

# Spatial Birth and Death Processes with Poisson Birth and Shot-Noise Deaths

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# SPATIAL POINT PROCESSES WITH POISSON BIRTHS AND DEATHS BY RANDOM CONNECTION

## Geometry, Point Processes, Dynamics & Content Delivery Networking

### I. Mathematical Analysis

1. Stochastic Model
2. Stochastic Stability
3. Repulsion
4. Birth Death Equations

### II. TCP Example

1. Dimensional Analysis
2. Limiting Regimes
3. Simulation
4. Super-scalability

Joint work with **F. Mathieu** and **I. Norros**, *Queueing Systems*, 2017

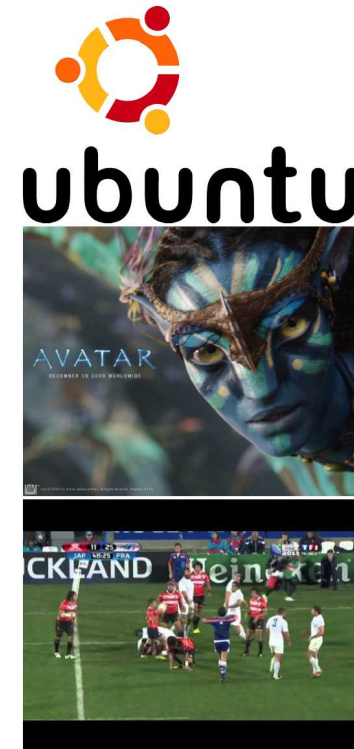
## PEER-TO-PEER CONTENT DISTRIBUTION

Content distribution is mainly:

- ▶ Filesharing
- ▶ Streaming
  - ▶ OnDemand (YouTube, Netflix)
  - ▶ Live (Sport events)

Things in common:

- ▶ Bandwidth is a key parameter
- ▶ Lot of stress for the network
- ▶ P2P solutions exist



**PEER-TO-PEER CONTENT DISTRIBUTION** (*continued*)**■ P2P Principles**

- Peers **join and leave** aiming each at downloading a very large file
- The file is cut in smaller **chunks**
- Peers exchange chunks on a **Tit for Tat** basis
- The swarm of peers solves the initialization problem
- The **bit rate** between two peers is determined by their **distance**

## I. SPATIAL BIRTH AND DEATH STOCHASTIC MODEL

- Nodes live in a subset  $D$  of the Euclidean space  $\mathbb{R}^d$  or on a torus
- Dynamics: **arrivals**
  - **Poisson rain**: new nodes arrive according to a Poisson process with time space intensity  $\lambda dxdt$  on  $D \times \mathbb{R}$
- **Service requirement**: each node  $p$  is born with an individual service requirement  $F_p > 0$  i.i.d. exponential with mean  $F$

## INTERACTION

### ■ Dynamics: **service rate**

- **Bit rate function:** two nodes at locations  $\mathbf{x}$  and  $\mathbf{y}$  serve each other at rate  $\mathbf{f}(\|\mathbf{x} - \mathbf{y}\|) \geq 0$
- **Service rate:** the service rate of a node at  $\mathbf{x}$  in configuration  $\phi$  is

$$\mu(\mathbf{x}, \phi) = \sum_{\mathbf{y} \in \phi \setminus \{\mathbf{x}\}} \mathbf{f}(\|\mathbf{x} - \mathbf{y}\|)$$

- **Service completion:** for a system with state history  $\{\phi_t\}_t$ , a node  $p$  born at point  $\mathbf{x}_p$  at time  $t_p$  leaves at time

$$\tau_p = \inf\left\{t > t_p : \int_{t_p}^t \mu(\mathbf{x}_p, \phi_s) ds \geq F_p\right\}$$

## SPATIAL BIRTH AND DEATH PROCESS

- $\mathcal{N}(\mathbf{D})$ : the space of counting measures in  $(\mathbf{D}, \mathcal{D})$
- In the finite domain case, the state  $\phi_t$  at time  $t$  is a **Markov process** living in the space  $\mathcal{N}(\mathbf{D})$ :
  - a node has **birth intensity**  $\lambda$  at  $\mathbf{x}$
  - a node located at  $x$  has **death intensity**  $\mu(\mathbf{x}, \phi_t)/F$
- **Spatial birth-and-death process** with a death rate defined as a **shot-noise** of the configuration

## CONSTRUCTION: FINITE CASE

### ■ Lemma

If  $D$  is compact and  $f$  is bounded from below by a positive constant on some non-degenerate interval, then the Markov process  $\{\phi_t\}_t$  is ergodic for any birth rate  $\lambda > 0$

### ■ Proof

- stochastic domination:  $M/M/p$  queue that is modified so that a lone customer cannot leave
- Harris-recurrence techniques

### ■ Remarks: in general

- non monotonic dynamical system



## THE $M/M/p2p$ SYSTEM

**Consider the following queueing system:**

- customers (‘peers’) arrive according to a Poisson process with parameter  $\lambda$
- every peer has an independent  $\text{Exp}(1)$  distributed service requirement, and each peer serves every other peer at rate  $\mu$ ; that is, if there are  $n$  peers in the system, each of them has stochastic intensity  $(n - 1)\mu$  to leave

**After arrival of first peer, the system is never empty**

**THE  $M/M/p2p$  SYSTEM** (continued)

**The system is a birth-death process with balance equations**

$$\pi_n \lambda = \pi_{n+1} \mu (n+1) n,$$

**whose solution is the queue's stationary probability measure**

$$\mathbb{P}[Q = n] = \pi_n = \frac{(\lambda/\mu)^{n-1}}{n!(n-1)!} \bigg/ \sum_{k=1}^{\infty} \frac{(\lambda/\mu)^{k-1}}{k!(k-1)!}, \quad n \geq 1$$

**The infinite sums are expansions of Bessel functions:**

$$\sum_{n=0}^{\infty} \frac{x^n}{n!n!} = I_0(2\sqrt{x}), \quad \sum_{n=0}^{\infty} \frac{x^n}{(n+1)!n!} = \frac{1}{\sqrt{x}} I_1(2\sqrt{x})$$

**In particular, the mean number of peers in system is**

$$\mathbb{E}Q = \frac{I_0(2\sqrt{\lambda/\mu})}{I_1(2\sqrt{\lambda/\mu})} \sqrt{\frac{\lambda}{\mu}}$$

## CONSTRUCTION: INFINITE CASE

- **Definition over finite time?**

- **Assume  $F = 1$**

- **Proposition**

If  $D = \mathbb{R}^d$  and

$$\int_1^{\infty} f(r)r^{d-1} dr < \infty$$

then the spatial birth and death point process is uniquely defined on all finite time intervals  $[t_0, t]$

