Reading group	
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Introduction	Reading group: Mathematical Control Theory by
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Discrete time	Ludurdo Dr Contug
Exercices	Alexandre Vieira
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Plenty application in engineering: automobile, microprocessors, robotics, airplanes, drones...



We assume small variations of  $\theta$ , and we end up with equation :

$$\ddot{\theta} - \theta = u$$

How to bring, from any point  $\theta(0)$  and  $\dot{\theta}(0)$  small enough, the system to  $(\theta, \dot{\theta}) = (0, 0)$  ?



$$\ddot{ heta}(t) - heta(t) = u(t) = 3e^{-2t}$$

with  $\theta(0) = 1$  and  $\dot{\theta}(0) = -2$ : stabilizes the system to (0,0).

But if we add a little disturbance to the initial conditions:  $\dot{\theta}(0) = -2 + \varepsilon$ , then the system  $\theta(t) \xrightarrow[t \to +\infty]{} +\infty$ .



$$\ddot{\theta} - \theta = u$$

If we choose a control proportional to the angle:  $u = \alpha \theta$ , then we will never approach (0,0). This can be shown for the whole system: left as an exercice! However, there exist positive scalars  $\alpha$  and  $\beta$  such that the control

$$u = -\alpha\theta - \beta\dot{\theta}$$

stabilizes the system at (0,0). Only problem : how to know  $\theta(t)$  and  $\dot{\theta}(t)$  ?



$$\ddot{\theta} - \theta = u$$

Writing  $x = (\theta, \dot{\theta})^{\mathsf{T}}$ , we rewrite the system as:

$$\dot{x} = \underbrace{\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}}_{=A} x + \underbrace{\begin{pmatrix} 0 \\ 1 \end{pmatrix}}_{=B} u$$

We apply a linear feedback u = Kx:

$$\dot{x} = (A + BK)x$$

How to choose K? Such that A + BK has negative eigenvalues!

# Description of a system

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### Definition

- A system or machine consists of:
  - A time set  $\mathcal{T}$ , which is a subgroup of  $(\mathbb{R},+)$
  - $\bullet$  A non empty-set  ${\mathcal X}$  called the state-space
  - ${\scriptstyle \bullet}$  A non empty-set  ${\cal U}$  called a control-value space
  - A map  $\phi : \mathcal{D}_{\phi} \to \mathcal{X}$  called the transition map, which is defined on a subset  $\mathcal{D}_{\phi}$  of

$$\left\{(\tau,\sigma,x,\omega)|\sigma,\tau\in\mathcal{T},\ \sigma\leq\tau,\ x\in\mathcal{X},\omega\in\mathcal{U}^{[\sigma,\tau[}\right\}\right\}$$

## Description of a linear system

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$$\dot{x}(t) = Ax(t) + Bu(t)$$
 for  $t \in \mathbb{R}$ 

For a given  $u(\cdot) \in \mathbb{R}^m$  and  $x(0) \in \mathbb{R}^n$ , we have a unique solution:

$$x(t) = e^{At}x(0) + e^{At}\int_0^t e^{-As}Bu(s)ds$$

• Time set:  $\mathcal{T} = \mathbb{R}$ 

- State space:  $\mathcal{X} = \mathbb{R}^n$
- Control-value space:  $\mathcal{U} = \mathbb{R}^m$
- Transition map:  $\phi(\tau, \sigma, x, \omega) = e^{A(\tau-\sigma)}x + e^{A\tau} \int_{\sigma}^{\tau} e^{-As} B\omega(s) ds$

## Description of a system

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### Definition

A system  $\Sigma(\mathcal{T}, \mathcal{X}, \mathcal{U}, \phi)$  must comply with:

- Nontriviality: for all state  $x \in \mathcal{X}$ , there exist at least one pair  $(\sigma, \tau) \in \mathcal{T}^2$ ,  $\sigma < \tau$  and  $\omega \in \mathcal{U}^{[\sigma, \tau]}$  such that  $\omega$  is admissibile for x, i.e.  $(\tau, \sigma, x, \omega) \in \mathcal{D}_{\phi}$ .
- Restriction: If  $\omega \in \mathcal{U}^{[\sigma,\mu[}$  is admissible for x, then for each  $\tau \in [\sigma,\mu[, \omega_1 = \omega_{\mid [\sigma,\tau[}$  is also admissible for x and  $\omega_{\mid [\tau,\mu[}$  is admissible for  $\phi(\tau,\sigma,x,\omega_1)$ .

