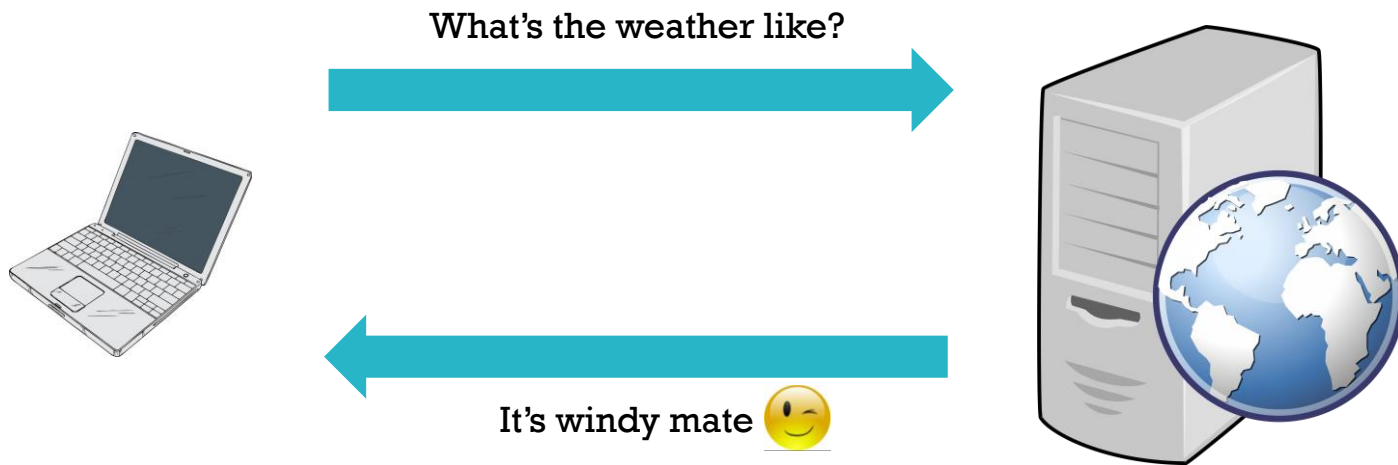


LIFE ON THE WEB

- A very common paradigm in Web development: APIs
 - I send some query to the server
 - The server replies with some data



HOW DOES THIS DATA LOOK LIKE?

- More often than not it is in JSON



```
{"city": "Chicago", "when": "now"}
```



```
{  
  "city": "Chicago",  
  "timestamp": "15-05-17-08-08",  
  "temp": 22,  
  "description": "windy"  
}
```



BOTTOM LINE

- JSON is used a lot:
 - Data exchange on the Web (APIs and the like)
 - Part of programming languages (Python, Ruby, JavaScript)
 - NoSQL databases (MongoDB, Neo4j)
- The how come:
 - Almost no research on the JSON data format and its usage
 - Very few studies available
 - No formal data model/query language
 - Some preliminary work on schema specification



JSON: DATA MODEL, (QUERY LANGUAGES) AND SCHEMA SPECIFICATION

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Thank to Domago Vrgoc for his slides

WHAT IS JSON?

- A fully compositional data format with:
 - Strings and numbers as atomic elements (and some other stuff)
 - Arrays allowing grouping
 - Objects/dictionaries allowing nesting

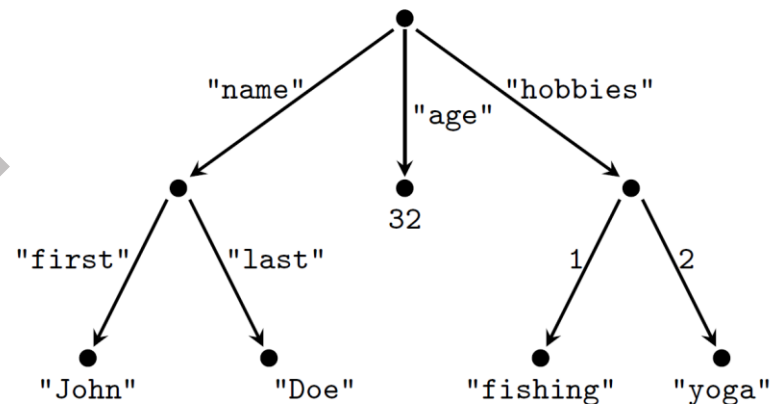
```
{  
  "name": {  
    "first": "John",  
    "last": "Doe"  
  },  
  "age": 32,  
  "hobbies": ["fishing", "yoga"]  
}
```



