Projection of a Smooth Space Curve:

Numeric and Certified Topology Computation

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Contraction Contr

informatics mathematics

Projection and apparent contour

3D curve = intersection of 2 implicit surfaces

$$\mathcal{C}: \begin{cases} p(x, y, z) = 0\\ q(x, y, z) = 0 \end{cases}, (x, y, z) \in \mathbb{R}^3 \\ \mathcal{B} = \pi_{(x, y)}(\mathcal{C}) \end{cases}$$



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 $\mathcal{B} = \{(x, y) \in \mathbb{R}^2 | f(x, y) = 0\}$ Singularities:

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• Restrict to box B_0



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- Compute local topology at singularities
- Connect boxes : Graph

Symbolic tool box

- Gröbner basis
- Triangular decomposition
- Rational univariate representation

• ...

- + Handle any types of singularities
- + Bit complexity analysis
- Global analysis
- High complexity
- Restricted to polynomial functions
- Generic case is typically the worst case

Numerical tool box

- Subdivision
- Homotopy
- + Local: analysis restricted to a box
- + Adaptative: running time sensitive to the local geometry
- + Fast limited precision computation
- + Certification via interval analysis
- + Not restricted to polynomials: only evaluation required
- Difficult to analyse the complexity
- Need generic assumptions: regular solutions

- Symbolic: f polynomial, squarefree part = $f/\gcd(f, f')$
- Numeric:
 - Newton iteration: $x_{n+1} = x_n f(x_n)/f'(x_n)$
 - Homotopy
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Subdivide interval I_0 until

- $0 \notin \Box f(I) \longrightarrow$ no solution in I
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Use instead $g=f+\lambda f'$, $\lambda\in\mathbb{R}$

- works only for solutions of multiplicity 2 $(g' = f' + \lambda f'' \neq 0)$
- adds spurious solutions

Interval analysis

Arithmetic operations $\begin{bmatrix} a, b \end{bmatrix} \oplus \begin{bmatrix} c, d \end{bmatrix} = \begin{bmatrix} a + c, b + d \end{bmatrix}$ $\begin{bmatrix} a, b \end{bmatrix} \otimes \begin{bmatrix} c, d \end{bmatrix} = \begin{bmatrix} \min(ac, ad, bc, bd), \max(ac, ad, bc, bd) \end{bmatrix}$

Interval function extension of $f:\mathbb{R}\longrightarrow\mathbb{R}$

- $\Box f(I) \supseteq \{f(x) | x \in I\}$
- convergence: $w(\Box f(I)) \longrightarrow 0$ as $w(I = [a, b]) = b a \longrightarrow 0$

Examples

- \bullet f polynomial, use interval arithmetic operations
- Mean value evaluation $\Box f(I) := f(mid(I)) + \Box f'(I)(I mid(I))$

Exclusion criterion:

- $0 \not\in \Box f(I) \Longrightarrow f$ has no solution in I
- ... But $0 \in \Box f(I)$ does NOT imply that f has no solution in I

Newton/Krawczyk operator

 $F: \mathbb{R}^n \longrightarrow \mathbb{R}^n$ F(X) = 0: system of n equations, n unknowns. Assume the solutions are regular: the determinant of the jacobian J_F does not vanish at the solutions.

Interval Newton operator: $m \in X \subset \mathbb{R}^n$, $N(X) = m - J_F^{-1}(m) \Box F(X)$ Krawczyk operator = mean value evaluation of N

$$K(X) = N(m) + \Box J_N(X)(X - m)$$

Lemma.

- $K(X) \subset X \Longrightarrow \exists !$ solution in X
- $\bullet \ K(X) \ \cap X = \emptyset \Longrightarrow \text{ no solution in } X$
- Quadratic convergence
- Neumaier, Interval methods for systems of equations, 1990
- Dedieu, Points fixes, zeros et la methode de Newton, 2006
- Rump, Verification methods: Rigorous results using floating point arithmetic, Acta Numerica, 2010

Subdivision algorithm

Input: $F: \mathbb{R}^n \to \mathbb{R}^n$, X_0 box of \mathbb{R}^n **Output:** A list R of boxes containing solutions in X_0 of F = 0 $L := \{X_0\}$ **Repeat:** X := L.popIf $0 \in F(X)$ then If $K_F(X) \subset Int(X)$ then insert X in RElse If $K_F(X) \cap X \neq \emptyset$ then bisect X and insert its sub-boxes in LEnd if End if Until $L = \emptyset$ **Return** R



WARNING:

- Square system: as many equations as unknowns
- Regular solutions: $det(J_F(s)) \neq 0$

Our problem: Isolate singularities





Projection of a 3D smooth curve: Generic singularities are Nodes = transverse intersection of 2 branches

Apparent contour: Generic singularities are Nodes and Cusps

z

Subresultant approach in 2D

 $\mathcal{B} = \{(x,y) \in \mathbb{R}^2 | r(x,y) = 0\}$, where $r(x,y) = Res(p,p_z,z)(x,y)$

Singularities of $\ensuremath{\mathcal{B}}$ are the solutions of

$$r(x,y) = \frac{\partial r}{\partial x}(x,y) = \frac{\partial r}{\partial y}(x,y) = 0$$

- Over-determined
- Cusps are solutions of multiplicity 2

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Let s_{10}, s_{11}, s_{22} be the coefficients of the subresultant sequence of p and p_z wrt z

$$(\mathcal{S}_2)$$
 $s_{10}(x,y) = s_{11}(x,y) = 0$ and $s_{22}(x,y) \neq 0$

Lemma.