## PROOF: A-GRAPHICAL REPRESENTATION

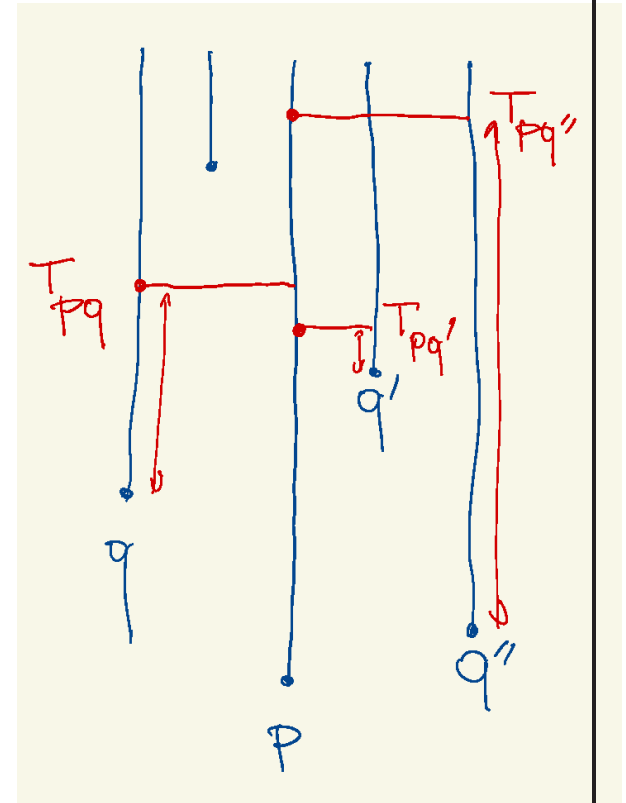
- $u$  fixed,  $t_0 < u$
- $\Psi_{t_0}$ : space-time arrival P.P.P. in  $[t_0, u]$
- Point  $p = (x_p, t_p)$  of  $\Psi_{t_0}$
- Graphical representation of Shot-Noise

For all pairs  $p \neq q \in \Psi_{t_0}$   
Killing times  $T_{pq}$

$$T_{pq} = T_{qp} \sim (t_p \vee t_q) + \mathbf{Exp}(2f(\|x_p - x_q\|))$$

Bernoulli directions of killing  $I_{pq}$

$$I_{pq} = 1 - I_{qp} \sim \mathbf{Bernoulli}\left(\frac{1}{2}\right)$$



**PROOF: A-GRAPHICAL REPRESENTATION** (continued)**■ Local finiteness in graphical construction****■ Lemma If**

$$\int_1^{\infty} f(\mathbf{r}) r^{d-1} dr < \infty,$$

**almost surely, none of the sets**

$$N_p = \{T_{pq} : q \in \Phi_{t_0}\}$$

**has accumulation points**

**■ Proof based on degree properties of the random connection model**

PROOF: A-GRAPHICAL REPRESENTATION (continued)

■ Let  $p = (\mathbf{x}, \mathbf{t}) \in \Psi_{\mathbf{t}_0}$

$$\begin{aligned}
 \mathbb{E}^p |\mathbf{N}_p \cap (\mathbf{t}_0, \mathbf{u}]| &= \mathbb{E} \int_{\mathbb{R}^d \times (\mathbf{t}_0, \infty)} \mathbf{1}_{T_{pq} \leq \mathbf{u}}(\Psi_{\mathbf{t}_0} - \delta_p)(d\mathbf{q}) \\
 &= \mathbb{E} \int_{\mathbb{R}^d \times (\mathbf{t}_0, \mathbf{u}] } \mathbf{1}_{T_{pq} \leq \mathbf{u}}(\Psi_{\mathbf{t}_0})(d\mathbf{q}) \\
 &= \lambda \int_{\mathbb{R}^d} \int_{\mathbf{t}_0}^{\mathbf{u}} \mathbb{P}[\mathbf{Exp}(f(\|\mathbf{x} - \mathbf{y}\|)) \leq \mathbf{u} - (\mathbf{t} \vee \mathbf{v})] d\mathbf{v} d\mathbf{y} \\
 &\leq \lambda(\mathbf{u} - \mathbf{t}_0) \int_{\mathbb{R}^d} \left( \mathbf{1} - e^{-(\mathbf{u} - \mathbf{t}_0)f(\|\mathbf{y}\|)} \right) d\mathbf{y} \\
 &\leq \lambda(\mathbf{u} - \mathbf{t}_0) \left( \nu_d + (\mathbf{u} - \mathbf{t}_0) \int_{\mathbb{R}^d \setminus \mathbf{B}(\mathbf{0}, 1)} f(\|\mathbf{y}\|) d\mathbf{y} \right) < \infty
 \end{aligned}$$

## PROOF: B-CONSTRUCTION ALGORITHM

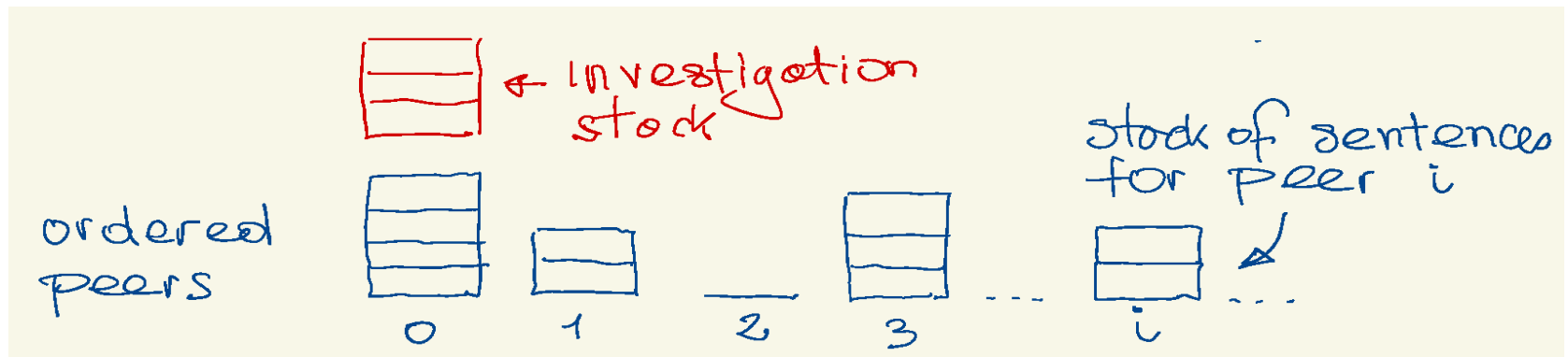
- Not all killing epochs lead to death: only living peers matter
- Death times solution of the **infinite recursive equation**

$$\delta_p = \inf \{ T_{pq} : q \in \Psi_{t_0}, \delta_q \geq T_{pq}, I_{pq} = 1 \}$$

- **The Construction Algorithm**  
gives the solution of this recursive equation on compacts of time
- **Principle**

pick a node, check its earliest killing time;  
determine whether the killer's death time is earlier or later than this  
time...