DATA MODEL FOR JSON – JSON TREES

- JSON document = a set of key-value pairs
 - Each value is again a JSON object
- Naturally suggests tree structure

```
{  
  "name": {  
    "first": "John",  
    "last": "Doe"  
  },  
  "age": 32,  
  "hobbies": ["fishing", "yoga"]  
}
```



SOME SPECIFICS OF JSON TREES

- Inherent typing
 - Each node is either an object or an array
 - Leaves are atomic values
- Labelled edges
 - Allows retrieving values deterministically (e.g. Doc["name"])
- Array nodes have random access to their children
- Fully compositional
 - Value of a node is always a valid JSON tree



HOW IS THIS DIFFERENT FROM XML?

- Main differences:
 - JSON trees are deterministic
 - Allows direct access
 - Value of a node is again a JSON tree
 - Comparing values amounts to comparing subtrees, not atoms
 - JSON trees mix ordered and unordered data
 - Random access for arrays
- Can we code JSON as XML?
 - Sure, but this might be an overkill
 - One of the main reasons JSON got popular = not to do this



SCHEMA DEFINITION FOR JSON

- We want to specify what sort of data our JSON has
- IETF has a draft proposal for JSON Schema
- Heavily studied in
 - Pezoa, et.al. *Foundations of JSON Schema*, WWW 2016
- A friendly introduction available at:
<http://cswr.github.io/JsonSchema/>



SCHEMA DEFINITION FOR JSON

```
{  
  "type": "object",  
  "properties": {  
    "first_name": { "type": "string" },  
    "last_name": { "type": "string" },  
    "age": { "type": "integer" },  
    "club": {  
      "type": "object",  
      "properties": {  
        "name": { "type": "string" },  
        "founded": { "type": "integer" }  
      },  
      "required": ["name"]  
    }  
  },  
  "required": ["first_name", "last_name", "age", "club"]  
}
```

```
{  
  "first_name": "Alexis",  
  "last_name": "Sanchez",  
  "age": 27,  
  "club": {  
    "name": "Arsenal FC",  
    "founded": 1886  
  }  
}
```



SCHEMA DEFINITION FOR JSON

- To capture JSON Schema we introduce JSON Schema logic (JSL)

$$\varphi, \psi := \top \mid \neg\varphi \mid \varphi \wedge \psi \mid \varphi \vee \psi \mid \psi \in \text{NodeTests} \mid \\ \square_e\varphi \mid \square_{i:j}\varphi \mid \diamond_e\varphi \mid \diamond_{i:j}\varphi$$

- NodeTests take care of basic stuff (typing and matching):
 - E.g. $\text{Int} \wedge \text{Min}(i)$ – a number greater than i
- Modal operators take care of arrays, objects and nesting:
 - E.g. $\square_{1: +\infty} \text{Str}$ – all elements of my array are strings



HOW DOES THIS WORK?

```
{  
  "type": "object",  
  "properties": {  
    "first_name": { "type": "string" },  
    "last_name": { "type": "string" },  
    "age": { "type": "integer" },  
    "club": {  
      "type": "object",  
      "properties": {  
        "name": { "type": "string" },  
        "founded": { "type": "integer" }  
      },  
      "required": ["name"]  
    },  
  },  
  "required": ["first_name", "last_name", "age", "club"]  
}
```

```
{  
  "first_name": "Alexis",  
  "last_name": "Sanchez",  
  "age": 27,  
  "club": {  
    "name": "Arsenal FC",  
    "founded": 1886  
  }  
}
```

$\text{Obj} \wedge \diamond_{\text{first_name}} \text{Str} \wedge \diamond_{\text{last_name}} \text{Str} \wedge \diamond_{\text{age}} \text{Int}$
 $\diamond_{\text{club}} (\text{Obj} \wedge \diamond_{\text{name}} \text{Str} \wedge \square_{\text{founded}} \text{Int})$



WHAT CAN OUR LOGIC DO?

