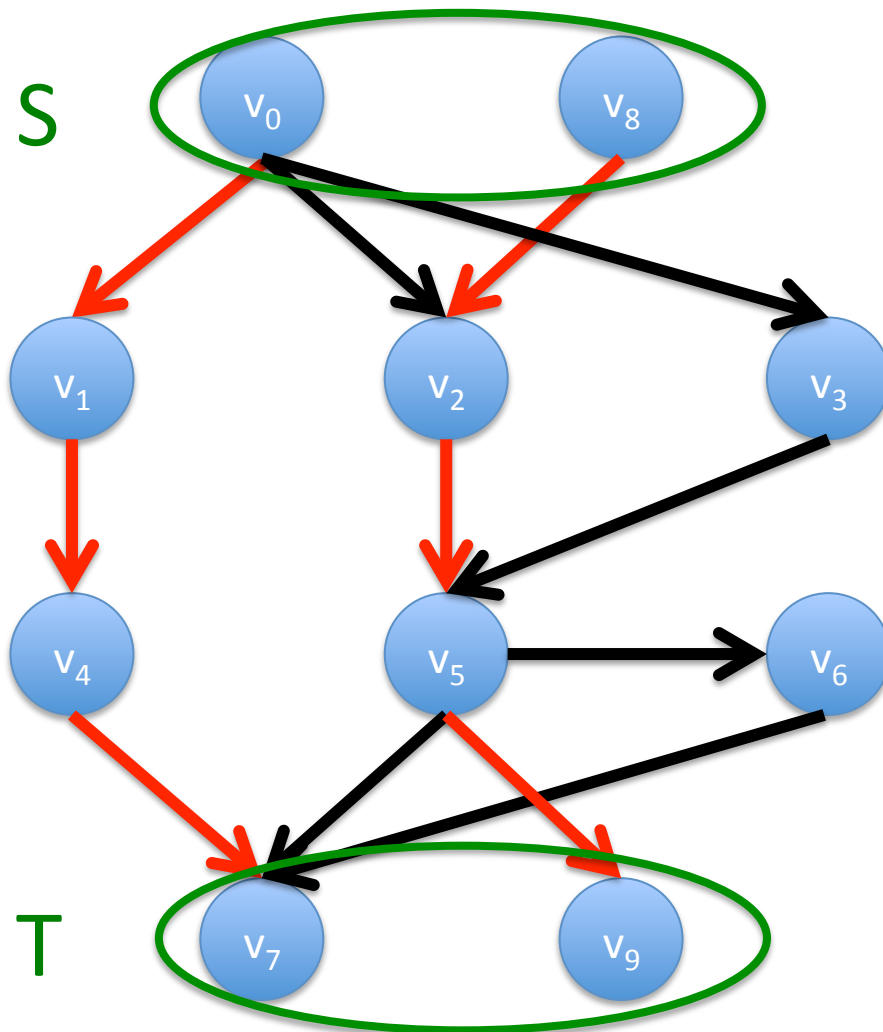


Set of s-t Paths with no common arcs

Outline

- Preliminaries
- Menger's Theorem for Disjoint Paths
 - **Vertex Disjoint S-T Paths**
 - Internally Vertex Disjoint s-t Paths
 - Arc Disjoint s-t Paths
- Path Packing

Vertex Disjoint S-T Paths

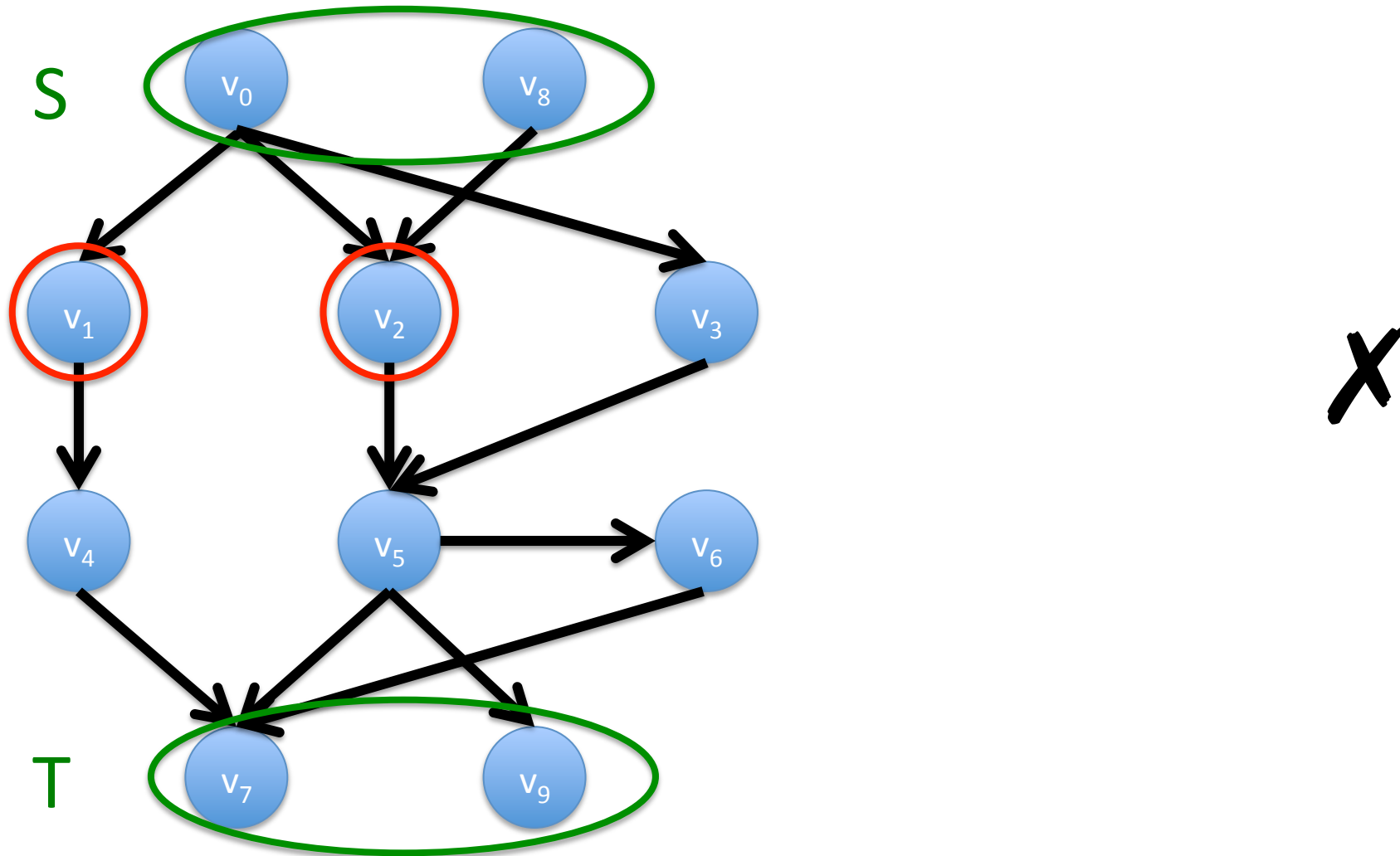


Maximum number of disjoint paths?

Minimum size of S-T disconnecting vertex set !!

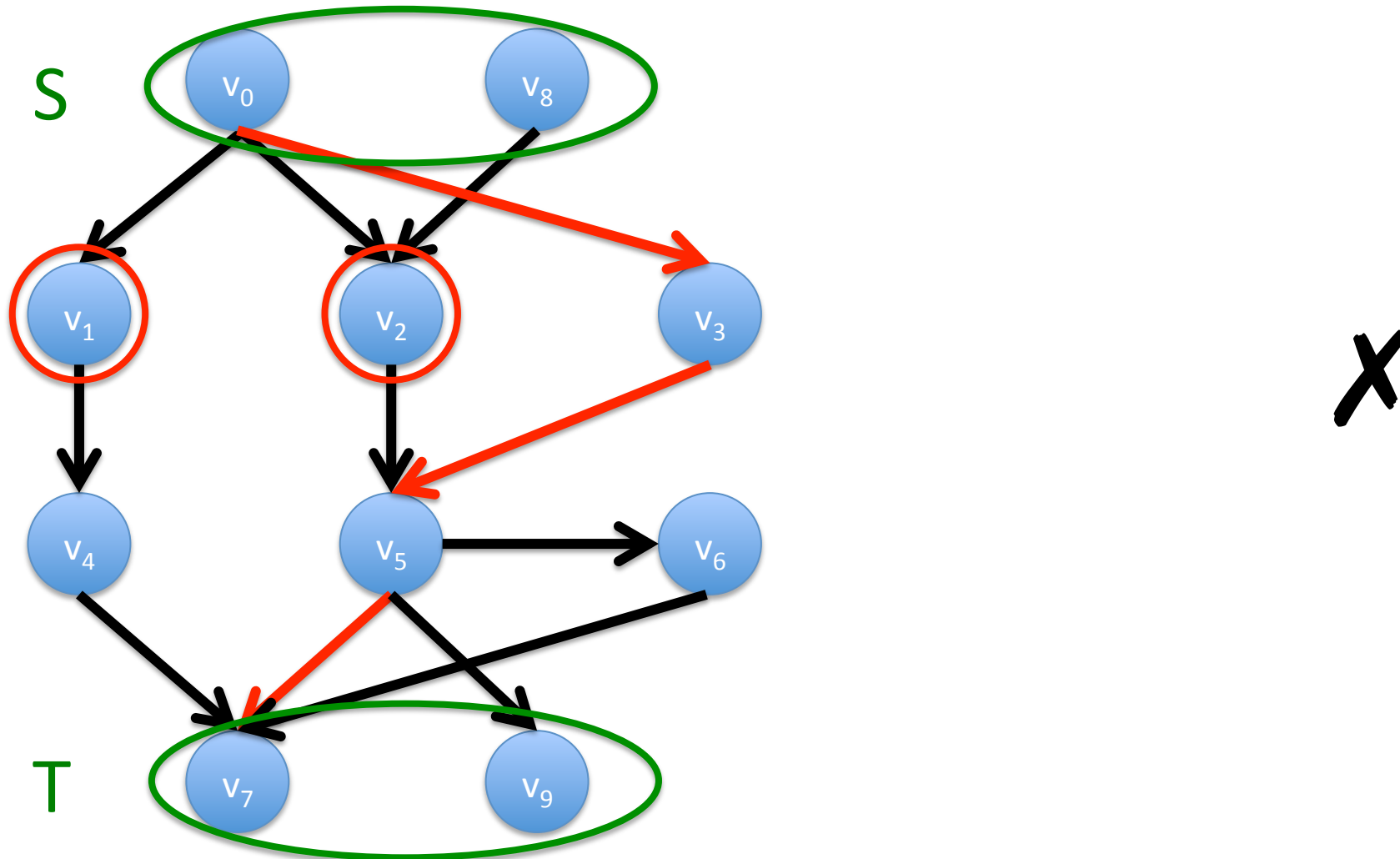
Set of S-T Paths with no common vertex

S-T Disconnecting Vertex Set



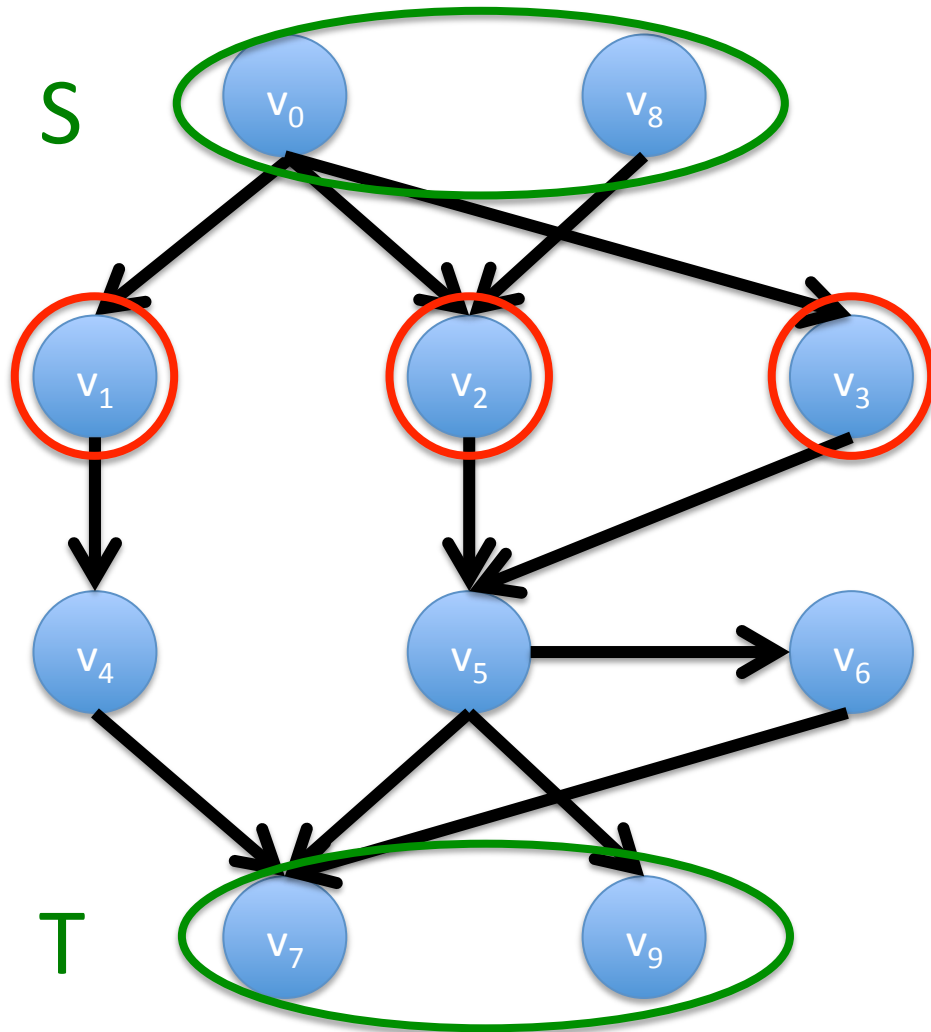
Subset U of V which intersects with all S-T Paths