PROOF: B-CONSTRUCTION ALGORITHM (continued)



**Algorithm : construction of death process on  $(t_0, \infty)$**

**1. Initialization:**

- each peer has a stack of ordered death sentences (peer, time, killer) - earliest on top
- there is a global investigation stack, initially empty



**PROOF: B-CONSTRUCTION ALGORITHM** (continued)

- 2. If investigation stack empty:  
pick first peer with top sentence with time  $< u$  and no certificate;  
move this sentence to the investigation stack;  
if no such sentence exists, stop;**
- 3. Look at the top of investigation stack, say  $(x, s, y)$ , and do**
  - If  $y$ 's stack has on top a death sentence or certificate later than  $s$ , then death happens: change the sentence  $(x, s, y)$  into death certificate with same date and return it to the top of  $x$ 's stack;**
  - If  $y$  has death certificate earlier than  $s$ , the sentence  $(x, s, y)$  is removed from investigation stack and deleted;**
  - Otherwise move top sentence of  $y$ 's stack to investigation stack;**
- 4. Go to 2.**

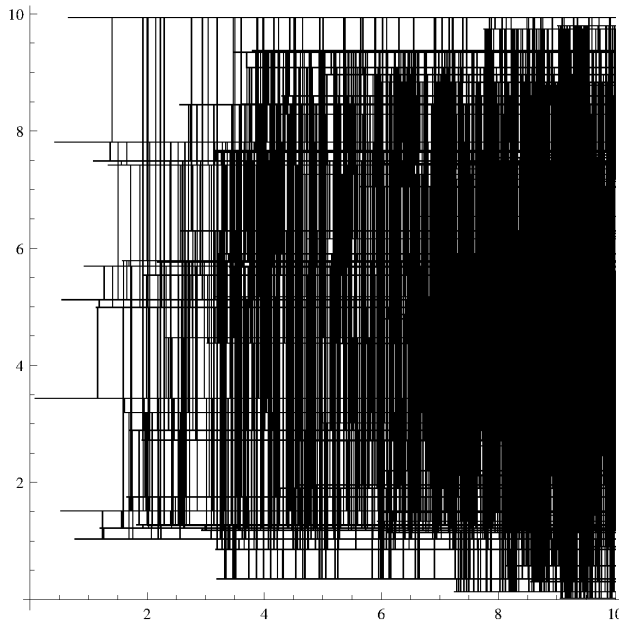
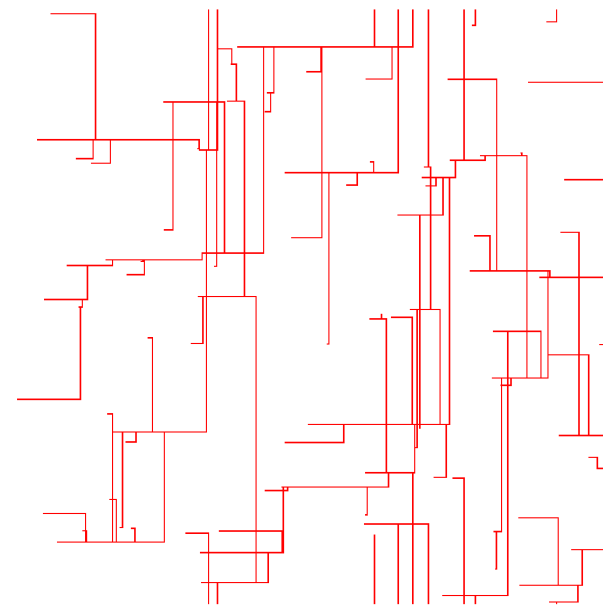
PROOF: B-CONSTRUCTION ALGORITHM (continued)

■ **Lemma** If

$$\int_0^{\infty} f(\mathbf{r}) r^{d-1} dr < \infty,$$

for each sentence, the sequence of peers  $p_1, p_2, \dots$  produced by its investigation is such that  $(T_{p_n p_{n+1}})$  is decreasing and a.s. finite

- **Proof** For  $u - t_0$  small enough, an upper-bound random connection model with connection function  $f$  does not percolate. For more general  $u$ , decompose  $[t_0, u]$  in small intervals and apply the last observation
- **Theorem** For every peer  $p$  born in  $(t_0, \infty)$  the construction algorithm determines a unique death time  $\delta_p \leq t$
- **Corollary** Almost surely, the death process is defined uniquely as a factor of  $\Psi_{t_0}$

**PROOF: B-CONSTRUCTION ALGORITHM** (*continued*)**Graphical Representation****State Construction**

## PROOF: C-CONSTRUCTION PROPERTIES

- When the process is a.s. well-defined, each peer  $p$  is a.s. killed by a uniquely determined peer  $\kappa(p)$
- Since  $\kappa$  is non-cyclic, its graph is a forest of infinite trees
- **Conjecture**  
When  $D = \mathbb{R}^2$ , the directed graph  $(\Psi_{t_0}, \{(p, \kappa(p)) : p \in \Psi_{t_0}\})$  is almost surely a tree
- **Proposition**  
For any peer  $p$ , the conditional distribution of the number of peers it kills, given the history of the process up to time  $t_p$ , is  $\text{Geom}(\frac{1}{2})$
- **Observation**  
This does not show yet that the death process is well-defined when  $t_0 = -\infty$

## II. CONSTRUCTION OF THE STATIONARY REGIME

### ■ Existence/uniqueness of stationary regimes

#### ■ **Theorem** Under the assumptions

(i)  $a := \int_{\mathbb{R}^d} f(\|x\|) dx < \infty$

(ii)  $f$  is non-increasing and bounded

there exists a unique stationary regime holding for all initial conditions made of a homogeneous Poisson point process of initial nodes.