- Square system
- Nodes and cusps are regular solutions

[IMP15] R. Imbach, G. Moroz, and M. Pouget. Numeric certified algorithm for the topology of resultant and discriminant curves. Research Report RR-8653, Inria, April 2015.

4D approach

 $\begin{aligned} \mathcal{C} &= \{ (x, y, z) \in \mathbb{R}^3 | p(x, y, z) = p_z(x, y, z) = 0 \} \\ \mathcal{B} &= \{ (x, y) \in \mathbb{R}^2 | \exists z \in \mathbb{R} \text{ s.t. } (x, y, z) \in \mathcal{C} \} \end{aligned}$



• Node: $(x, y, z_1), (x, y, z_2) \in \mathcal{C}$, with $z_1 \neq z_2$

• Cusp: $(x, y, z_1), (x, y, z_2) \in \mathcal{C}$, with $z_1 = z_2$

Set $z_1 = c - \sqrt{r}, z_2 = c + \sqrt{r}$, with c center and \sqrt{r} radius of $[z_1 z_2]$

4D approach

$$\mathcal{C} = \{(x, y, z) \in \mathbb{R}^3 | p(x, y, z) = p_z(x, y, z) = 0\}$$
$$\mathcal{B} = \{(x, y) \in \mathbb{R}^2 | \exists z \in \mathbb{R} \text{ s.t. } (x, y, z) \in \mathcal{C}\}$$

$$(S_4) \begin{cases} \frac{1}{2}(p(x,y,c+\sqrt{r}) + p(x,y,c-\sqrt{r})) = 0\\ \frac{1}{2\sqrt{r}}(p(x,y,c+\sqrt{r}) - p(x,y,c-\sqrt{r})) = 0\\ \frac{1}{2}(p_z(x,y,c+\sqrt{r}) + p_z(x,y,c-\sqrt{r})) = 0\\ \frac{1}{2\sqrt{r}}(p_z(x,y,c+\sqrt{r}) - p_z(x,y,c-\sqrt{r})) = 0 \end{cases}$$

Lemma. S_4 is regular, its solutions project to cusps and nodes of \mathcal{B} R. Imbach, G. Moroz, M. Pouget. Numeric and Certified Isolation of the Singularities of the Projection of a Smooth Space Curve, MACIS 2015

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Certified numerical tracking in 3D

- Enclose $\mathcal{C} \subset \cup C_k$, with $C_k = (x_k, y_k, z_k)$ 3D box
 - More efficient than a classic 3D subdivision
 - Correct topology
 - Certification via a parametric interval Krawczyk test

[MGGJ13] B. Martin, A. Goldsztejn, L. Granvilliers, and C. Jermann. Certified parallelotope continuation for one-manifolds. SIAM Journal on Numerical Analysis, 2013.



Certified numerical tracking in 3D

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- Enclose $\pi_{(x,y)}(\mathcal{C}) = \mathcal{B} \subset \cup \pi_{(x,y)}(C_k)$



Certified numerical tracking in 3D

- Enclose $\mathcal{C} \subset \cup C_k$, with $C_k = (x_k, y_k, z_k)$ 3D box
- Enclose $\pi_{(x,y)}(\mathcal{C}) = \mathcal{B} \subset \cup \pi_{(x,y)}(C_k)$
- ullet Restrict the 4D solving domain of \mathcal{S}_4
 - Cusp in $\mathcal{B}_k \longleftrightarrow$ sol. in $(x_k, y_k, z_k, [0(\frac{w(z_k)}{2}])^2])$
 - Node in $\mathcal{B}_{ij} = \mathcal{B}_i \cap \mathcal{B}_j \neq \emptyset \longleftrightarrow$ sol. in $(x_{ij}, y_{ij}, \frac{z_i + z_j}{2}, [0, (\frac{z_i z_j}{2})^2])$



Experiments: isolation of singularities

Degree	RSCube	\mathcal{S}_2 (Sub-resultant)	\mathcal{S}_4 in	\mathcal{S}_4 with
of $p(x,y,z)$	\mathbb{R}^2	in $[-1, 1]^2$	$[-1,1]^2 imes \mathbb{R} imes \mathbb{R}^+$	curve tracking
5	3.1	0.05	24.8	1.25
6	32	0.50	8.40	2.36
7	254	4.44	43.8	4.13
8	1898	37.9	70.2	5.91
9	9346	23.1	45.6	5.30

Average running times in seconds for 5 random dense polynomials of degree d, bitsize 8

- Symbolic method becomes intractable (RSCube via triangular decomposition by F. Rouillier)
- Subdivision: working in 4D is more expensive than in 2D
- Subdivision: tracking the curve is efficient