S-T Disconnecting Vertex Set



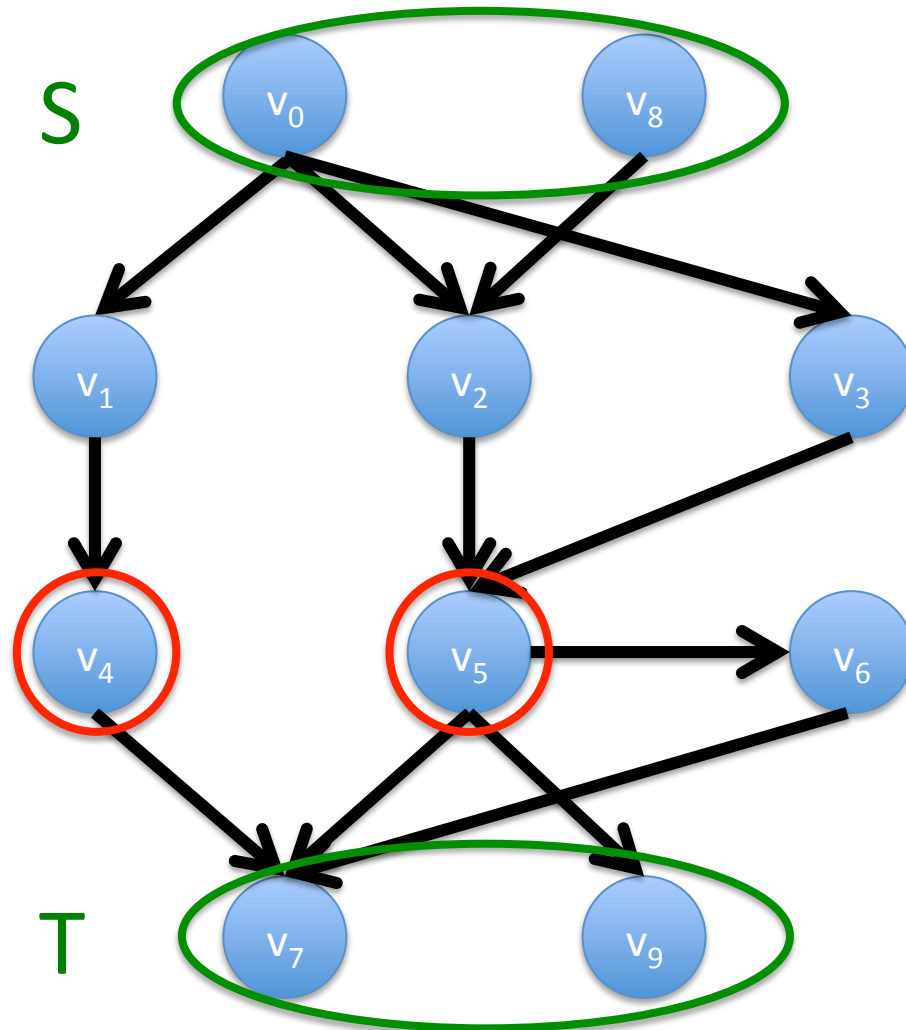
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S-T Disconnecting Vertex Set



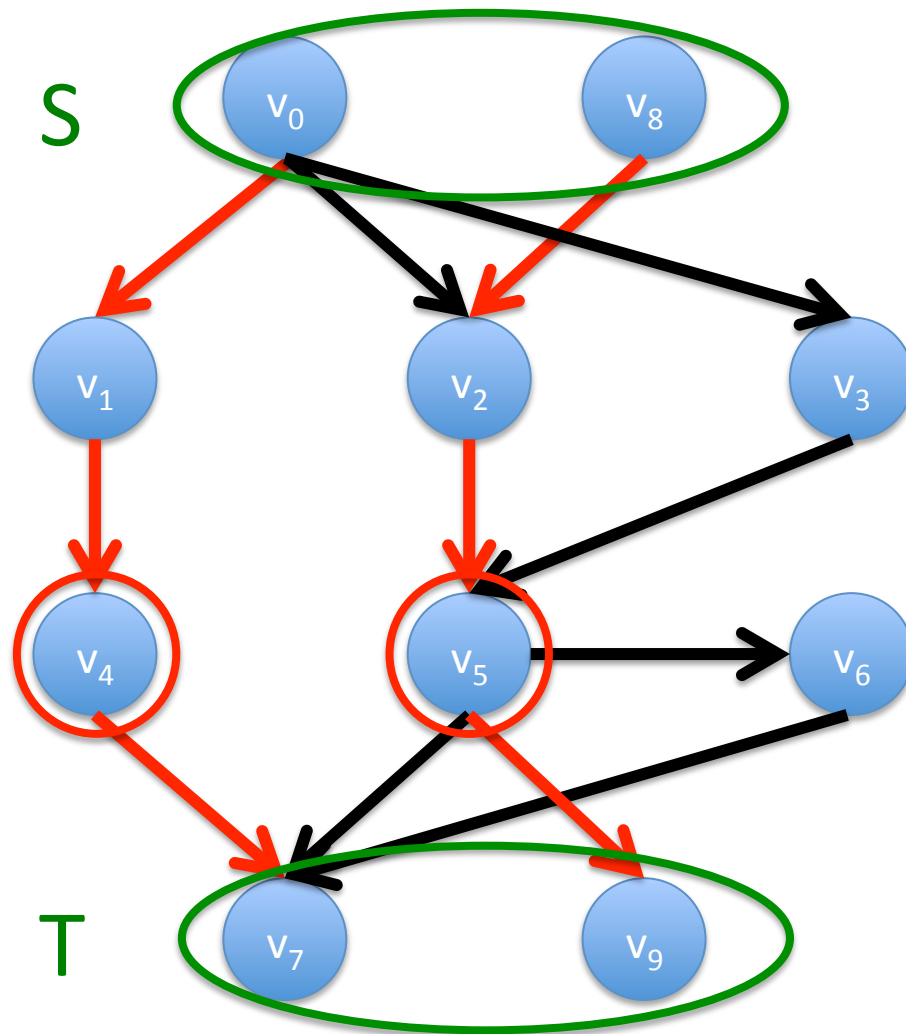
Subset U of V which intersects with all S-T Paths

S-T Disconnecting Vertex Set



Subset U of V which intersects with all S-T Paths

Connection

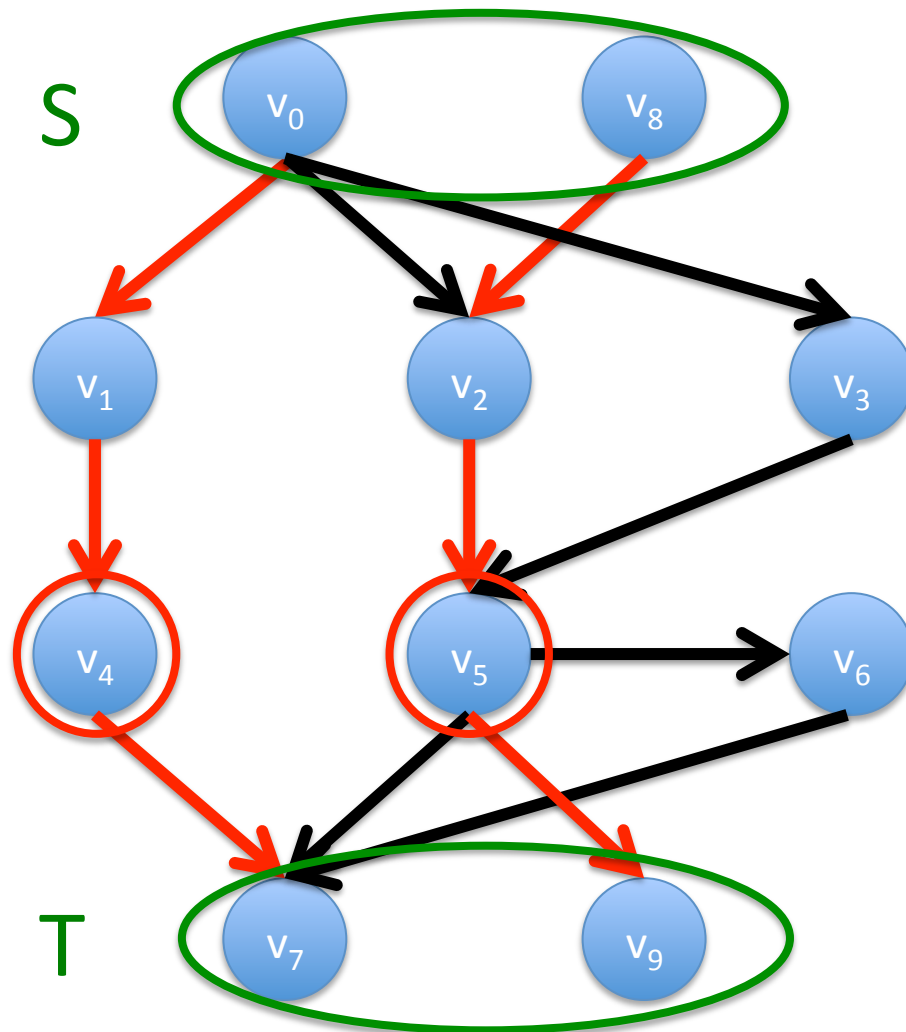


Maximum number of
disjoint paths

\leq

Minimum size of
S-T disconnecting
vertex set !!

Menger's Theorem



Maximum number of
disjoint paths

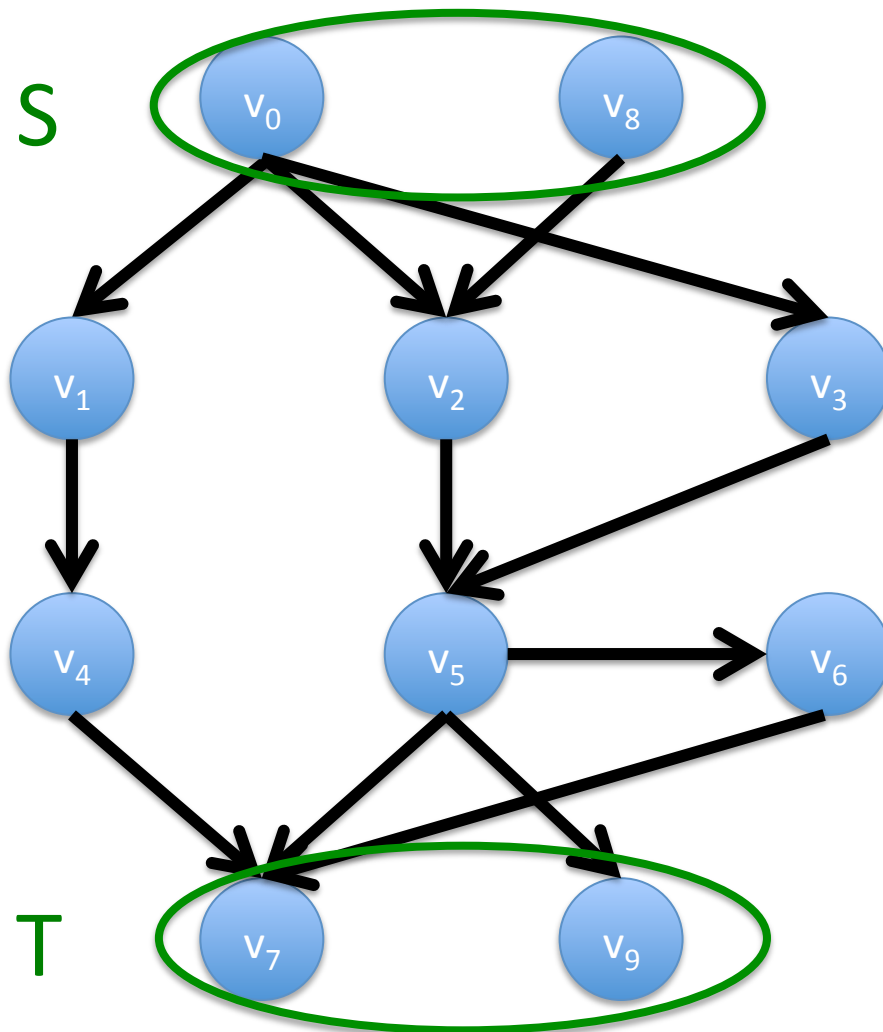
=

Minimum size of
S-T disconnecting
vertex set !!

Proof ?

Mathematical Induction on $|A|$

Menger's Theorem

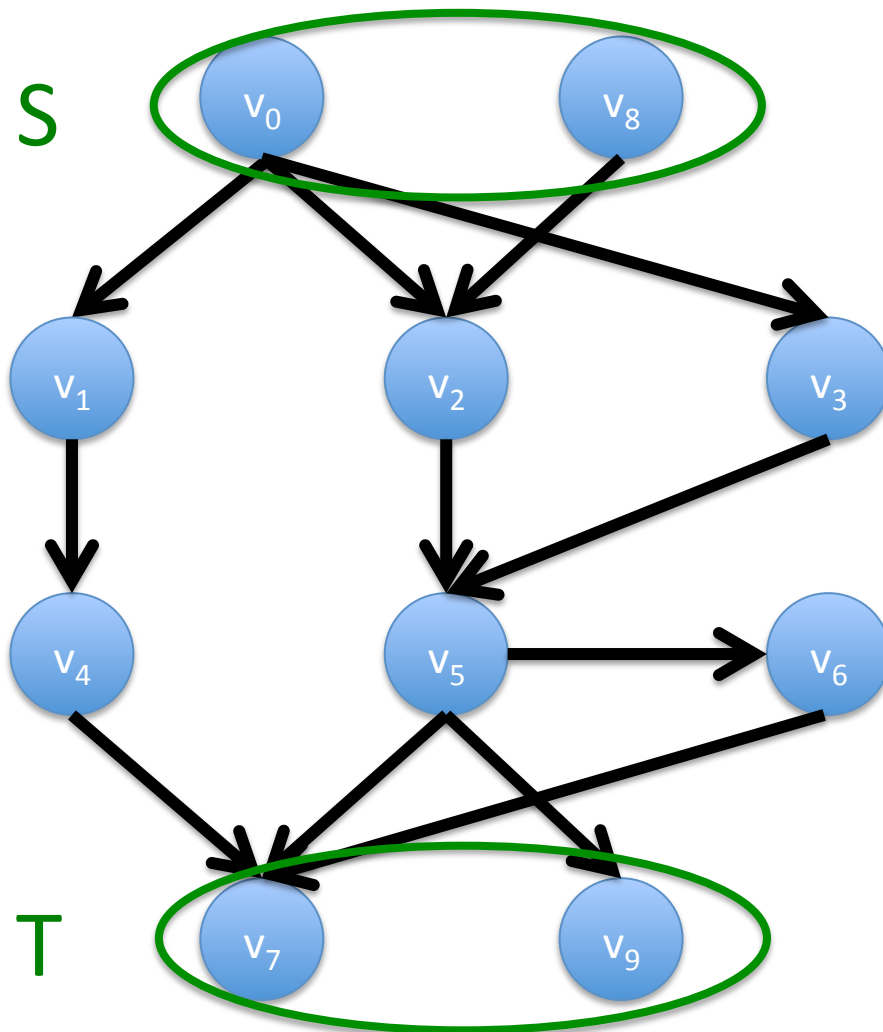


True for $|A| = 0$

Assume it is
true for $|A| < m$

Mathematical Induction on $|A|$

Menger's Theorem



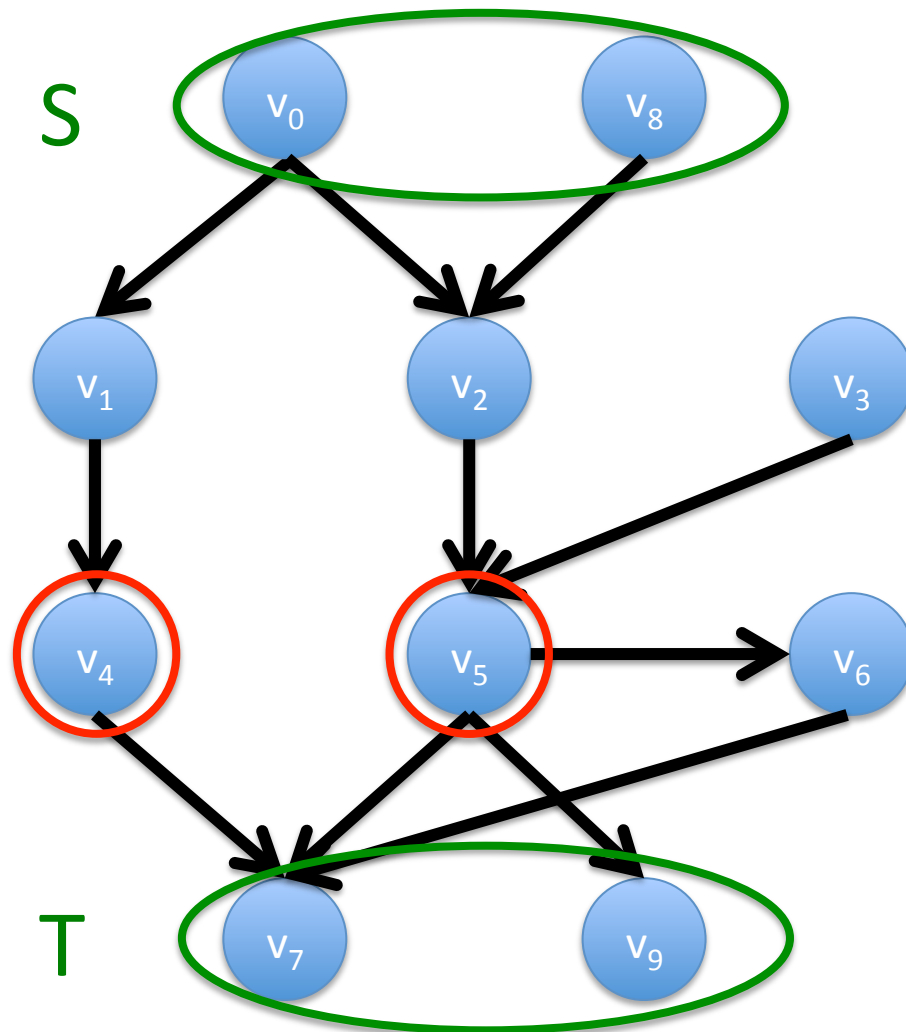
Consider $|A| = m$

Let minimum size
disconnecting vertex
set = k

Remove an arc (u,v)