**II. CONSTRUCTION OF THE STATIONARY REGIME** (*continued*)**■ Proof Structure: 2 coupling steps**

- **Coupling 1:** couple the two process defined on  $(0, \infty)$  **with and without** the addition of a P.P. of initial nodes at time 0 and show that the effect of this addition **vanishes at infinity**
- **Coupling 2:** Coupling from the past:  $(0, \infty) \rightarrow (-\infty), 0)$

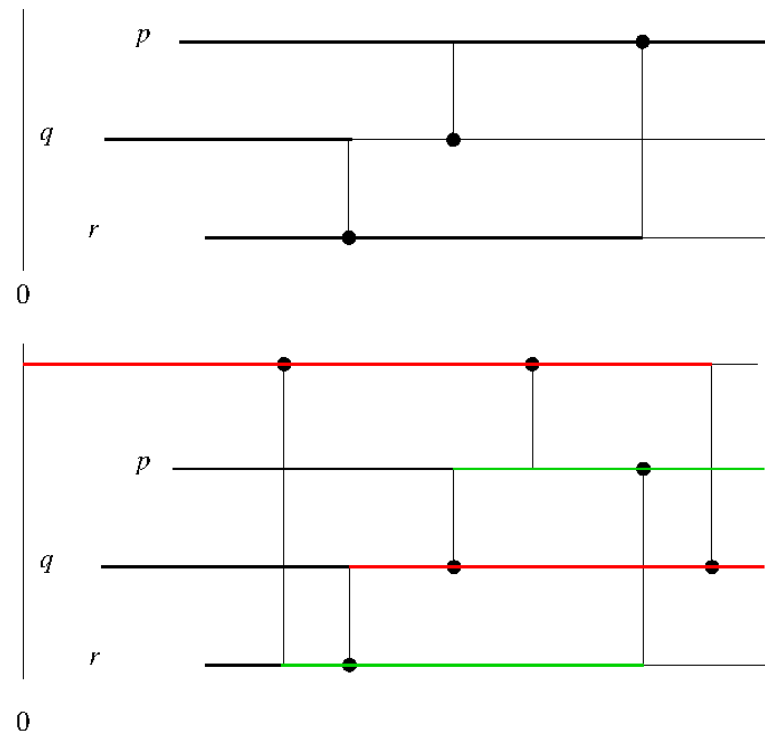
**■ Class of additional point processes  $Z$ :**

**motion-invariant P.P. with a few additional properties, e.g.**

- **homogeneous P.P.P.**
- **dependent thinning of homogeneous P.P.P.**

## COUPLING 1

- **Two parallel histories**
  - **Black + Green**  
without add. point
  - **Black + Red**  
with add. point
- **Built by a variant of the graphical representation and construction algorithm**



**COUPLING 1** (*continued*)**Effect of the addition of an independent homogeneous P.P.P. of peers at time 0 on the two coupled histories on  $(0, \infty)$** **■ Lemma**

Under the assumptions of the theorem, the number of special points (**red** and **green**) created by any single point of the additional P.P.P. is a.s. finite.

**■ Proof** supermartingale argument on cardinality of special progeny of this point**■ Lemma**

Under the assumptions of the theorem, the intensity of special points (**red** and **green**) created by the additional P.P.P. decreases exponentially fast with time

**■ Proof** Abel & Cain argument



COUPLING 1 (continued)

### Proof of First Lemma (Sketch)

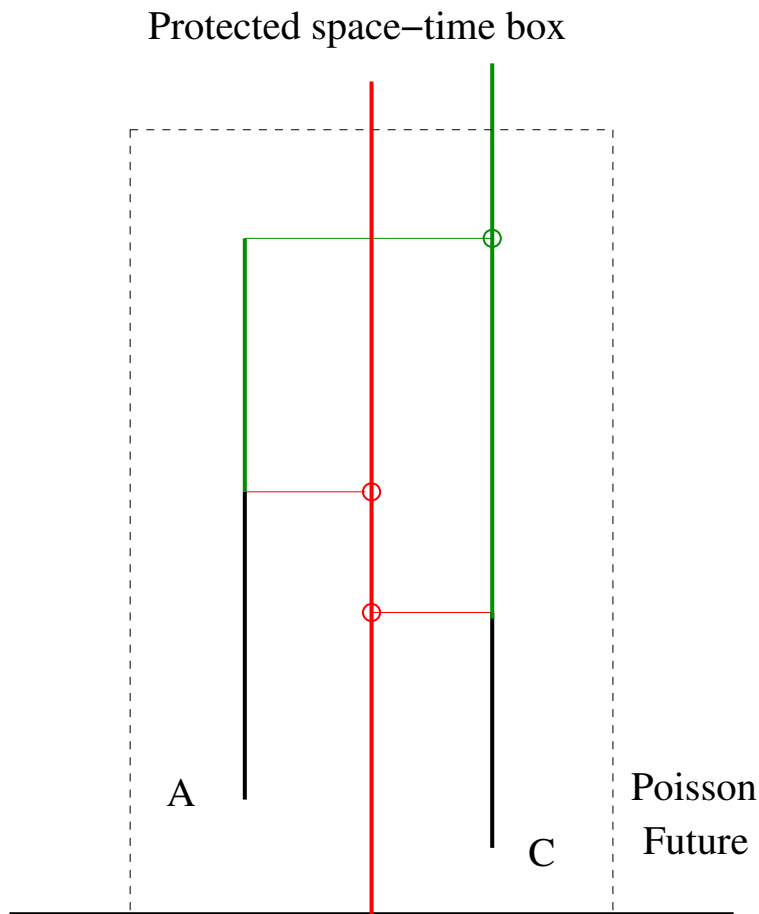
- $S_n(z)$  set of special points stemming from  $z$  at  $n$ -th duel concerning this population
- If duel between **regular and special**:

$$\text{card}(S_n(z)) := \text{card}(S_{n-1}(z)) + J_n$$

with  $J_n \sim +1$  w.p.  $1/2$  and  $-1$  w.p.  $1/2$ , cond. independent on the past

- If duel between **special and special**:
  - either  $\text{card}(S_n(z))$  is unchanged (specials of the same kind)
  - or it evolves as above (different kinds and same progeny)
  - or its evolution is bounded above by the last equation (different kind and different progeny)

COUPLING 1 (continued)



### Proof of Second Lemma (Sketch)

uniformly over time and space, the offspring cardinality of a special point (here **red**) has a strictly negative drift through the **Cain kills Abel** scheme

- This requires proving the Palm expectation of the death pressure on specials is uniformly bounded;
- This implies exponential decrease of the density of special points.