- In terms of expressive power:

Theorem:

JSL captures the *core* JSON Schema

- For every schema there is an equivalent formula and vice versa

- In terms of algorithmic properties:

Theorem:

a) Testing if JSL formula is true on a JSON is in $O(|T|^{2x}|\varphi|)$

- Drops to $O(|T|x|\varphi|)$ when uniqueness constraint is not permitted

b) The satisfiability problem for JSL is in EXPSPACE

- It is PSPACE-complete when uniqueness constraint is not permitted



CAPTURING FULL JSON SCHEMA

- Full JSON Schema allows references and definitions
 - Can get crazy if not *well-formed*
 - To capture this we need recursion
- We add this to JSL by allowing datalog-like definitions
 - Example: every path from the root to the leaves is even:

$$\begin{aligned}\gamma_1 &= \square_{\Sigma} * \gamma_2 \\ \gamma_2 &= (\diamond_{\Sigma} * \top) \wedge (\square_{\Sigma} * \gamma_1) \\ \gamma_1\end{aligned}$$



RECURSIVE JSL

- We can now capture all of (well-formed) JSON Schema

Theorem:

Recursive JSL captures well-formed JSON Schema

- Algorithmically we get a jump:

Theorem:

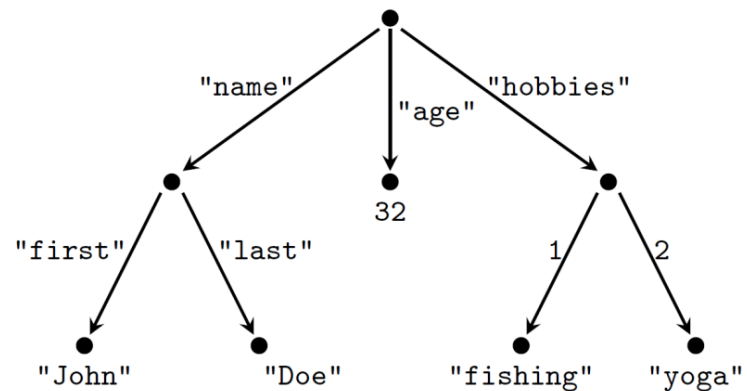
- a) Evaluating a recursive JSL formula is PTIME-complete
- b) The satisfiability problem for recursive JSL is in 2EXPTIME
 - It is EXPTIME-complete when uniqueness constraint is not allowed



NAVIGATIONAL QUERIES OVER JSON

- Inspired by:
 - Python (dictionaries), MongoDB (find function)
 - JSONPath (non-determinism, recursion)

`Doc["hobbies"][1]`



- Compare values (and retrieve):
`Doc["name"] == {"first": "John", "last": "Doe"}`



JSON NAVIGATIONAL LOGIC (JNL)

- A query language capturing this

$$\alpha, \beta := \langle \varphi \rangle \mid X_w \mid X_i \mid \alpha \circ \beta \mid \varepsilon$$

$$\varphi, \psi := \top \mid \neg \varphi \mid \varphi \wedge \psi \mid \varphi \vee \psi \mid [\alpha] \mid EQ(\alpha, A) \mid EQ(\alpha, \beta)$$

- Our previous queries:
 - $X_{\text{hobbies}} \circ X_1$
 - $EQ(X_{\text{name}}, \{ \text{"first": "John", "last": "Doe"} \})$



ALGORITHMIC PROPERTIES OF JNL

Main problems:

- Evaluation:
 - Input: A JSON tree T , a node n , a formula φ
 - Question: Will φ return n when evaluated over T ?
- Satisfiability:
 - Input: A JNL formula φ
 - Question: If there a JSON tree where φ will return at least one node?

Theorem:

- a) Evaluation of JNL can be solved in time $O(|T| \times |\varphi|)$
- b) Satisfiability is NP-complete (even with no negation or EQ).



EXTENDING JNL

- JSONPath allows non-determinism and recursion – JNL*

$$\alpha, \beta := \langle \varphi \rangle \mid X_e \mid X_{i:j} \mid \alpha \circ \beta \mid (\alpha)^* \mid \varepsilon$$
$$\varphi, \psi := \top \mid \neg \varphi \mid \varphi \wedge \psi \mid \varphi \vee \psi \mid [\alpha] \mid EQ(\alpha, A) \mid EQ(\alpha, \beta)$$

Theorem:

- Evaluating JNL* is $O(|T|^3 \mathbf{x} |\varphi|)$
- Without $EQ(\alpha, \beta)$ evaluating JNL* is $O(|T| \mathbf{x} |\varphi|)$
- Satisfiability of JNL* is undecidable
- Without $EQ(\alpha, \beta)$ satisfiability of JNL* is EXPSPACE-complete.
- Without $EQ(\alpha, \beta)$ and without recursion it is PSPACE-complete.



NAVIGATIONAL QUERIES OVER JSON

- Can be shown same to JSON Schema logic
- Check our paper for this 😊



FUTURE DIRECTIONS

Most interesting questions are driven by practical use-cases:

- MongoDB – a full fledged JSON database
- Streams
- Generating API documentation



THANK YOU!