# Description of a system

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### Definition

A system  $\Sigma(\mathcal{T}, \mathcal{X}, \mathcal{U}, \phi)$  must comply with:

• Semigroup: If  $\sigma$ ,  $\tau$ ,  $\mu$  are any three elements of  $\mathcal{T}$  so that  $\sigma < \tau < \mu$ , if  $\omega_1 \in \mathcal{U}^{[\sigma,\tau[}$  and  $\omega_2 \in \mathcal{U}^{[\tau,\mu[}$ , and if x is a state so that

$$\phi(\tau, \sigma, x, \omega_1) = x_1$$
 and  $\phi(\mu, \tau, x_1, \omega_2) = x_2$ 

then  $\omega = \omega_1 \omega_2$  is also admissible for x and  $\phi(\mu, \tau, x, \omega) = x_2$ .

• Identity: For each  $\sigma \in \mathcal{T}$  and each  $x \in \mathcal{X}$ , the empty sequence  $\diamond \in \mathcal{U}^{[\sigma,\sigma[}$  is admissible for x and  $\phi(\sigma,\sigma,x,\diamond) = x$ .

## Output

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### Definition

A system with output is given by a system  $\boldsymbol{\Sigma}$  together with:

- $\bullet$  A set  ${\mathcal Y}$  called the measurement-value or output-value space
- A map  $h: \mathcal{T} \times \mathcal{X} \to \mathcal{Y}$  called the readout or measurement map

# Trajectory

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### Definition

A trajectory  $\Gamma$  for the system  $\Sigma$  on the interval  $\mathcal{I} \subseteq \mathcal{T}$  is a pair of function  $(\xi, \omega) \in \mathcal{X}^{\mathcal{I}} \times \mathcal{U}^{\mathcal{I}}$  such that

$$\xi(\tau) = \phi(\tau, \sigma, \xi(\sigma), \omega_{|[\omega, \tau[})$$

holds for each pair  $\sigma, \tau \in \mathcal{I}, \sigma < \tau$ . For a given  $\omega$ , we call such  $\xi$  a path. We call  $\xi(\sigma)$  the initial state,  $\xi(\tau)$  the terminal state.

### Definition

 $\boldsymbol{\Sigma}$  is said complete if every input is admissible for every state:

$$\mathcal{D}_{\phi} = \left\{ ( au, \sigma, x, \omega) | \sigma, au \in \mathcal{T}, \ \sigma \leq au, \ x \in \mathcal{X}, \omega \in \mathcal{U}^{[\sigma, au[} 
ight\}$$

	Discrete systems
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Introduction	A system is said discrete if $\mathcal{T}=\mathbb{Z}.$
Systems Discrete time	We define the transition mapping by
Exercices	$\begin{array}{ll} \mathcal{P}: & \mathcal{E} \to & \mathcal{X} \\ & (t,x,u) \mapsto & \mathcal{P}(t,x,u) \end{array}$ where $\mathcal{E}$ is a subset of $\mathbb{Z} \times \mathcal{X} \times \mathcal{U}$ . For all $t$ in $\mathbb{Z}$ : $x(t+1) = \mathcal{P}(t,x(t),\omega(t))$

## Linear discrete systems

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### Definition

A discrete time system  $\Sigma$  is linear (over field  $\mathbb{K}$ ) if:

- $\Sigma$  is complete
- $\mathcal{P}(t,\cdot,\cdot)$  is linear for each  $t\in\mathbb{Z}$
- If it has output in addition:
  - $\mathcal{Y}$  is a vector space
  - $h(t, \cdot)$  is linear for each  $t \in \mathbb{Z}$ .