## COUPLING 2: FROM THE PAST

- $\Psi_{s,t}^Z$ : configuration at  $t$  built by Coupling 1 when initial time is  $s$  and initial condition is  $Z$
- **Corollary** of second Lemma  
For all compacts  $K$  of space, there is a finite expectation time  $T_K(s)$  such that for all  $t > T_K(s)$ ,  $\Psi_{s,t}^Z$  has no special points in  $K$
- $\Psi_{s,t}$ : configuration at  $t$  when initial time is  $s$  and initial condition is  $\emptyset$
- From Lemma, for all  $t$

$$\exists \lim_{s \rightarrow -\infty} \Psi_{s,t} = \Phi_t$$

with  $\Phi_t$  translation invariant w.r.t. space and time

- Theorem 1 follows

### 3. f-REPULSION

#### ■ Theorem

For all  $f$  such that there exists a non degenerate translation invariant stationary regime  $\Phi = \Phi_0$ , in the stationary regime,

$$\mathbb{E}\left[\sum_{\mathbf{x}_i \in \Phi} f(\|\mathbf{x}_i\|)\right] \geq \mathbb{E}_0\left[\sum_{\mathbf{x}_i \in \Phi \setminus \mathbf{0}} f(\|\mathbf{x}_i\|)\right]$$

where  $\mathbb{P}_0$  is the Palm probability w.r.t.  $\Phi$

■ **Proof:** rate conservation principle + Papangelou theorem for point processes with stochastic intensity

## SKETCH OF PROOF - TORUS

- $\Phi_t$ : state of the SBD at time  $t$ .
- $\Lambda_t$ : total rate

$$\Lambda_t = \sum_{\mathbf{X} \in \Phi_t} \Lambda_t(\mathbf{X}),$$

with, for all  $\mathbf{X} \in \Phi_t$ :

$$\Lambda_t(\mathbf{X}) = \sum_{\mathbf{Y} \in \Phi_t, \mathbf{Y} \neq \mathbf{X}} f(\|\mathbf{X} - \mathbf{Y}\|)$$

SKETCH OF PROOF - TORUS (continued)

■ **Rate conservation principle applied to  $\mathbb{A}_t$ :**

- $\mathbb{E}^\uparrow$ : (time) Palm probability of the SBD at birth epochs
- $\mathbb{E}^\downarrow$  at death epochs

$$\mathbf{r}^\uparrow \mathbb{E}^+(\mathcal{I}) = \mathbf{r}^\downarrow \mathbb{E}^\downarrow(|\mathcal{D}|)$$

**with**

- $\mathcal{I} = \mathbb{A}_{0+} - \mathbb{A}_0$  the total rate increase,  $\mathbf{r}^\uparrow$  the inc. intensity
- $\mathcal{D} = \mathbb{A}_{0+} - \mathbb{A}_0$  the total rate decrease,  $\mathbf{r}^\downarrow$  the dec. intensity

SKETCH OF PROOF - TORUS (continued)

- Since  $r^\uparrow = r^\downarrow$ ,

$$\mathbb{E}^\uparrow(\mathcal{I}) = \mathbb{E}^\downarrow(\mathcal{D})$$

- From PASTA

$$\mathbb{E}^\uparrow(\mathcal{I}) = 2\mathbb{E}(\mathbf{n}_0) \frac{\mathbf{a}}{|\mathbf{D}|}$$

with  $\mathbf{n}_0$  the total population and

$$\mathbf{a} = \int_{\mathbf{T}} \mathbf{f}(\|\mathbf{x}\|) \mathbf{m}(d\mathbf{x})$$

with  $\mathbf{T}$  the torus of area  $|\mathbf{D}|$

SKETCH OF PROOF - TORUS (continued)

- The (total) death point process admits a stochastic intensity w.r.t. the filtration  $\mathcal{F}_t = \sigma(\Phi_s, s \leq t)$  equal to  $\Lambda_t$
- From **Papangelou's theorem**  $\frac{d\mathbb{P}^\downarrow}{d\mathbb{P}} \Big|_{\mathcal{F}_{0-}} = \frac{\Lambda_0}{\mathbb{E}(\Lambda_0)}$
- Since the decrease (in state  $\Phi_{0-}$ ) is of magnitude  $\Lambda_0(\mathbf{X})$  (w.r.t.  $\Phi_{0-}$ ) with probability  $\frac{\Lambda_0(\mathbf{X})}{\Lambda_0}$  (w.r.t.  $\Phi_{0-}$ ),

$$\begin{aligned} \mathbb{E}^\downarrow(|\mathcal{D}|) &= 2\mathbb{E} \left( \frac{\Lambda_0}{\mathbb{E}(\Lambda_0)} \sum_{\mathbf{X} \in \Phi_0} \frac{\Lambda_0(\mathbf{X})}{\Lambda_0} \Lambda_0(\mathbf{X}) \right) = 2 \frac{\mathbb{E} \left( \sum_{\mathbf{X} \in \Phi_0} (\Lambda_0(\mathbf{X}))^2 \right)}{\mathbb{E} \left( \sum_{\mathbf{X} \in \Phi_0} \Lambda_0(\mathbf{X}) \right)} \\ &= 2 \frac{\mathbb{E}_0 \left( (\Lambda_0(\mathbf{0}))^2 \right)}{\mathbb{E}_0 \left( \Lambda_0(\mathbf{0}) \right)} \end{aligned}$$



**SKETCH OF PROOF - TORUS** (continued)**■ Rate conservation principle for total rate:**

$$\mathbb{E}(\mathbf{n}_0) \frac{\mathbf{a}}{|\mathbf{D}|} = \frac{\mathbb{E}_0((\mathbf{A}_0(\mathbf{0}))^2)}{\mathbb{E}_0(\mathbf{A}_0(\mathbf{0}))}$$

**■ Using the fact that**

$$\mathbb{E}_0((\mathbf{A}_0(\mathbf{0}))^2) \geq \mathbb{E}_0(\mathbf{A}_0(\mathbf{0}))^2,$$

**we get**

$$\mathbb{E}(\mathbf{n}_0) \frac{\mathbf{a}}{|\mathbf{D}|} \geq \mathbb{E}_0(\mathbf{A}_0(\mathbf{0}))$$

## 4. BIRTH AND DEATH EQUATIONS FOR MOMENT MEASURES

- Factorial moment measure densities e.g.

$$\mathbb{E}_{\Phi}^0[\Phi^!(\mathbf{B})] = \frac{1}{\beta} \int_{\mathbf{B}} m_{\text{mi}}^{[2]}(\|\mathbf{x}\|) d\mathbf{x}$$

- **Theorem**

The time stationary SBD satisfies the following balance relations:

$$\int_{\mathbf{R}^d} m_{\text{mi}}^{[2]}(\|\mathbf{x}\|) f(\|\mathbf{x}\|) d\mathbf{x} = \lambda$$