Mathematical Induction on $|A|$

Menger's Theorem

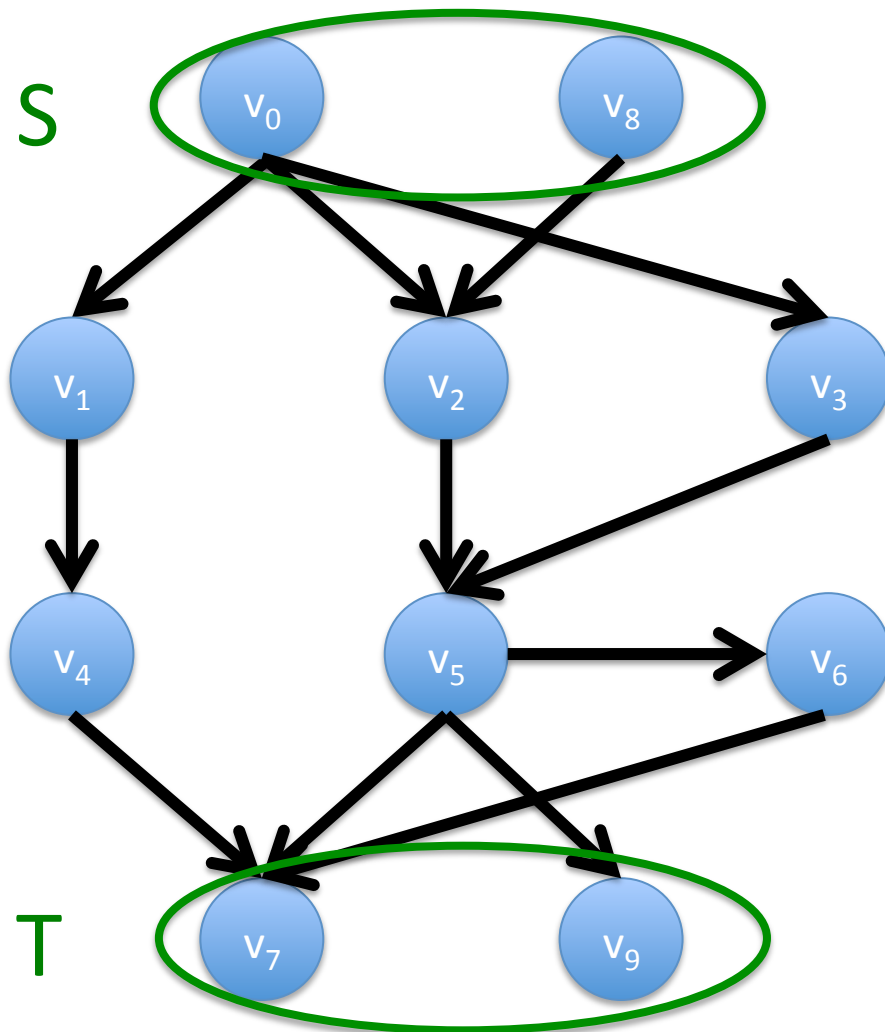


Minimum disconnecting
vertex set = U . $|U| = k$

Proved by induction

Mathematical Induction on $|A|$

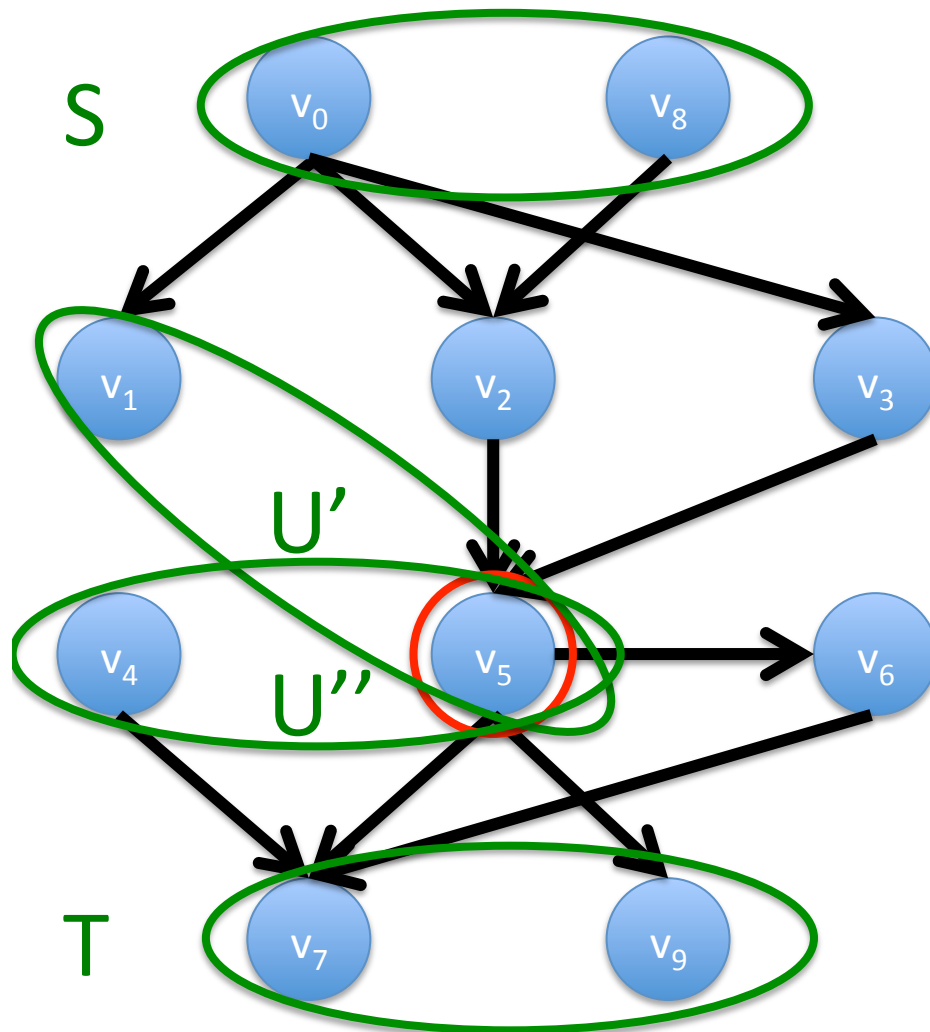
Menger's Theorem



Remove an arc (u,v)

Mathematical Induction on $|A|$

Menger's Theorem



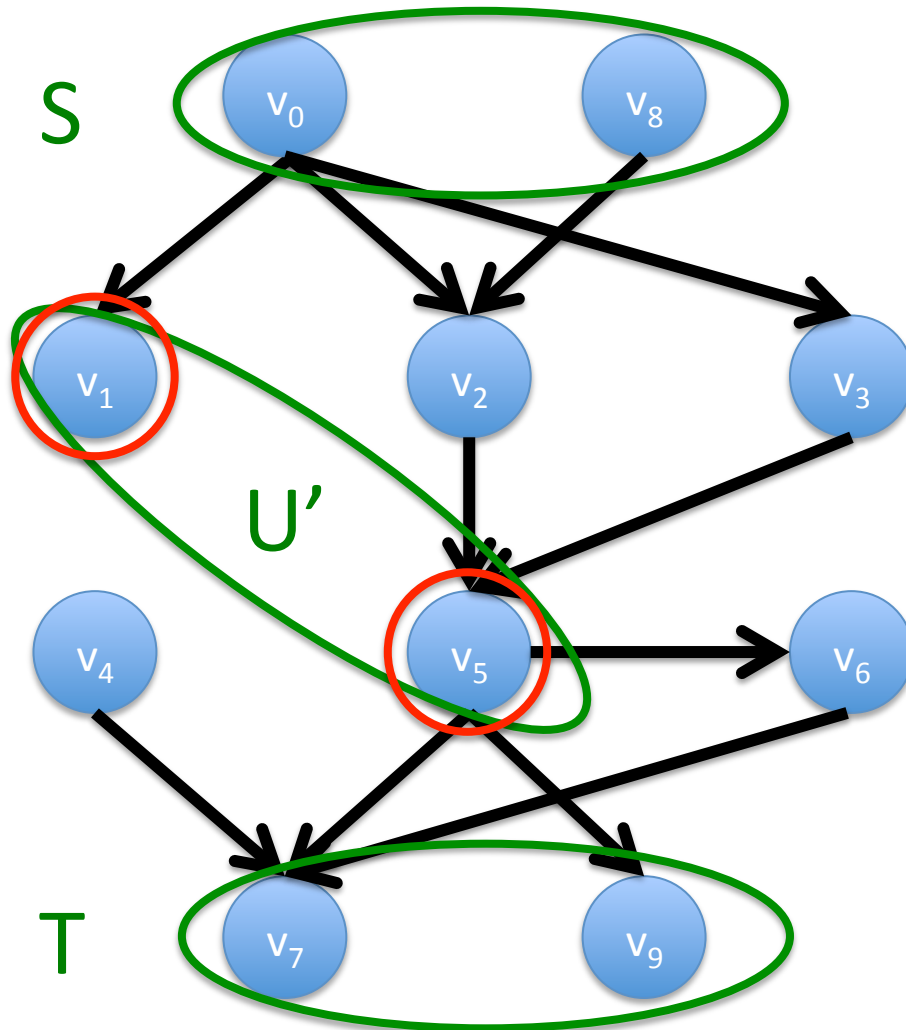
Minimum disconnecting
vertex set U . $|U| = k-1$

$U' = U \cup \{u\}$

$U'' = U \cup \{v\}$

Mathematical Induction on $|A|$

Menger's Theorem



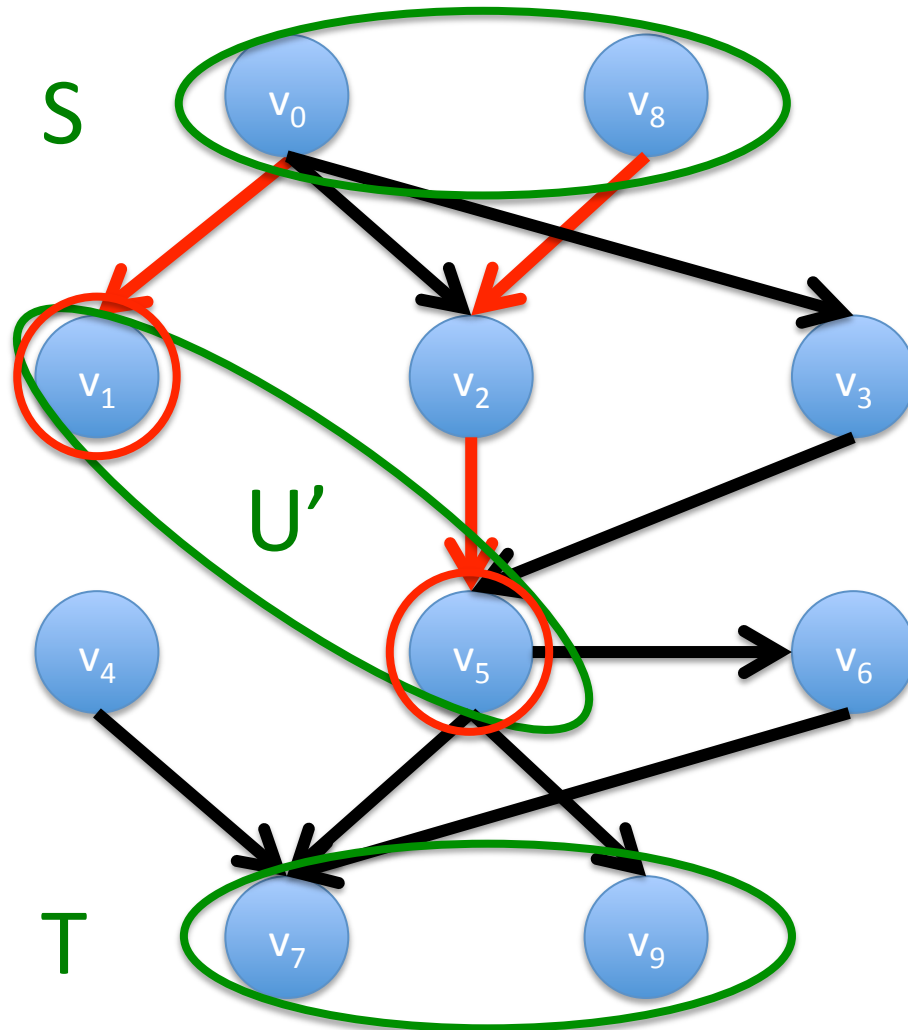
U' is an S-T disconnecting vertex set of size k in D

Any S- U' disconnecting vertex set has size $\geq k$

Graph contains k disjoint S- U' paths (induction)