The system is finite dimensional if  $\mathcal{X}$  and  $\mathcal{U}$  (and, if it exists,  $\mathcal{Y}$ ) are finite dimensional. In that case, we call dimension of  $\Sigma$  the dimension of  $\mathcal{X}$ .

### Linear discrete systems

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If  $\Sigma$  is linear, there exists linear maps  $A(t) : \mathcal{X} \to \mathcal{X}$ ,  $B(t) : \mathcal{U} \to \mathcal{X}$  and  $C(t) : \mathcal{X} \to \mathcal{Y}$  such that :

$$\mathcal{P}(t, x, u) = A(t)x + B(t)u$$
  
 $h(t, x) = C(t)x$ 

## Smooth discrete systems

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We denote  $\mathcal{E}_t = \{(x, u) | (t, x, u) \in \mathcal{E}\}$  the domain of  $\mathcal{P}(t, \cdot, \cdot)$ .

### Definition

A discrete time system (over field  $\mathbb{K}$ )  $\Sigma$  is  $\mathcal{C}^k$  if, for some intergers n and m:

- $\mathcal{X}$  is an open subset of  $\mathbb{K}^n$
- $\circ \ \mathcal{U}$  is an open subset of  $\mathbb{K}^m$
- For each t ∈ Z, the set E<sub>t</sub> is open and the map P(t, ·, ·) is of class C<sup>k</sup> there.

If it has output in addition:

- $\mathcal Y$  is an open subset of  $\mathbb K^p$
- $h(t, \cdot)$  is of class  $C^k$ .

## Smooth discrete systems

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### Definition

Let  $\Sigma$  be a  $\mathcal{C}^1$  discrete time system over  $\mathbb{R}$ , and assume that  $\Gamma = (\overline{\xi}, \overline{\omega})$  is a trajectory on an interval  $\mathcal{I}$ . The linearization of  $\Sigma$  along  $\Gamma$  is the discrete-time linear system  $\Sigma_*[\Gamma]$  with description  $(\mathbb{R}^n, \mathbb{R}^m, \mathcal{P}_*)$  where :

$$\forall (x, u) \in \mathbb{R}^n \times \mathbb{R}^m, \ \mathcal{P}_*(t, x, u) = A(t)x + B(t)u$$

where  $A(t) = \mathcal{P}_x(t, \overline{\xi}(t), \overline{\omega}(t))$  and  $B(t) = \mathcal{P}_u(t, \overline{\xi}(t), \overline{\omega}(t))$  for all  $t \in \mathcal{I}$  and A(t) = B(t) = 0 for all  $t \notin \mathcal{I}$ . If  $\Sigma$  is a system with outputs, then the discrete-time linear system  $\Sigma_*[\Gamma]$  admits as readout map:

$$h_*(t,x)=C(t)x$$

where  $C(t) = h_x(t, \overline{\xi}(t))$  if  $t \in \mathcal{I}$ , C(t) = 0 otherwise.

### Exercices

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- Exercice 1.2.1, 1.4.2, 2.4.2 seem pretty cool!
- Another one to test if you got it well: start with the transition mapping of a linear discrete-time system

$$x(t+1) = A(t)x(t) + B(t)u(t)$$

and check if it actually defines a system (meaning, make sure that the transition map is:

$$\begin{aligned} x(\tau) &= \phi(\tau, \sigma, x(\sigma), u) \\ &= A^{(\tau-\sigma)}x(\sigma) + \sum_{k=\sigma}^{\tau-1} A^{k-\sigma} Bu(\tau-\sigma-1-k) \end{aligned}$$

and check if it complies with all the axioms in the definition of a system).