- RCP for first moment measure

#### 4. BIRTH AND DEATH EQUATIONS FOR MOMENT MEASURES *(continued)*

### ■ RCP Hierarchy of Equations for Moment Measures

For all  $k \geq 2$ , for all  $\mathbf{x}_1, \dots, \mathbf{x}_k$  in  $\mathbb{R}^d$ ,

$$\begin{aligned}
 & m^{[k]}(\mathbf{x}_1, \dots, \mathbf{x}_k) \left( \sum_{i=1,k} \sum_{j=1,k, j \neq i} f(\|\mathbf{x}_i - \mathbf{x}_j\|) \right) \\
 & + \int_{\mathbb{R}^d} m^{[k+1]}(\mathbf{x}_1, \dots, \mathbf{x}_k, \mathbf{z}) \left( \sum_{i=1,k} f(\|\mathbf{x}_i - \mathbf{z}\|) \right) d\mathbf{z} \\
 & = \lambda \sum_{i=1,k} m^{[k-1]}(\mathbf{x}_1, \dots, \mathbf{x}_{i-1}, \mathbf{x}_{i+1}, \dots, \mathbf{x}_k)
 \end{aligned}$$

## II. TCP EXAMPLE

- **TCP model:**  $D$  is the Euclidean plane  $\mathbb{R}^2$

$$f(\mathbf{r}) = \frac{C}{r} \mathbf{1}_{r \leq R}$$

- [Mathis, Semke, Mahdavi & Ott, 97]
- Most results extend to the general  $f$  case

## 1. DIMENSIONAL ANALYSIS

### ■ 4 basic parameters:

- $R$  in meters (m),
- $C$  in  $\text{bit} \cdot \text{m} \cdot \text{s}^{-1}$ ,
- $\lambda$  in  $\text{m}^{-2} \cdot \text{s}^{-1}$ ,
- $F$  in bits.

### ■ $\pi$ -Theorem

In the TCP case, all system properties only depend on the dimensionless parameter

$$\rho = \frac{\lambda F R^3}{C}$$

### ■ Extension for more general $f$ s.t. $\int f(r) r dr < \infty$

## 1. DIMENSIONAL ANALYSIS (continued)

### ■ Sketch of proof

- choose  $R$  as a new distance unit, then
  - \* the arrival intensity becomes  $l = \lambda R^2$
  - \* the download speed constant becomes  $c = C/R$
- now define  $F$  as an information unit, then
  - \* the download speed constant becomes  $c = C/(RF)$
- take a time unit such that the download speed constant is 1, then
  - \* all parameters are equal to 1
  - \* the arrival intensity becomes  $l = \frac{\lambda FR^3}{C}$

**1. DIMENSIONAL ANALYSIS** (*continued*)**■ Terminology: Three cases**

- $\rho \gg 1$  is called **fluid**
- $\rho \ll 1$  is called **hard core**
- $\rho$  inbetween is called **intermediate**

## NOTATION

■ **In the steady state regime of the P2P dynamics:**

- $\beta_o$  the density of the node point process
- $\mu_o$  the mean rate of a typical node
- $W_o$  the mean latency of a typical node
- $N_o$  the mean number of nodes in a ball of radius  $R$  around a typical node



## 2. FIRST ORDER APPROXIMATION

- **Fluid, or Poisson heuristic:** obtained when ignoring the Palm expectation of the rate, and using the mean rate at a typical location instead:

$$\mu_f = \beta_f 2\pi \int_{r=0}^R (C/r) r dr = \beta_f 2\pi CR$$

with  $\beta_f$  the density of nodes in this heuristic

## FIRST ORDER APPROXIMATION AS AN ASYMPTOTIC

### ■ Theorem

When  $\rho$  tends to infinity:

- The fluid heuristic is asymptotically tight:

$$\beta_o \rightarrow \beta_f, \mathbf{W}_o \rightarrow \mathbf{W}_f, \mu_o \rightarrow \mu_f \dots$$

- The law of the latency of a typical node converges weakly to an exponential random variable of parameter  $\mathbf{W}_f = \frac{\mathbf{F}}{\mu_f}$

### ■ Proof: fluid limit techniques extended to spatial processes

FIRST ORDER APPROXIMATION AS AN ASYMPTOTIC (continued)

■ In this heuristic/limit

$$\beta_f = \sqrt{\frac{\lambda F}{2\pi CR}},$$

$$\mu_f = \sqrt{\lambda F 2\pi CR},$$

$$W_f = \sqrt{\frac{F}{\lambda 2\pi CR}},$$

$$N_f = \sqrt{\frac{\pi}{2}} \sqrt{\frac{\lambda FR^3}{C}} = \sqrt{\frac{\pi}{2}} \sqrt{\rho}$$

■ **Proof:**  $W_f = F/\mu_f$  and  $\beta_f = \lambda W_f$  (Little's law) and  $\mu_f = \beta_f 2\pi CR$   
Hence

$$\beta_f \mu_f = \lambda F \quad \Leftrightarrow \quad \beta_f \beta_f 2\pi CR = \lambda F$$

## COMMENTS

- $\rho$  is large when

- either the arrival intensity, or the file size, or the range are large
- or if the download speed constant  $C$  is small

- The time scale of a peer is  $W_f = \sqrt{F/(\lambda 2\pi CR)}$

If two nodes are at a distance  $r_0$  such that

$$\frac{F}{C r_0} \ll W_f = \sqrt{\frac{F}{\lambda 2\pi CR}} \Leftrightarrow r_0 \ll \sqrt{\frac{C}{2\pi \lambda FR}} = \frac{R}{\sqrt{2\pi\rho}}$$

then there is little chance to see these two nodes in the steady state:  
**hard exclusion** below that scale.

- $r_0$  tends to 0 in configurations where  $\rho$  tends to infinity and  $R$  is fixed

## FIRST ORDER APPROXIMATION AS A BOUND

- In the TCP case, the repulsion theorem is equivalent to saying that

$$\beta_o 2\pi CR \geq \mu_o$$

- It follows from the relations  $W_o \geq F/\mu_o$  and  $\beta_o = \lambda W_o$  that

$$\beta_o \geq \lambda \frac{F}{\beta_o 2\pi CR}$$

- **Corollary**

$$\beta_o \geq \sqrt{\frac{\lambda F}{2\pi CR}} = \beta_f \quad \text{and} \quad W_o \geq W_f$$

## HARD CORE REGIME

- A stationary point process is **hard-core** for balls of radius  $R$  if there are no other points in a ball of radius  $R$  centered on any point
- **Conjecture** When  $\rho$  tends to 0,
  - the stationary node point process tends to a hard-core point process for balls of radius  $R$  with intensity  $\beta_h$  and latency  $W_h$ :

$$\beta_h = \frac{1}{\pi R^2}, \quad W_h = \frac{1}{\lambda \pi R^2}$$

- the cdf of the latency converges weakly to

$$1 - \frac{e^{-\frac{t}{2W_h}}}{2}, \quad t > 0$$

HARD CORE REGIME (continued)