Mathematical Induction on $|A|$

Menger's Theorem



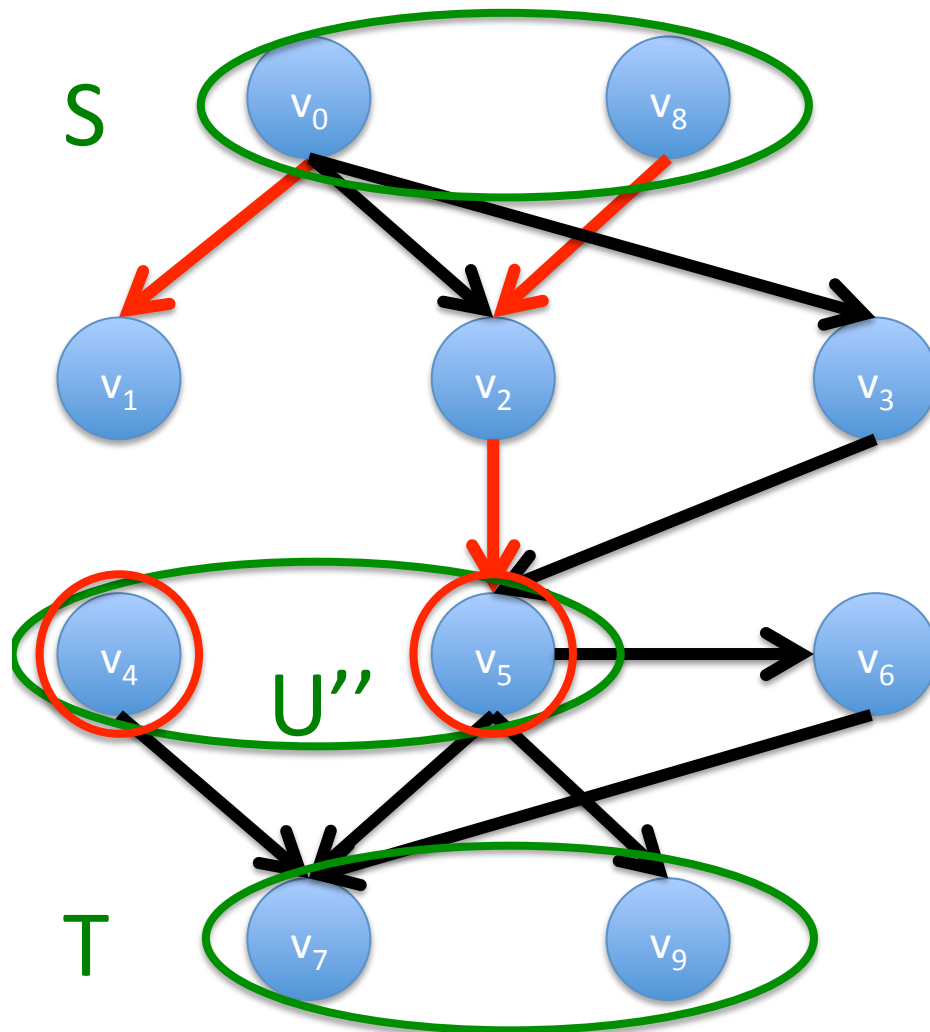
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Mathematical Induction on $|A|$

Menger's Theorem



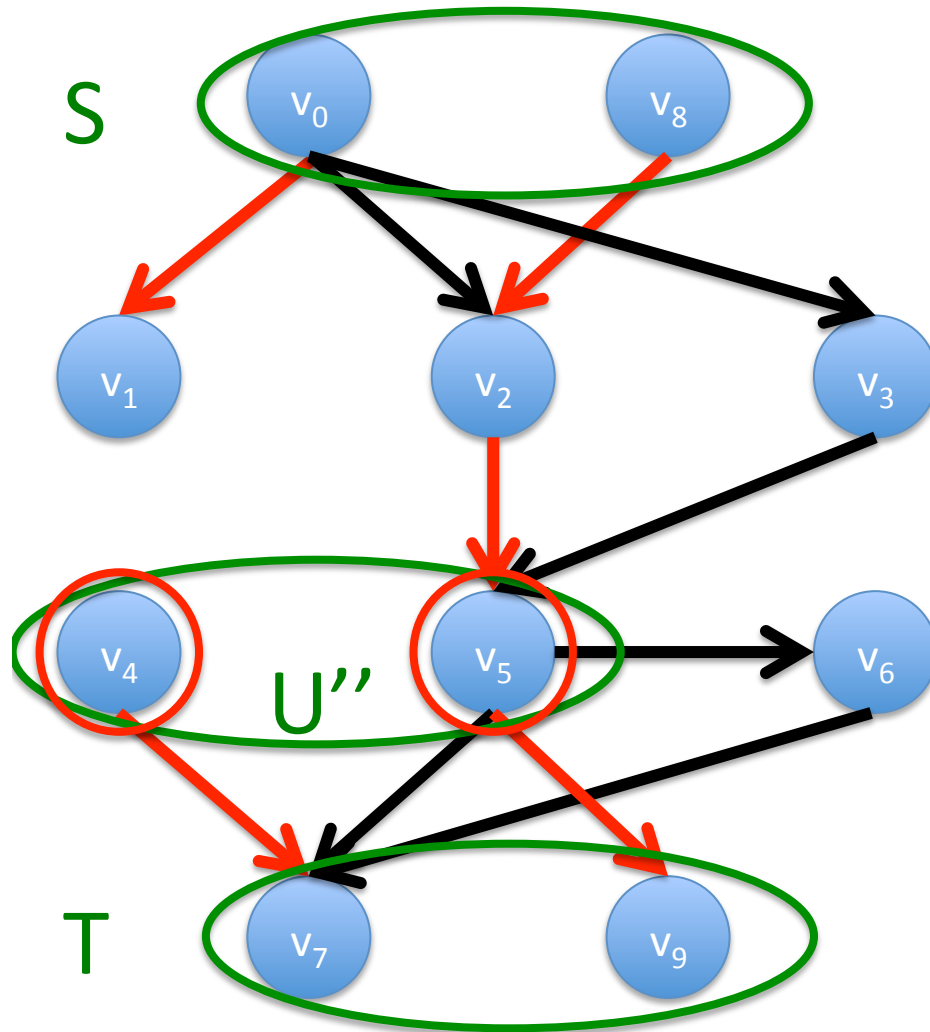
U'' is an S-T disconnecting vertex set of size k in D

Any U'' -T disconnecting vertex set has size $\geq k$

Graph contains k disjoint U'' -T paths (induction)

Mathematical Induction on $|A|$

Menger's Theorem



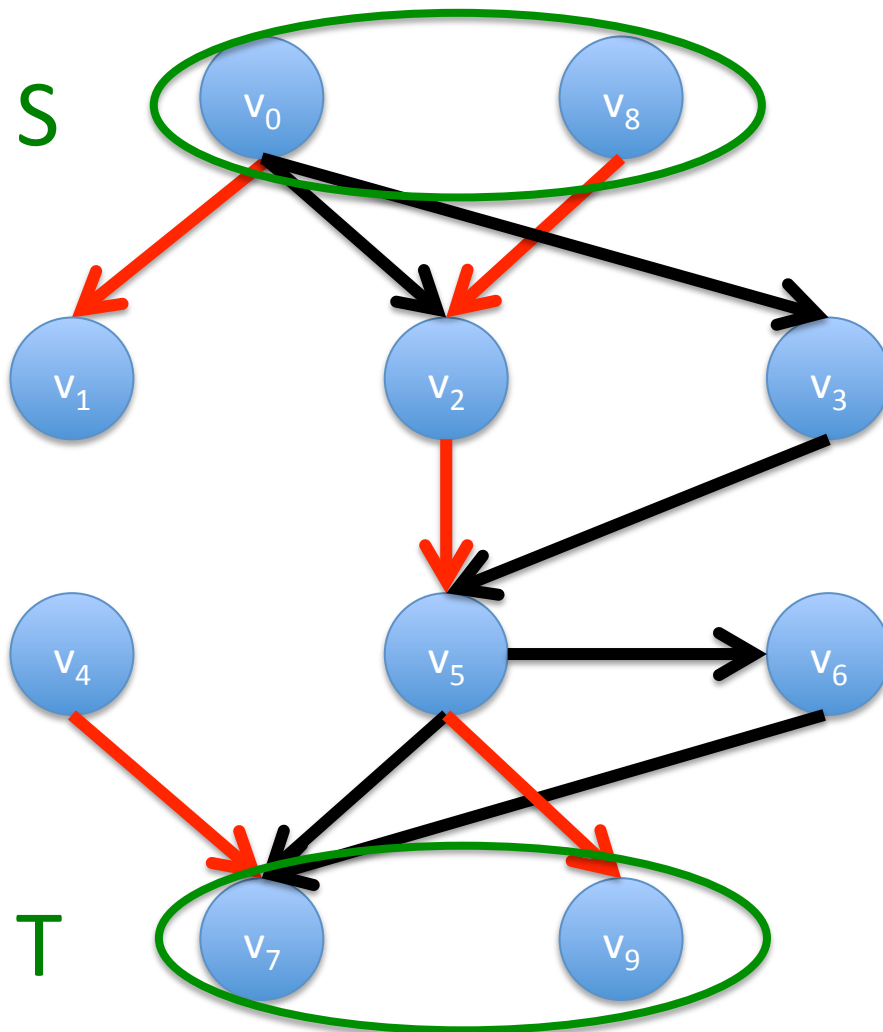
U'' is an S-T disconnecting vertex set of size k in D

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Graph contains k disjoint U'' -T paths (induction)

Mathematical Induction on $|A|$

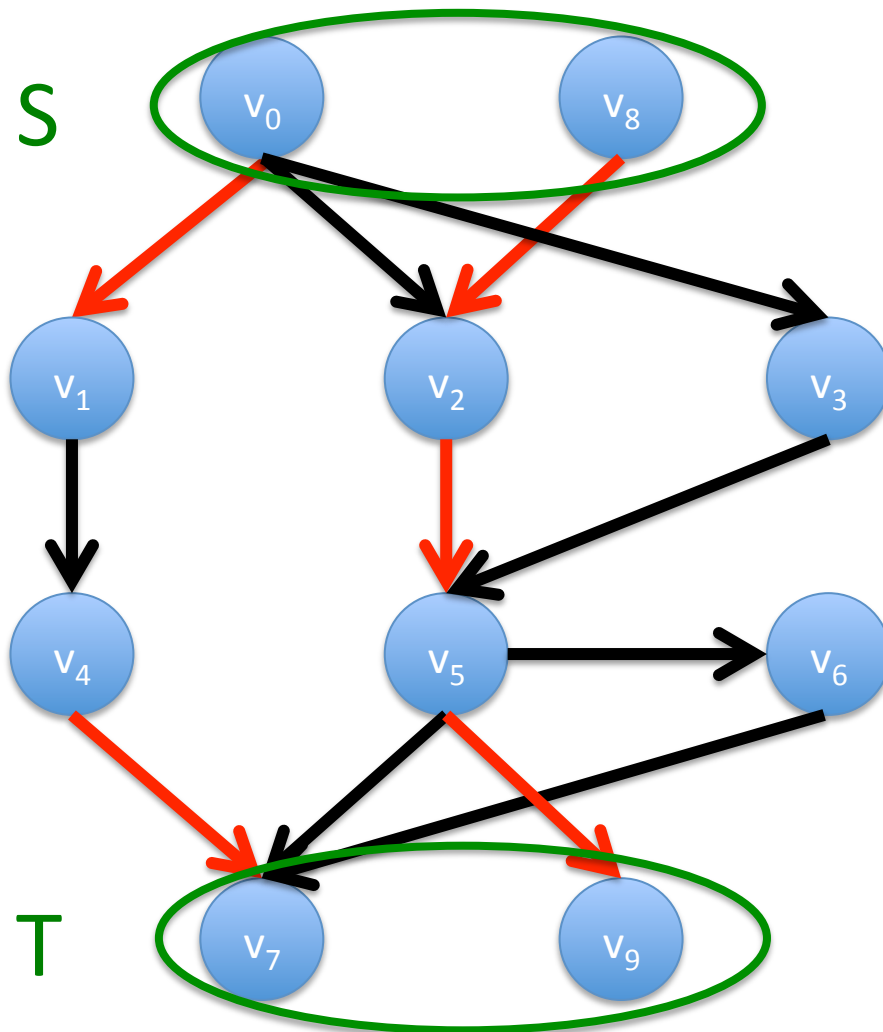
Menger's Theorem



Add arc (u,v) back

Mathematical Induction on $|A|$

Menger's Theorem



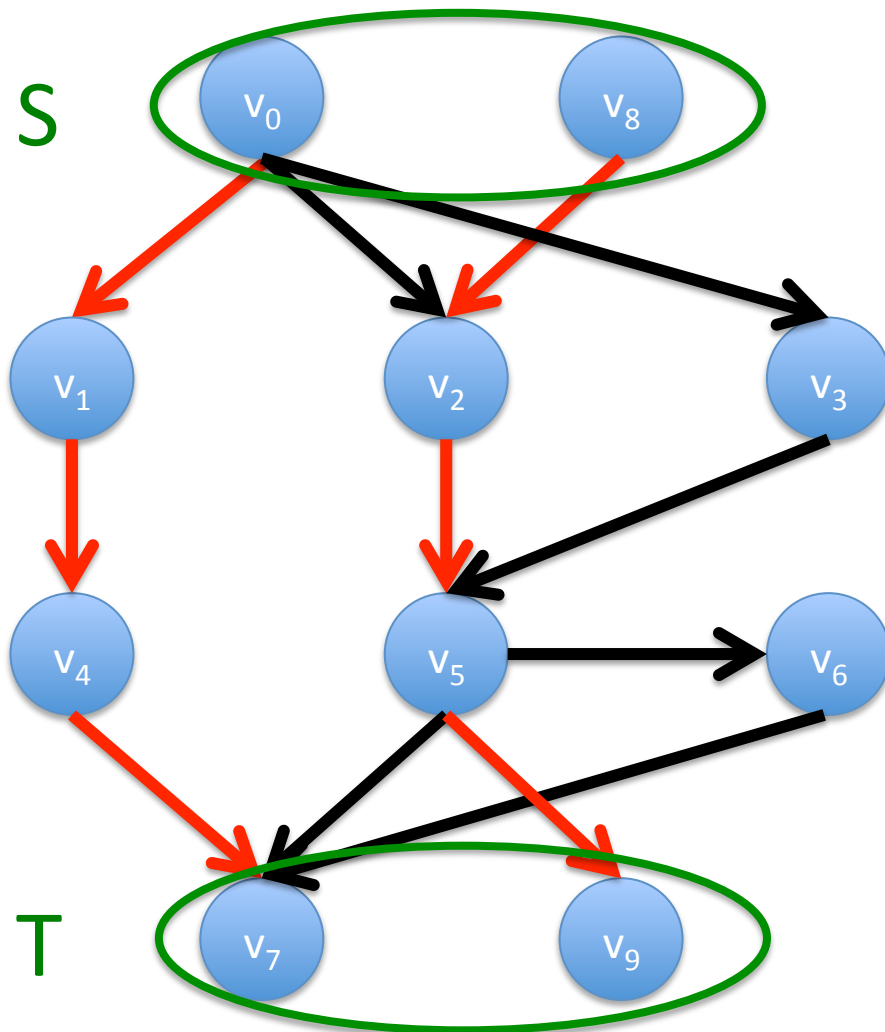
Add arc (u,v) back

$(k-1)$ disjoint pairs of paths intersecting in U

1 disjoint pair of paths now connected by (u,v)

Mathematical Induction on $|A|$

Menger's Theorem



Total k disjoint paths

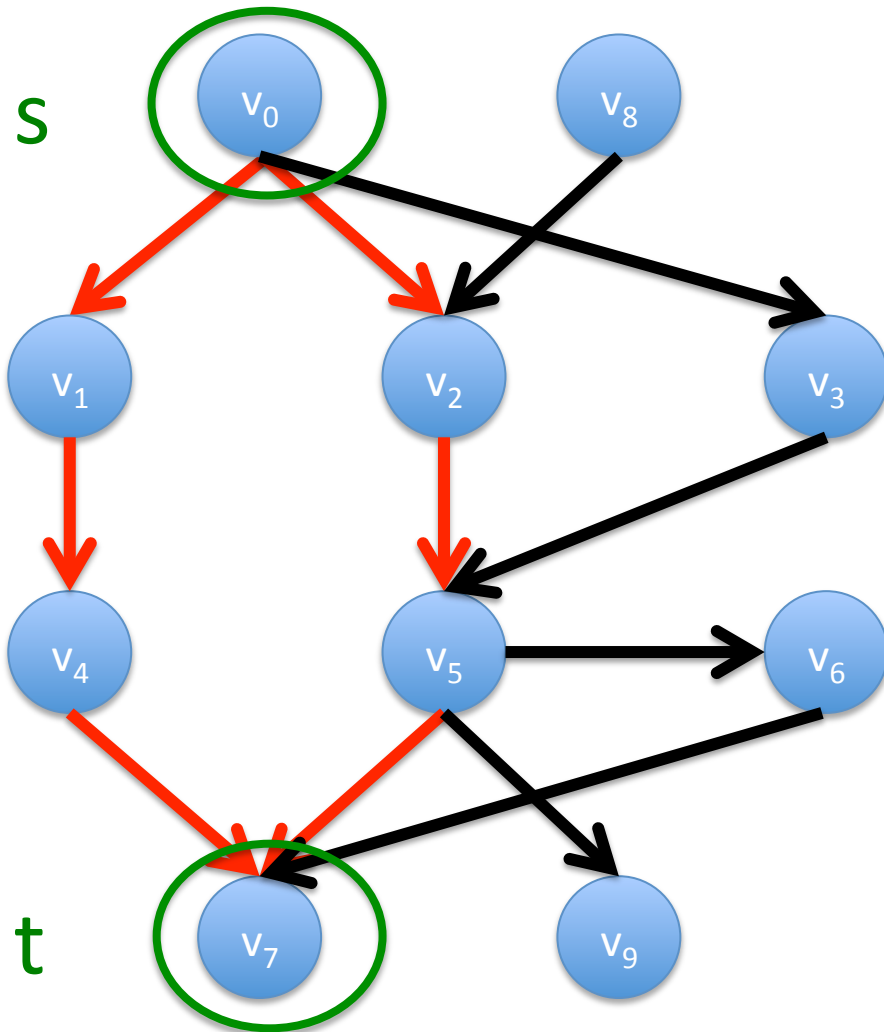
Hence proved

Mathematical Induction on $|A|$

Outline

- Preliminaries
- Menger's Theorem for Disjoint Paths
 - Vertex Disjoint S-T Paths
 - **Internally Vertex Disjoint s-t Paths**
 - Arc Disjoint s-t Paths
- Path Packing

Internally Vertex Disjoint s-t Paths

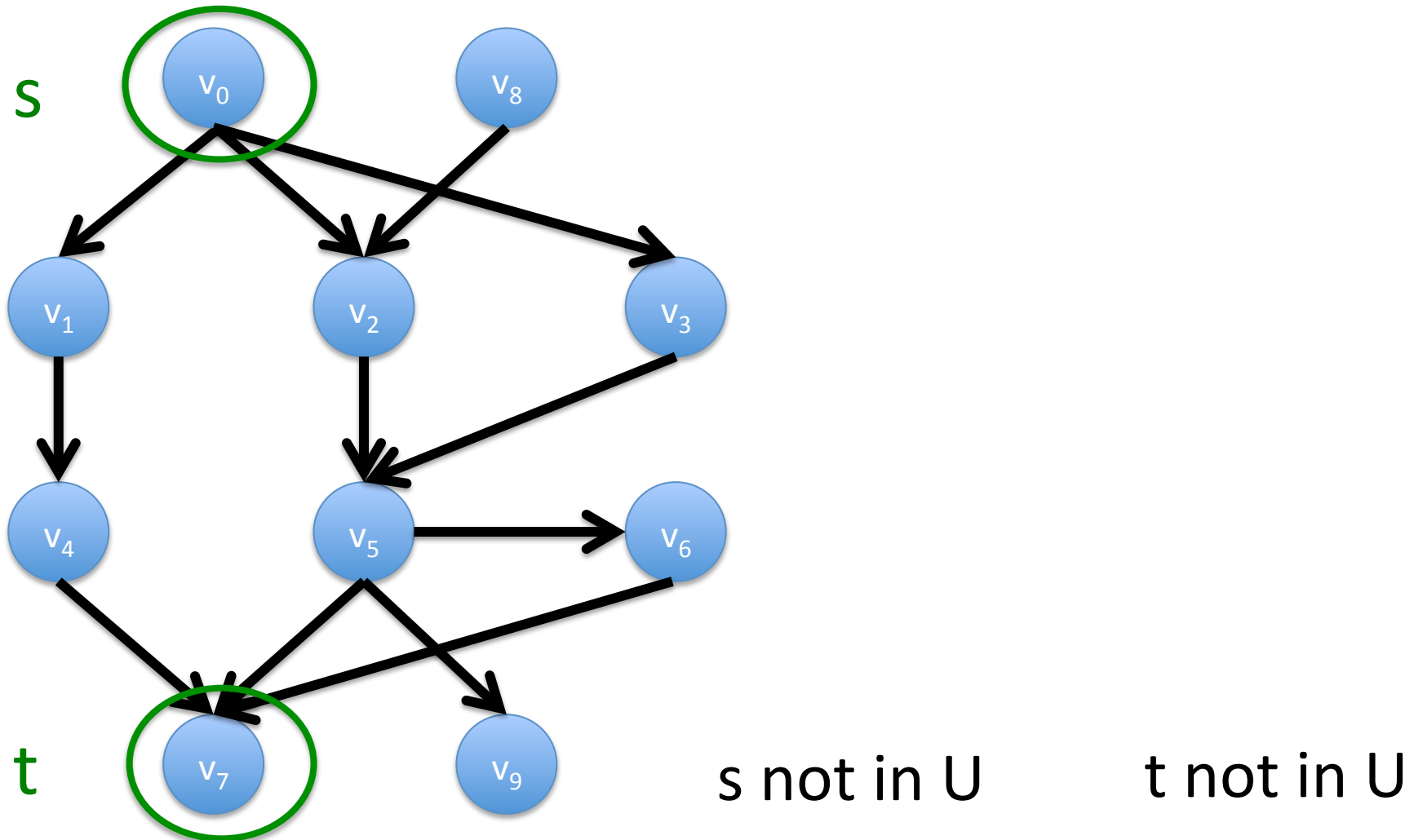


Maximum number of disjoint paths?

Minimum size of s-t vertex-cut !!

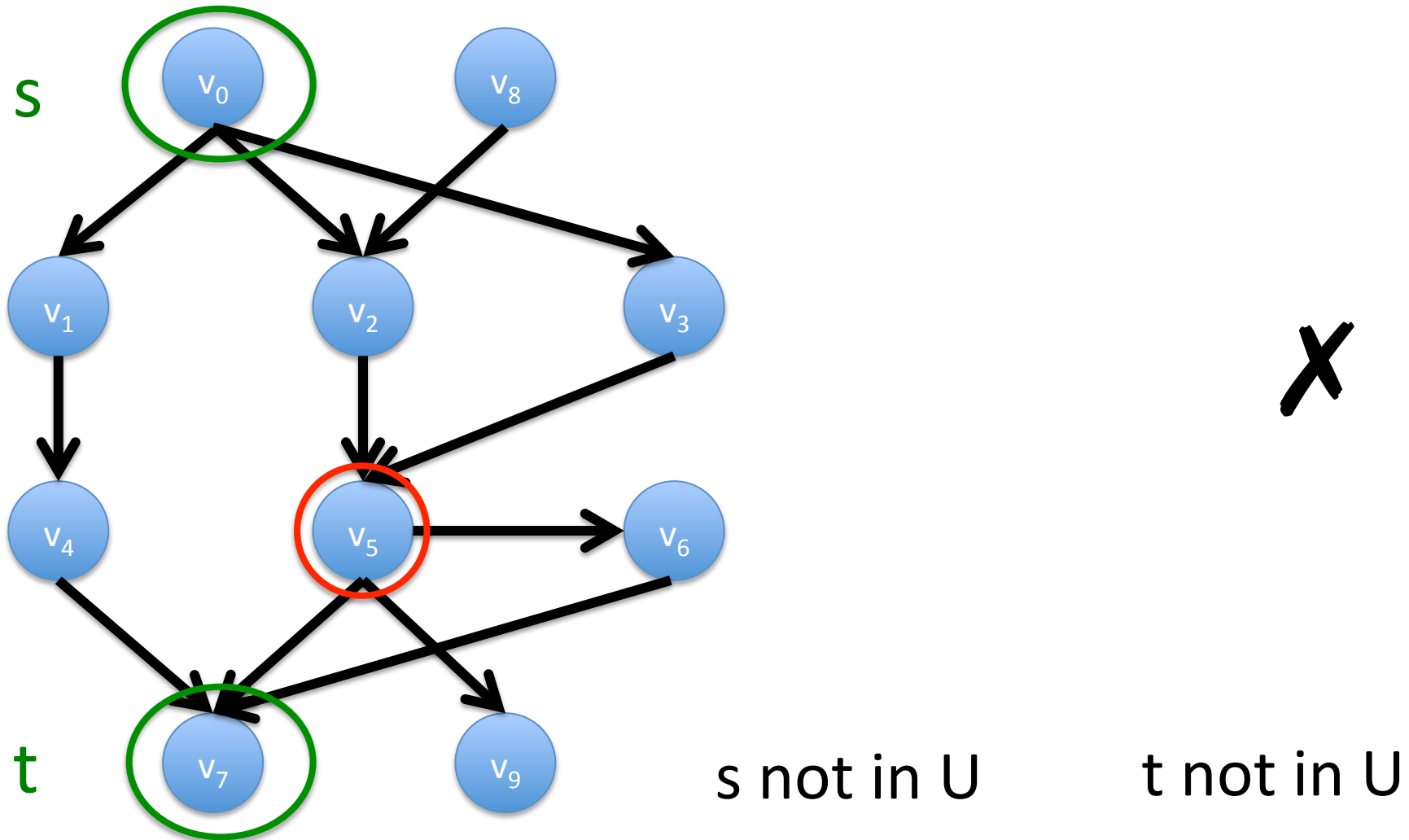
Set of s-t Paths with no common internal vertex

s-t Vertex-Cut



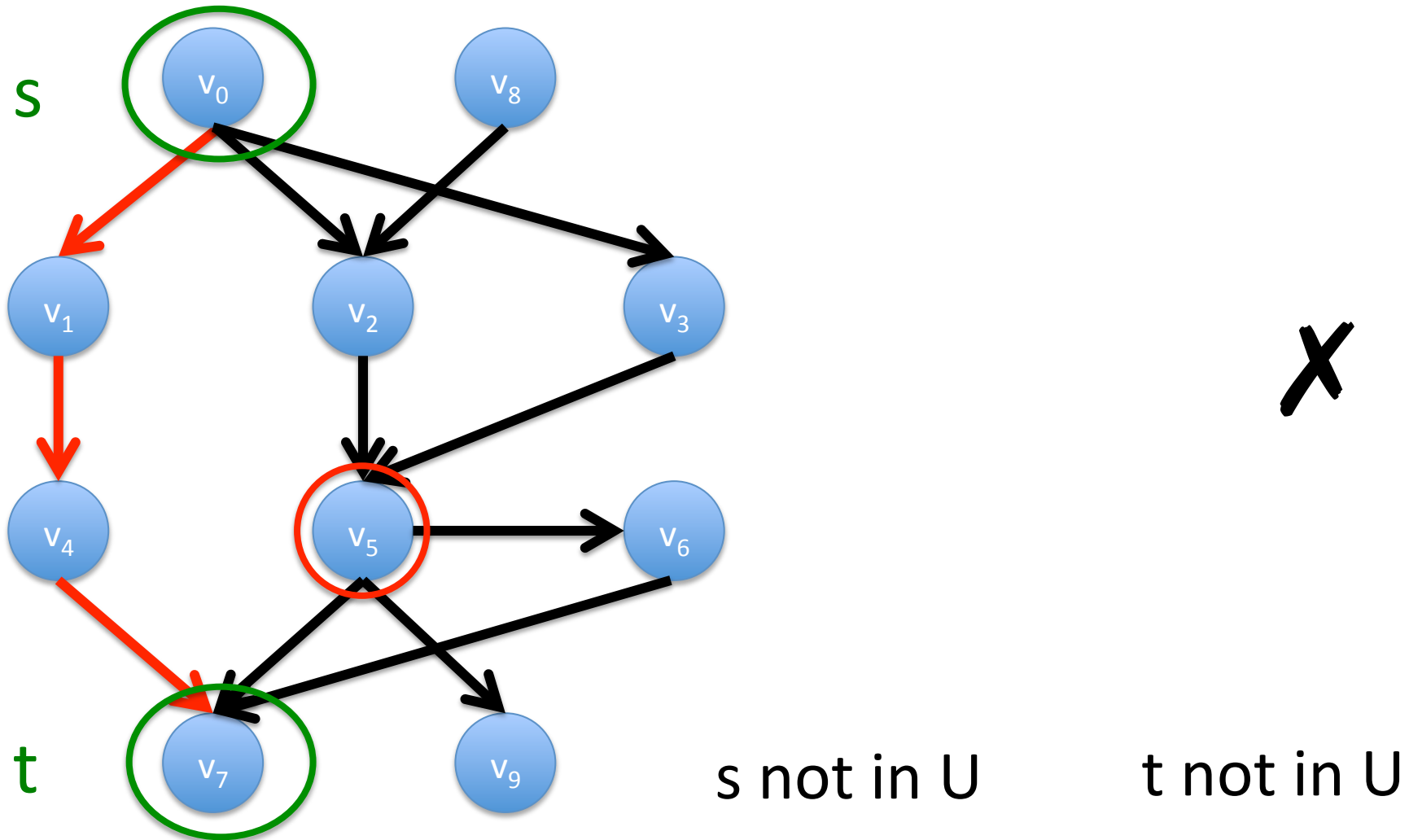
Subset U of V which intersects with all s-t Paths