### Rationale

$$N_f \ll 1$$



$$\sqrt{\frac{\lambda F R^3}{C}} \ll 1$$



$$\sqrt{\frac{\lambda R C F^2 R^2}{F C^2}} \ll 1$$



$$\frac{R F}{C} \ll \sqrt{\frac{F}{2\pi \lambda R C}} = W_f \leq W_o$$

**The latency of two nodes within range is negligible w.r.t. the mean latency**

## SECOND ORDER APPROXIMATION

### ■ Second order heuristic:

– considers  $\hat{\mu}$ , the unique solution of

$$\hat{\mu}^2 = \mu_f^2 \left( 1 - \frac{C}{\hat{\mu}R} \ln \left( 1 + \frac{\hat{\mu}R}{C} \right) \right)$$

– then defines

$$\hat{\beta} = \lambda F / \hat{\mu}, \quad \hat{W}_h = F / \hat{\mu}$$



SECOND ORDER APPROXIMATION (continued)

- Factorization of the factorial moment measure of order 3
- Balance equation for the second order factorial moment density, which reads

$$2\beta_o\lambda = 2m_{[2]}(\mathbf{x}, \mathbf{y}) \frac{C}{F} \frac{\mathbf{1}_{\|\mathbf{x}-\mathbf{y}\|\leq R}}{\|\mathbf{x}-\mathbf{y}\|} + \frac{C}{F} \int_D m_{[3]}(\mathbf{x}, \mathbf{y}, \mathbf{z}) \left( \frac{\mathbf{1}_{\|\mathbf{x}-\mathbf{z}\|\leq R}}{\|\mathbf{x}-\mathbf{z}\|} + \frac{\mathbf{1}_{\|\mathbf{y}-\mathbf{z}\|\leq R}}{\|\mathbf{y}-\mathbf{z}\|} \right) d\mathbf{z}$$

for all  $\mathbf{x}$  and  $\mathbf{y}$ .

- Approximations:

$$m_{[3]}(\mathbf{x}, \mathbf{y}, \mathbf{z}) \approx \frac{m_{[2]}(\mathbf{x}, \mathbf{y})m_{[2]}(\mathbf{x}, \mathbf{z})}{\beta_o}$$

$$m_{[3]}(\mathbf{x}, \mathbf{y}, \mathbf{z}) \approx \frac{m_{[2]}(\mathbf{x}, \mathbf{y})m_{[2]}(\mathbf{y}, \mathbf{z})}{\beta_o}$$

**SECOND ORDER APPROXIMATION** *(continued)*

■ Then

$$\begin{aligned} \beta_o \lambda \approx & \mathbf{m}_{[2]}(\mathbf{x}, \mathbf{y}) \frac{\mathbf{C} \mathbf{1}_{\|\mathbf{x}-\mathbf{y}\| \leq \mathbf{R}}}{\mathbf{F} \|\mathbf{x} - \mathbf{y}\|} \\ & + \mathbf{m}_{[2]}(\mathbf{x}, \mathbf{y}) \frac{\mathbf{C} \mathbf{1}}{\mathbf{F} \mathbf{2}} \int_{\mathbf{D}} \frac{\mathbf{1}_{\|\mathbf{x}-\mathbf{z}\| \leq \mathbf{R}} \mathbf{m}_{[2]}(\mathbf{x}, \mathbf{z})}{\|\mathbf{x} - \mathbf{z}\| \beta_o} d\mathbf{z} \\ & + \mathbf{m}_{[2]}(\mathbf{x}, \mathbf{y}) \frac{\mathbf{C} \mathbf{1}}{\mathbf{F} \mathbf{2}} \int_{\mathbf{D}} \frac{\mathbf{1}_{\|\mathbf{y}-\mathbf{z}\| \leq \mathbf{R}} \mathbf{m}_{[2]}(\mathbf{y}, \mathbf{z})}{\|\mathbf{y} - \mathbf{z}\| \beta_o} d\mathbf{z} \end{aligned}$$

that is

$$\mathbf{m}_{[2]}(\mathbf{x}, \mathbf{y}) \approx \lambda \mathbf{F} \frac{\beta_o}{\frac{\mathbf{C} \mathbf{1}_{\|\mathbf{x}-\mathbf{y}\| \leq \mathbf{R}}}{\|\mathbf{x}-\mathbf{y}\|} + \mu_o}$$

with  $\mu_o =: \mathbf{C} \int_{\mathbf{B}(0, \mathbf{R})} \frac{\mathbf{m}_{[2]}(0, \mathbf{z})}{\beta_o} \frac{1}{\|\mathbf{z}\|} d\mathbf{z}$

SECOND ORDER APPROXIMATION (continued)

So

$$\begin{aligned}\mu_o &\approx \lambda F 2\pi C \int_0^R \frac{1}{\mu_o + \frac{C}{r}} dr \\ &= \lambda F 2\pi C \left( \frac{R}{\mu_o} - \frac{C}{\mu_o^2} \ln\left(1 + \frac{\mu_o R}{C}\right) \right)\end{aligned}$$