s-t Vertex-Cut



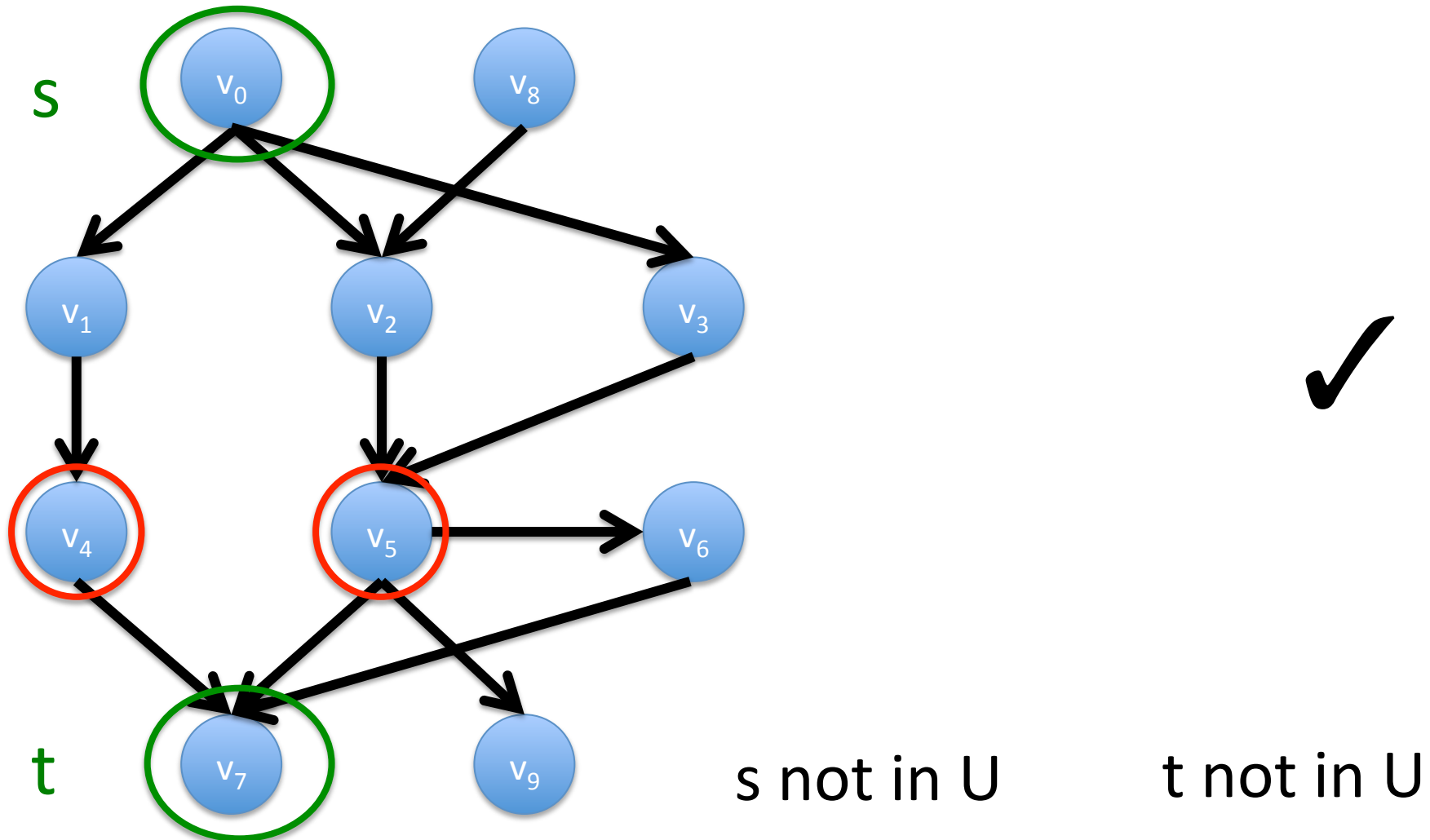
Subset U of V which intersects with all s-t Paths

s-t Vertex-Cut



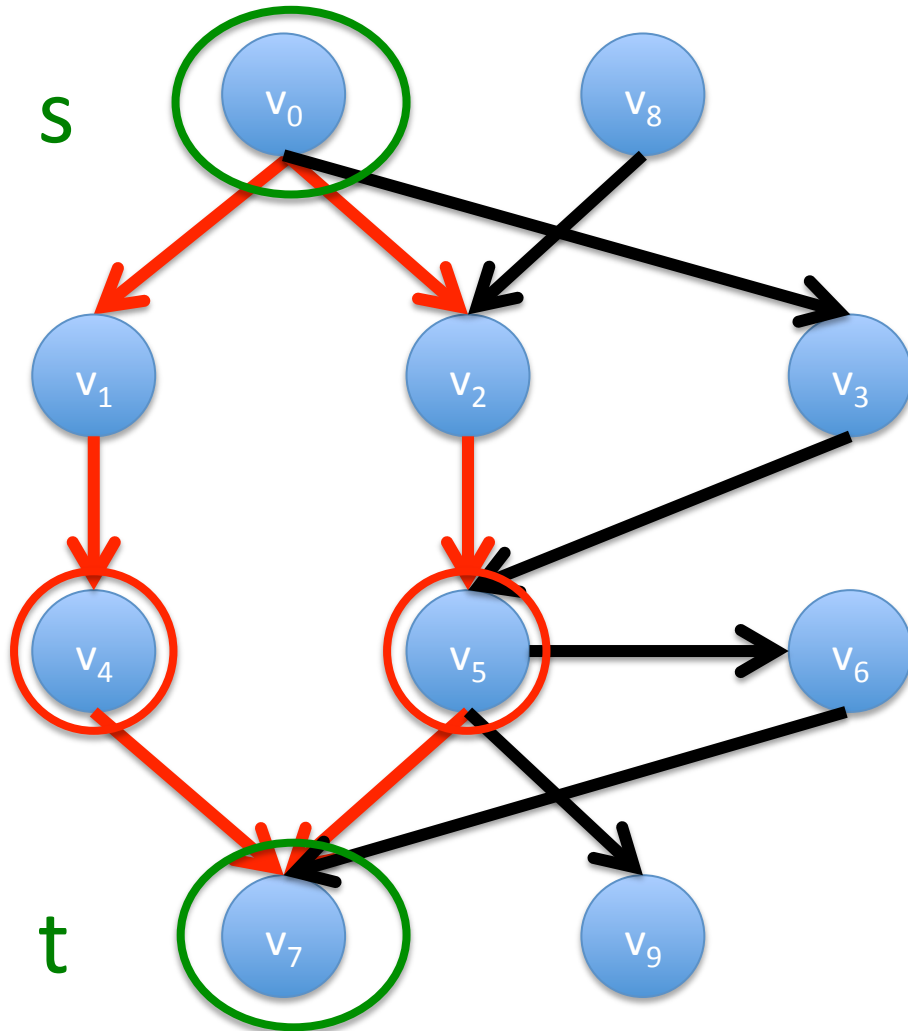
Subset U of V which intersects with all s-t Paths

s-t Vertex-Cut



Subset U of V which intersects with all s - t Paths

Connection

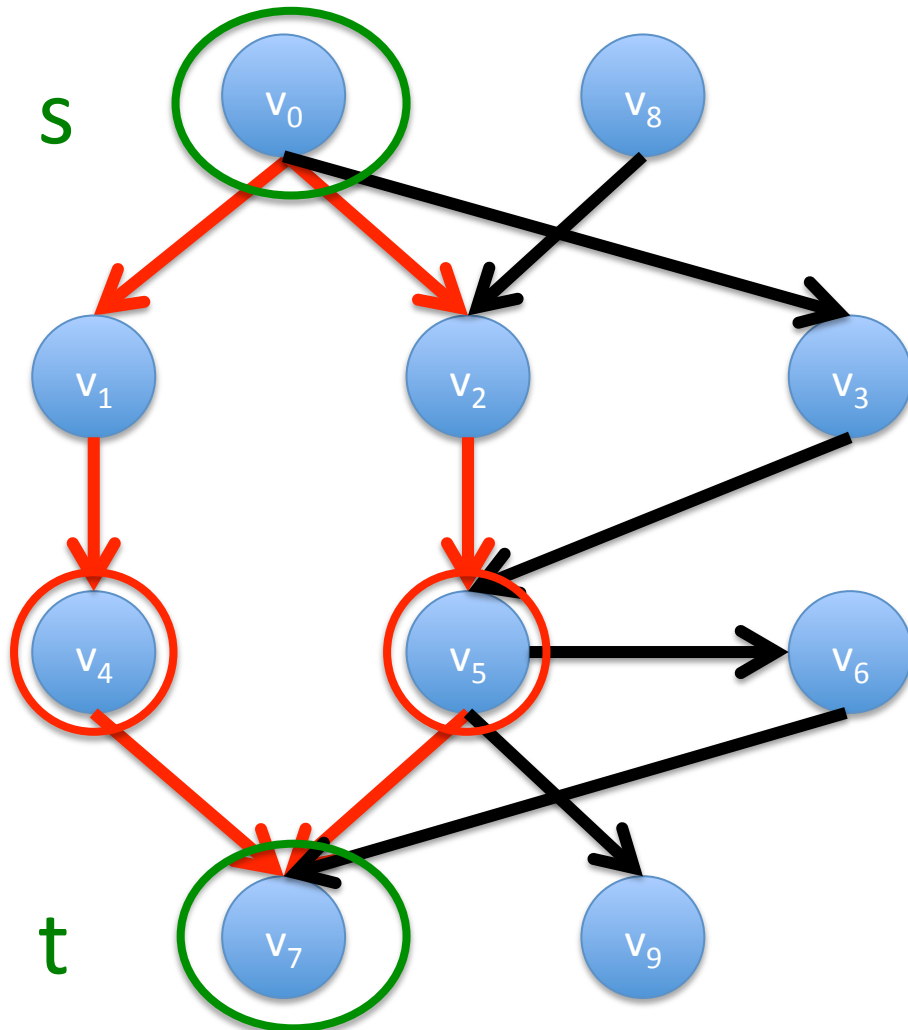


Maximum number of disjoint paths?

\leq

Minimum size of s-t vertex-cut !!

Menger's Theorem



Maximum number of
disjoint paths?

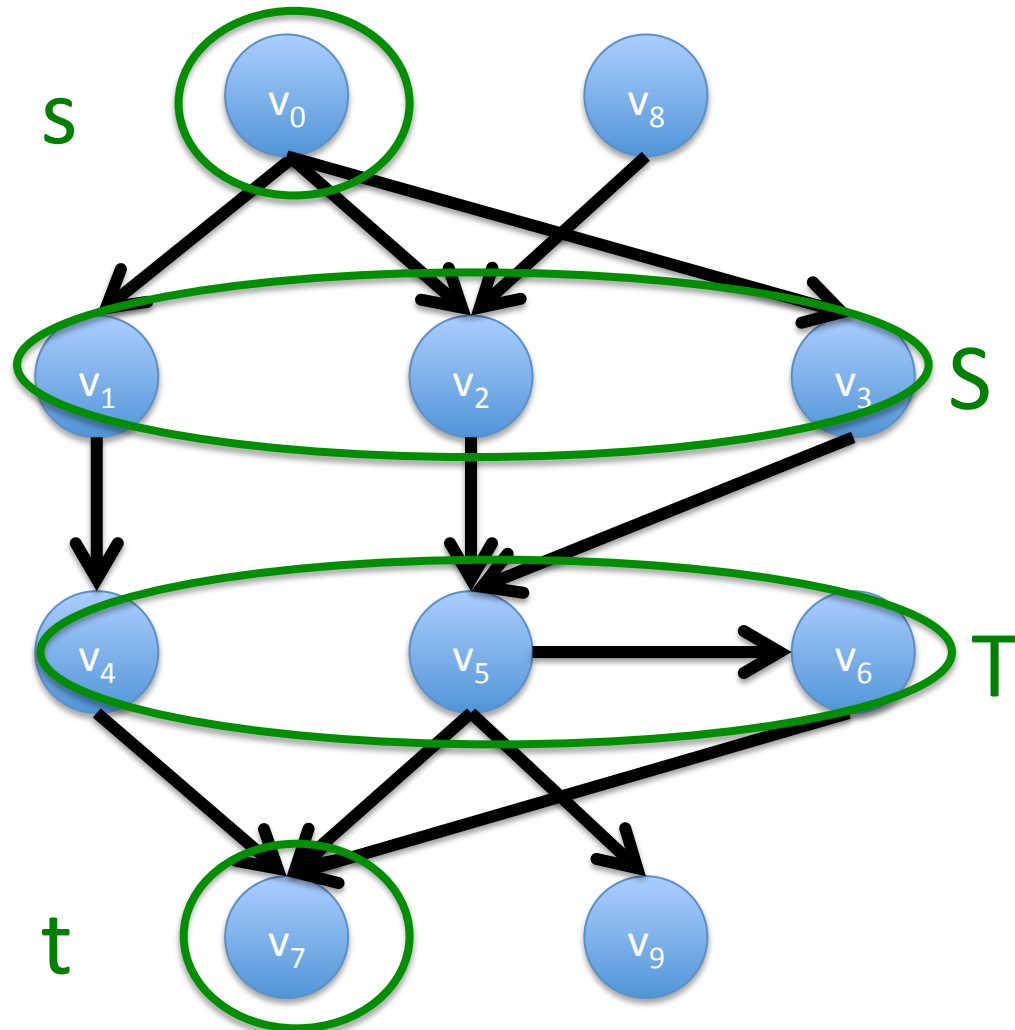
=

Minimum size of
s-t vertex-cut !!

Proof ?

Menger's Theorem for vertex disjoint S-T paths

Menger's Theorem



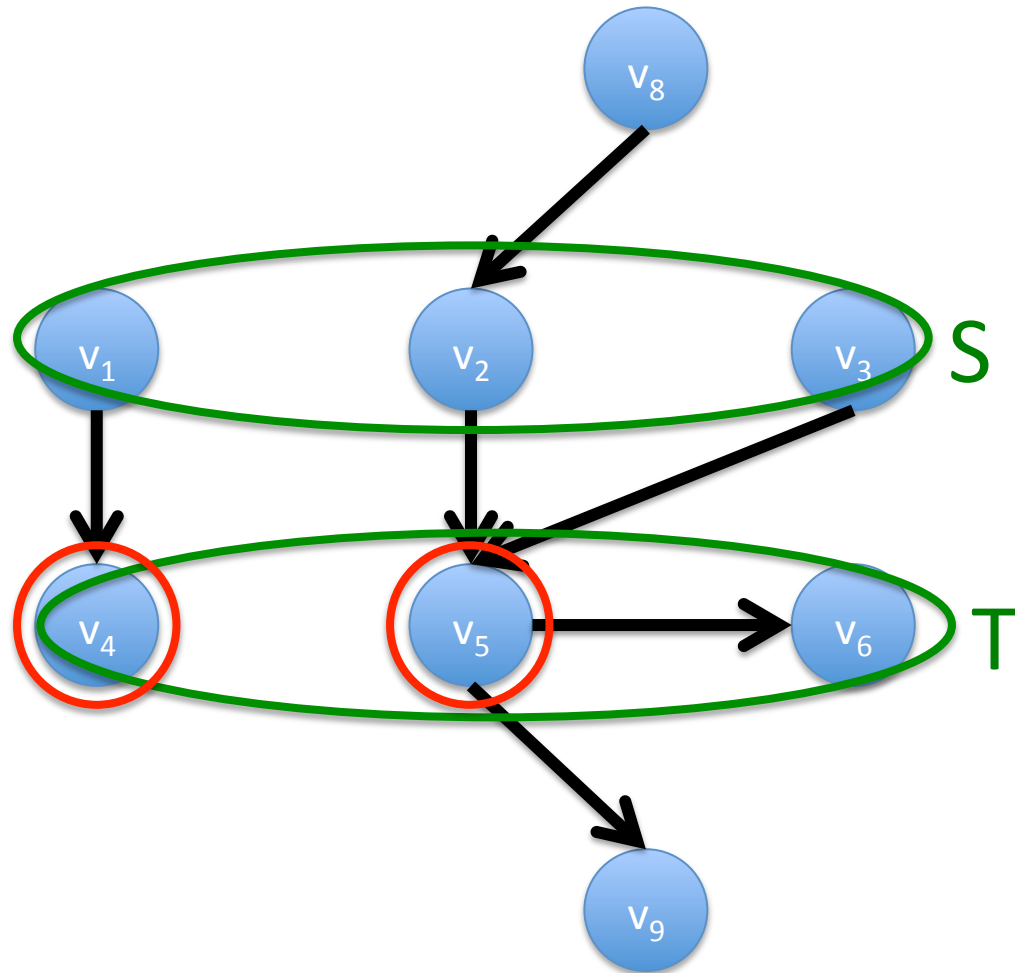
S = outneighbors of s

T = inneighbors of t

Remove s and t

Menger's Theorem for vertex disjoint S-T paths

Menger's Theorem



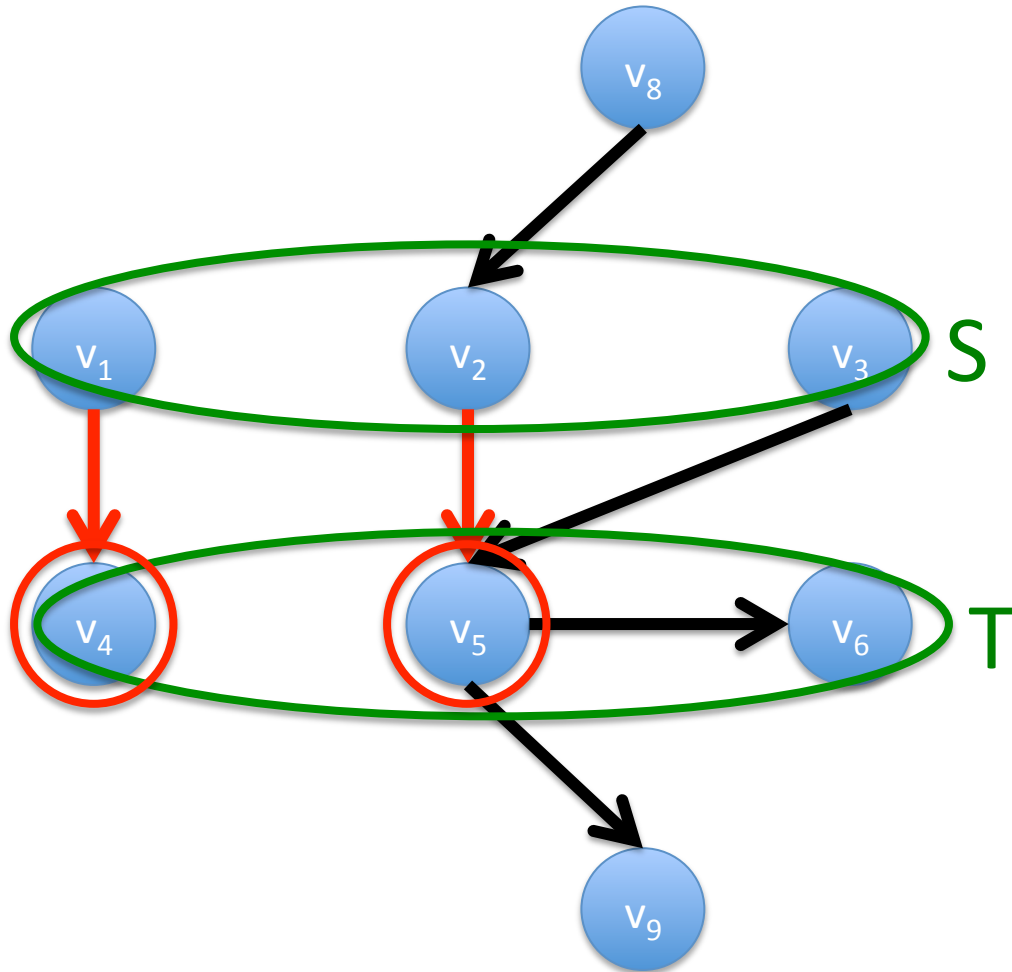
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Menger's Theorem for vertex disjoint S-T paths

Menger's Theorem



S = outneighbors of s

T = inneighbors of t

Remove s and t

Hence proved

Menger's Theorem for vertex disjoint S-T paths

Menger's Theorem

Theorem for vertex disjoint S-T paths

implies

Theorem for internally vertex disjoint s-t paths

Menger's Theorem

Theorem for vertex disjoint S-T paths

is implied by

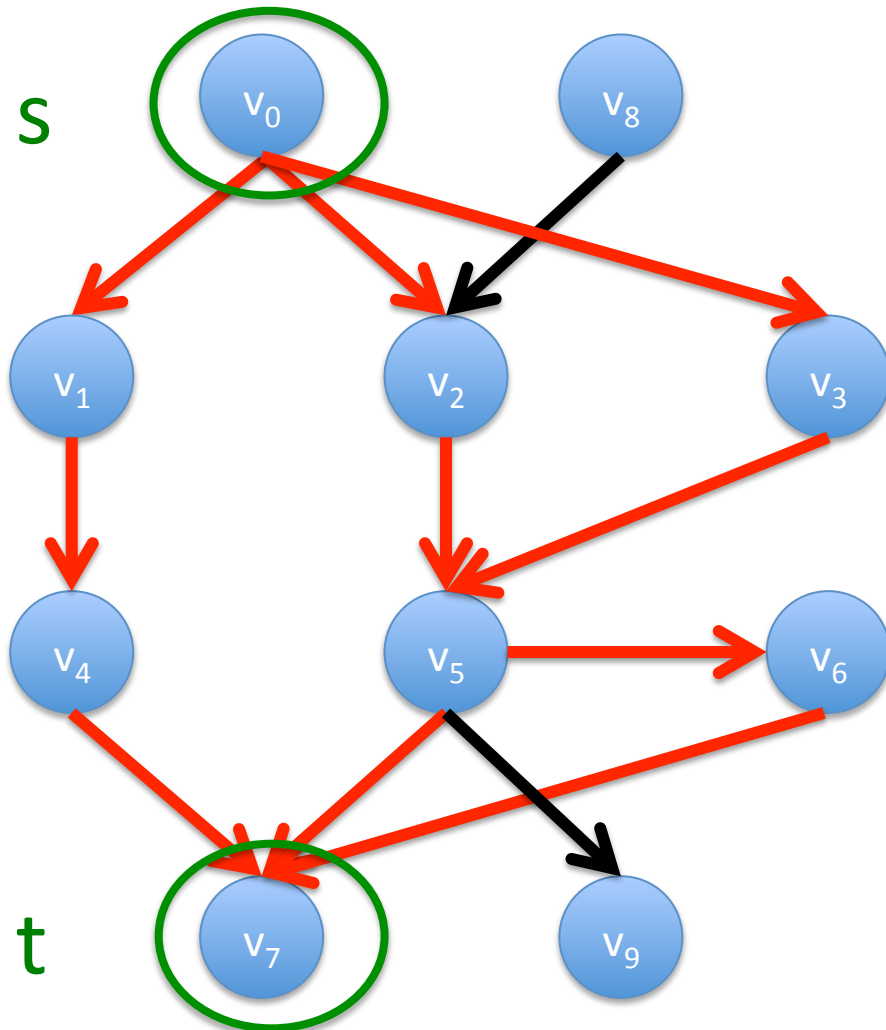
Theorem for internally vertex disjoint s-t paths

Left As Exercise !!!

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- Preliminaries
- Menger's Theorem for Disjoint Paths
 - Vertex Disjoint S-T Paths
 - Internally Vertex Disjoint s-t Paths
 - **Arc Disjoint s-t Paths**
- Path Packing

Arc Disjoint s-t Paths

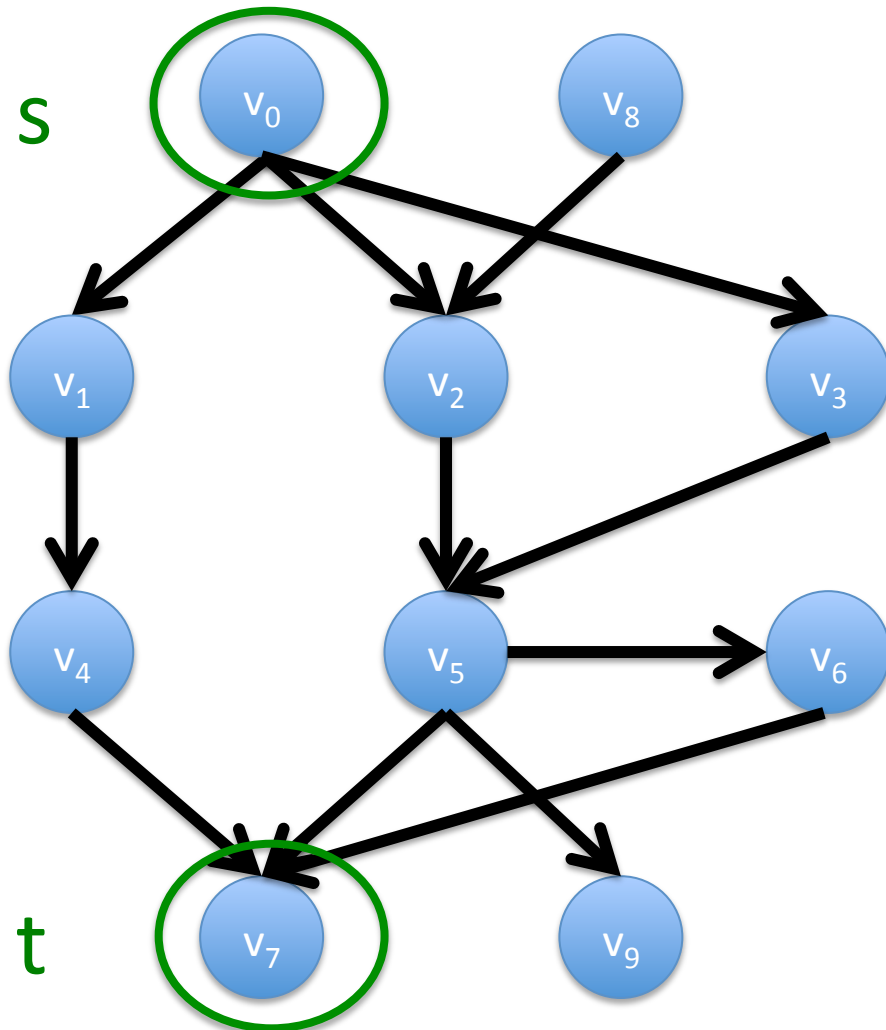


Maximum number of disjoint paths?

Minimum size of s-t cut !!