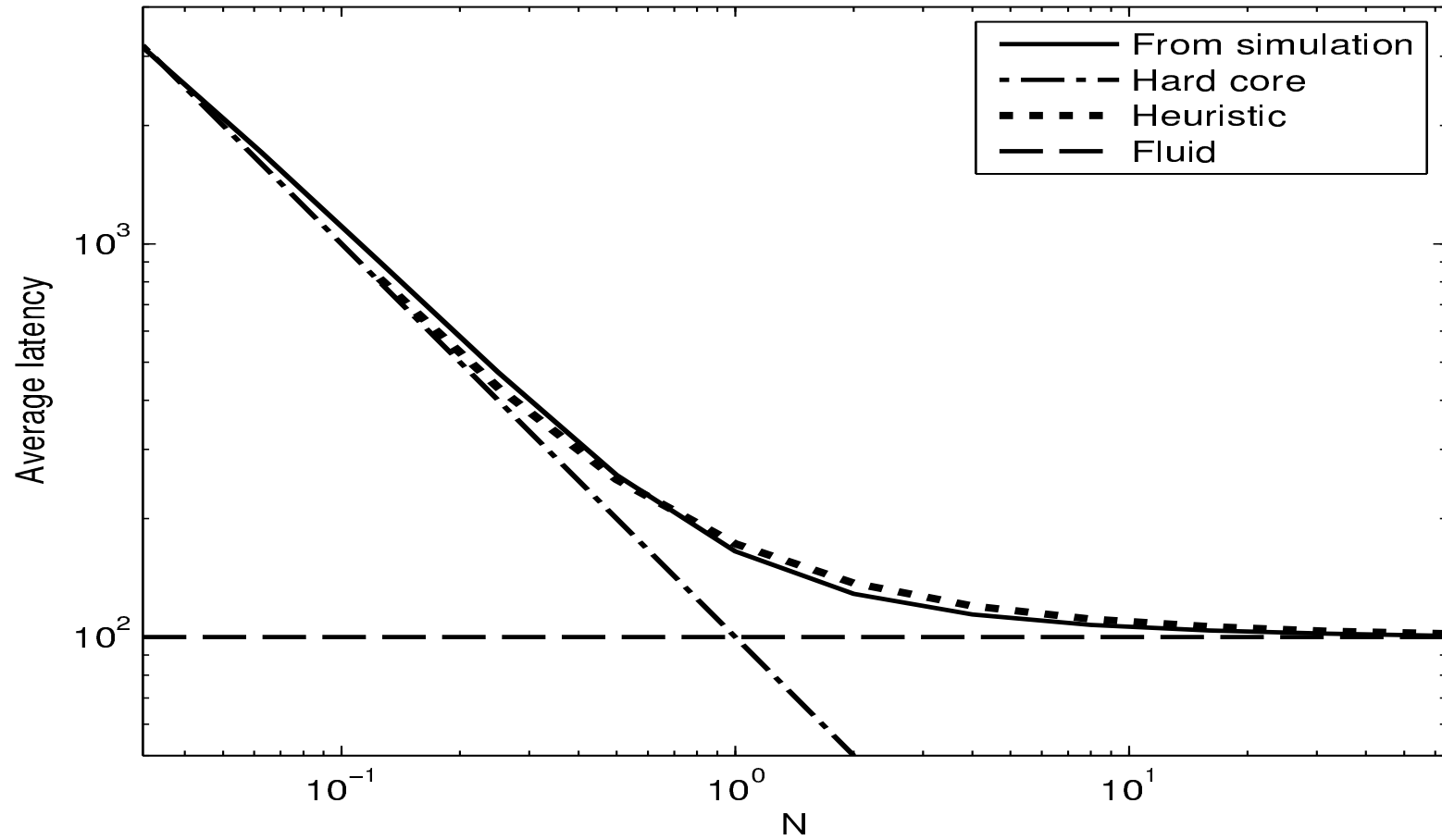
and

$$\hat{\mu}^2 = \mu_f^2 \left( 1 - \frac{C}{\hat{\mu}R} \ln \left( 1 + \frac{\hat{\mu}R}{C} \right) \right)$$

### 3. SIMULATION

- Fix 3 independent parameters and use the 4-rth one to run through all possible scenarios.
- The two first fixed parameters are  $R = .1$  and  $C = 1$
- Set  $W_f$  to 100. This implies that for all simulations, the fluid model will predict the same mean latency.
- Then, we use  $N_f$  as the variable parameter: We use  $N_f$  instead of  $\rho$  as main dimensionless parameter
- The remaining input parameters of the system are then completely defined:

$$\lambda = \frac{N_f}{\pi R^2 W_f}, \quad F = \frac{2N_f C W_f}{R}$$

Latency in function of  $N_f$ .

## 4. SUPER-SCALABILITY

- Dimensional analysis tells us that

$$\begin{aligned} W_o(\lambda, F, C, R) &= M \left( \sqrt{\frac{\pi \lambda F R^3}{2C}} \right) W_f(\lambda, F, C, R) \\ &= M \left( \sqrt{\frac{\pi \lambda F R^3}{2C}} \right) \sqrt{\frac{F}{\lambda 2\pi C R}} \end{aligned}$$

where  $M$  only depends on  $N_f = \sqrt{\frac{\pi \lambda F R^3}{2C}}$  and is **decreasing**.

- $\lambda$  and  $R$  are both **win-win** parameters. As they increase, both terms in the RHS decrease and the mean latency hence tends towards 0, while the behavior of the system becomes more and more fluid.

## SCALABILITY & SUPER SCALABILITY

**Single Server**  
**M/M/1 Queue**  
**Does not scale**

$$W = \frac{1}{\mu - \lambda}, \lambda < \mu$$

**Infinite Server**  
**M/M/ $\infty$  Queue**  
**Scales**

$$W = \frac{1}{\mu}$$

**Network Limited P2P**  
**Spatial B & D P2P**  
**Super Scales**

$$W = \frac{m(\lambda)}{\sqrt{\lambda}}, m(\cdot) \downarrow$$

## ADAPTING THE PEERING RADIUS

- **Mean Constant Number of Nearest Nodes:** take as neighbors the nodes in a ball with a radius  $R$  such that the mean number of other nodes in the ball is  $L$  i.e.  $\pi R^2 \beta_o = L$ , where  $\beta_o$  is the (unknown) steady state intensity of the point process  $\phi_t$ . Then

$$f(\mathbf{r}) = \frac{C}{r} \mathbf{1}_{r \leq R}, \quad R = \sqrt{\frac{L}{\pi \beta_o}}$$

- **General Case**

$$f(\mathbf{r}) = \frac{C}{r} \mathbf{1}_{r \leq R}, \quad R = \kappa \beta_o^{-\alpha}$$

- **(DA)** All system properties only depend on the parameter

$$\rho = \frac{\lambda F}{C} \kappa^{\frac{3}{1-2\alpha}}$$



## ASYMPTOTIC BEHAVIOR

- **General  $\alpha$  case:  $R = \kappa\beta^{-\alpha}$**
- **think of all parameters fixed and let  $\lambda$  tend to infinity**
  - $\beta$  is of the order  $\lambda^b$  with  $b = \frac{1}{2-\alpha}$  the **density exponent**
  - $W$  is of the order  $\lambda^l$  with  $l = \frac{\alpha-1}{2-\alpha}$  the **latency exponent**
  - $R$  is of the order  $\lambda^r$  with  $r = \frac{\alpha}{\alpha-2}$  the **radius exponent**
  - $N$  is of the order  $\lambda^n$  with  $n = \frac{1-2\alpha}{2-\alpha}$  the **swarm exponent**
- **2 regimes, both compatible with fluid:**
  - For  $\alpha > 2$ , we get a node density and a latency which both tend to 0 when  $\lambda$  tends to  $\infty$ : **Heaven's-flash**
  - For  $\alpha < 1$ , we get a density and swarm that tend to infinity and a latency which tends to zero when  $\lambda$  tends to  $\infty$ : **Swarm-flash**

## EXTENSIONS & OPEN QUESTIONS

- **Extensions (P2P version to be presented at IEEE Infocom 13)**
  - Seeders
  - Wireless
- **Challenges (Applied Probability Journal Paper in preparation)**
  - Hard core regime
  - Higher order approximations based on moment measures
  - Chunks

## CONCLUSION

- A **new point process dynamics** for handling P2P and CDN
- Has a unique stationary regime which exhibits a form of **repulsion**
- Satisfies a hierarchy of birth and death type integral relations which lead to **good approximations**