Set of s-t Paths with no common arcs

s-t Cut

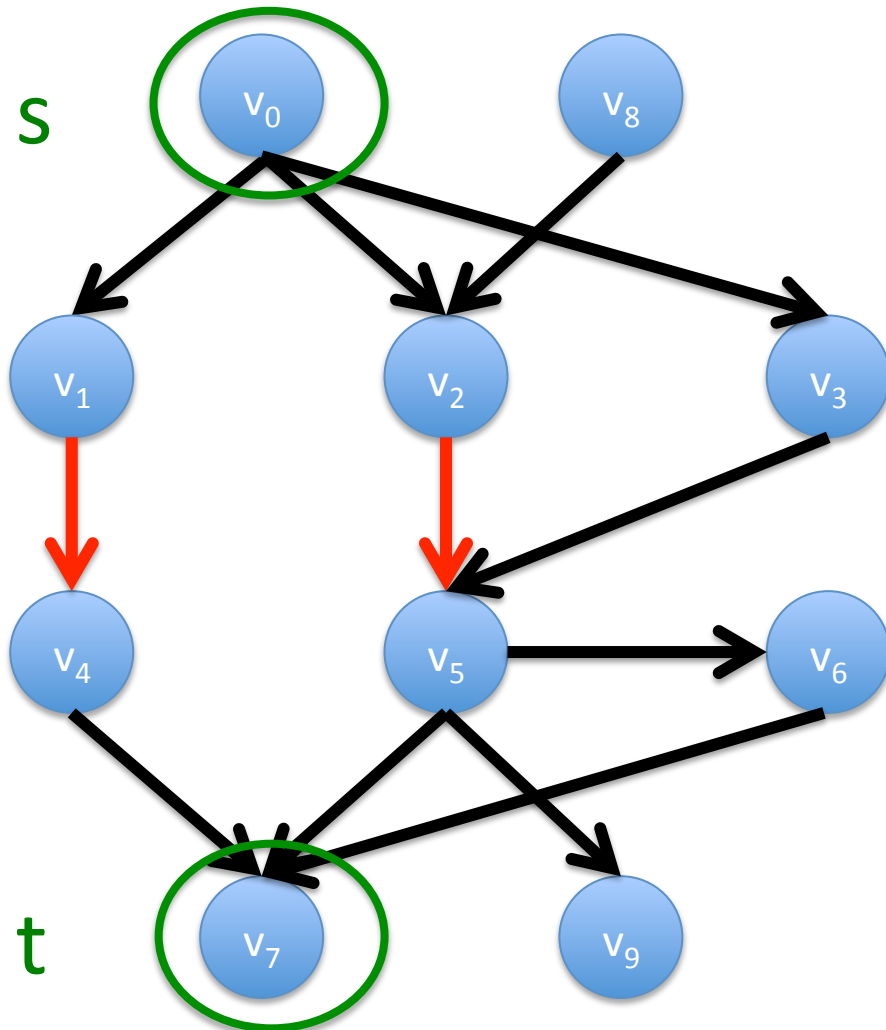


$C = \text{out-arcs}(U)$

$s \in U$

$t \notin U$

s-t Cut

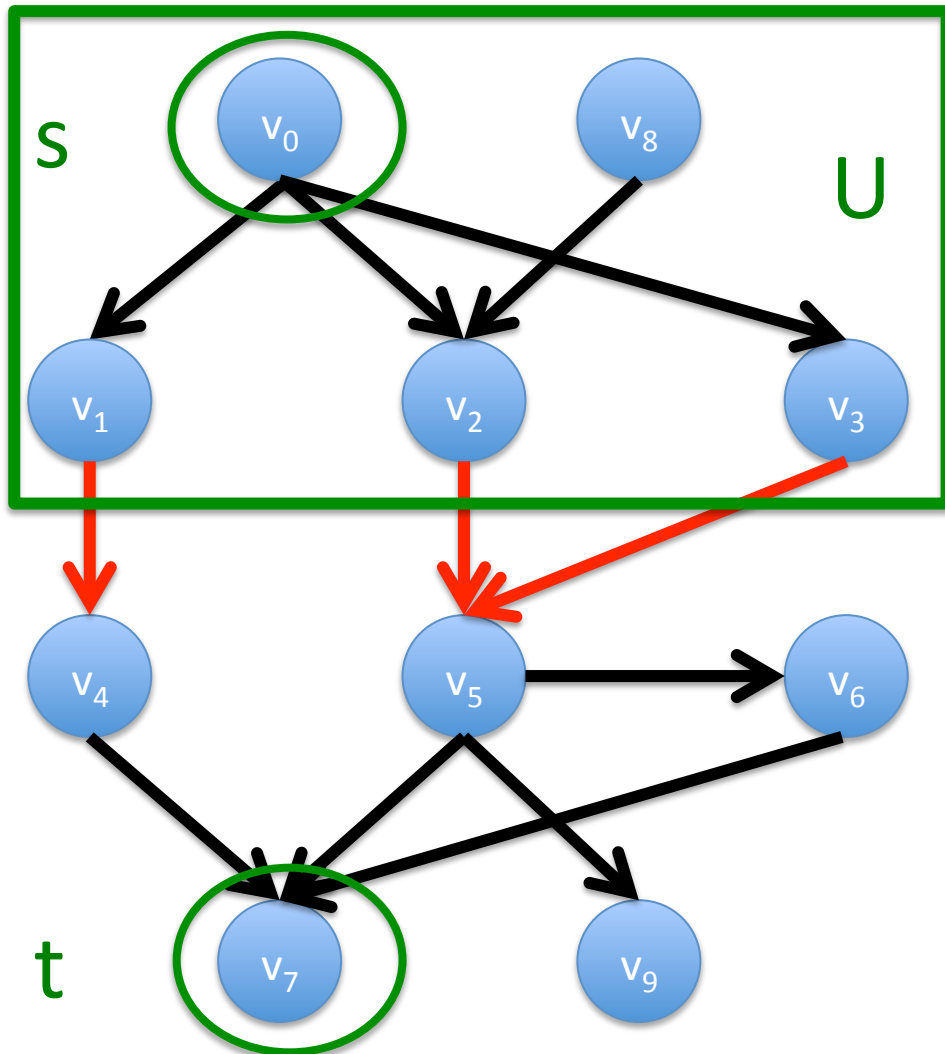


$C = \text{out-arcs}(U)$

$s \text{ in } U$

$t \text{ not in } U$

s-t Cut

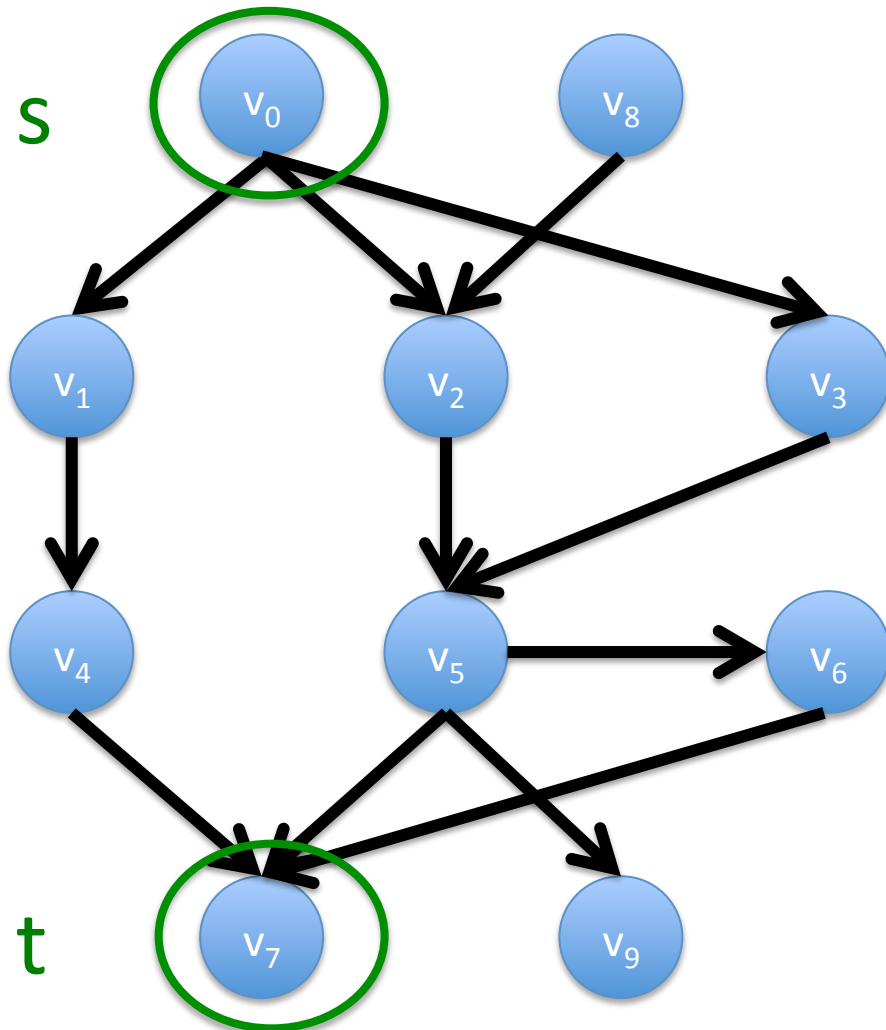


$C = \text{out-arcs}(U)$

$s \text{ in } U$

$t \text{ not in } U$

Connection



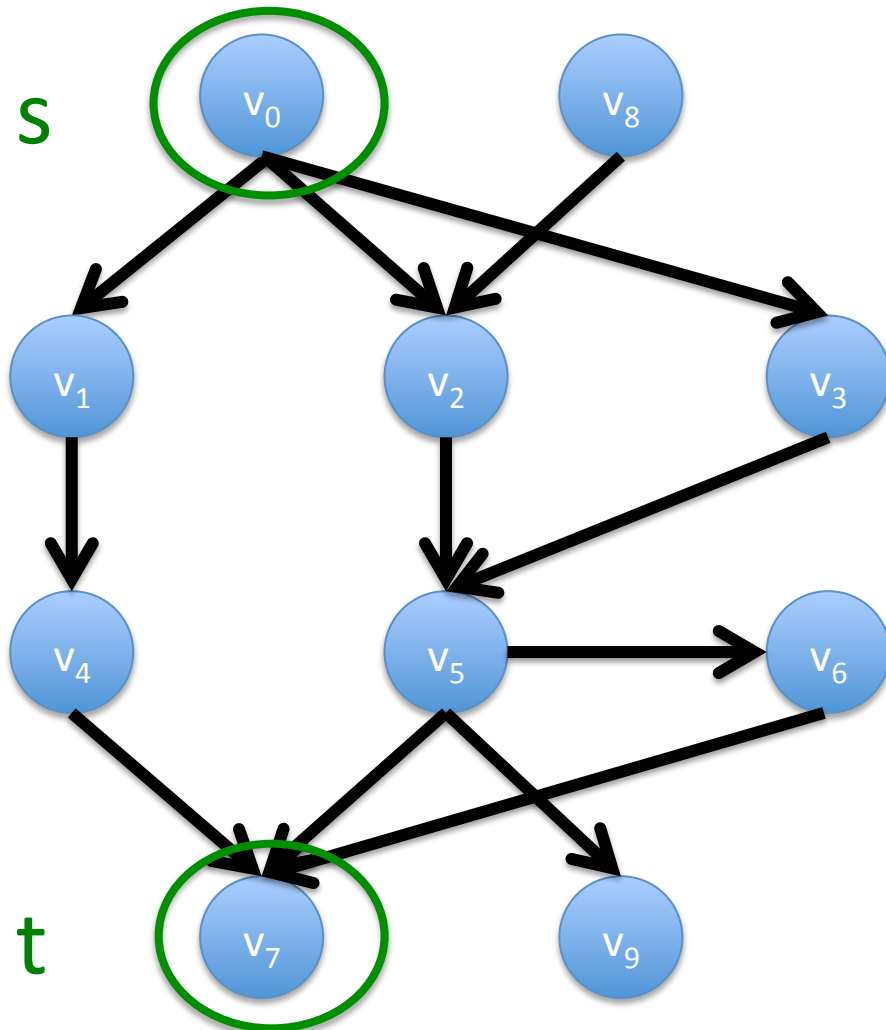
Maximum number of disjoint paths?

\leq

Minimum size of s-t cut !!

Minimum set of arcs intersecting all s-t paths is a cut

Menger's Theorem



Maximum number of
disjoint paths?

=

Minimum size of
s-t cut !!

Proof ?

Menger's Theorem for vertex disjoint S-T paths

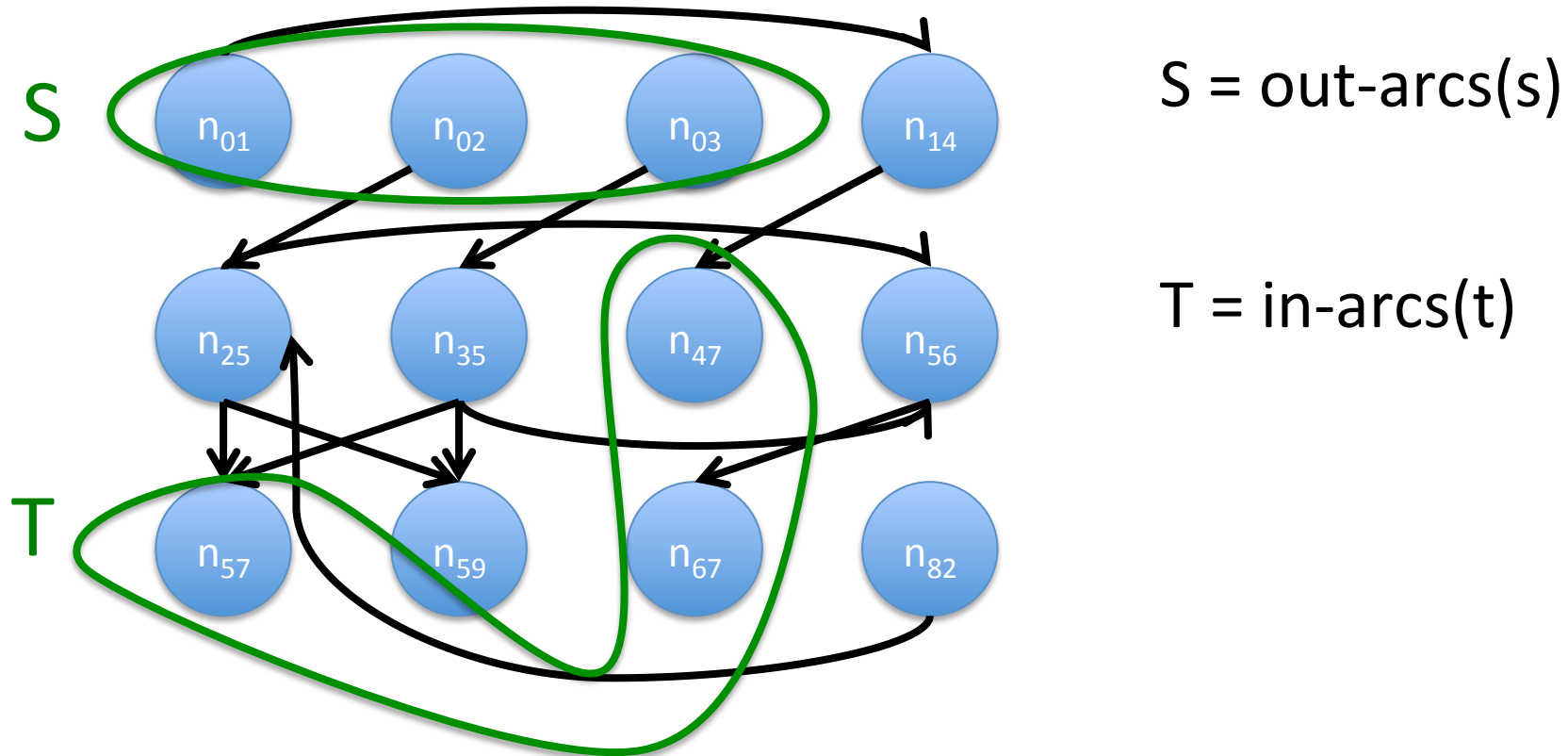
Line Digraph

$L(D)$ is a line digraph of D

Node n_{ij} of $L(D)$ corresponds to arc (v_i, v_j) in D

Arc (n_{ij}, n_{kl}) exists if and only if $j = k$

Menger's Theorem

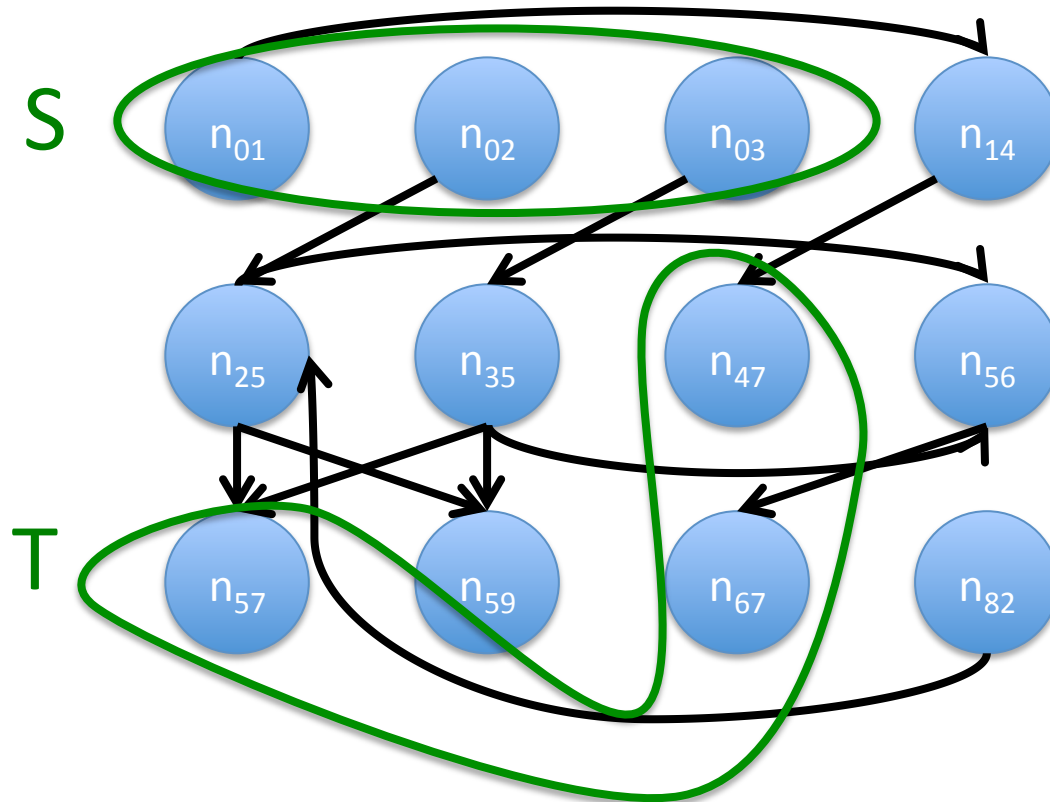


Vertex disjoint S-T path in $L(D)$

implies

Arc disjoint s-t path in D

Menger's Theorem



$S = \text{out-arcs}(s)$

$T = \text{in-arcs}(t)$

Hence proved

Minimum size S-T disconnecting vertex set

implies

s-t cut

Menger's Theorem

Theorem for vertex disjoint S-T paths

implies

Theorem for arc disjoint s-t paths

Menger's Theorem

Theorem for vertex disjoint S-T paths

is implied by

Theorem for arc disjoint s-t paths

Left As Exercise !!!

Outline

- Preliminaries
- Menger's Theorem for Disjoint Paths
- **Path Packing (todo)**
 - Description of the Algorithm
 - Analysis of the Algorithm