

Computing DBF/RBF for Timed/Task Automata

Nan Guan

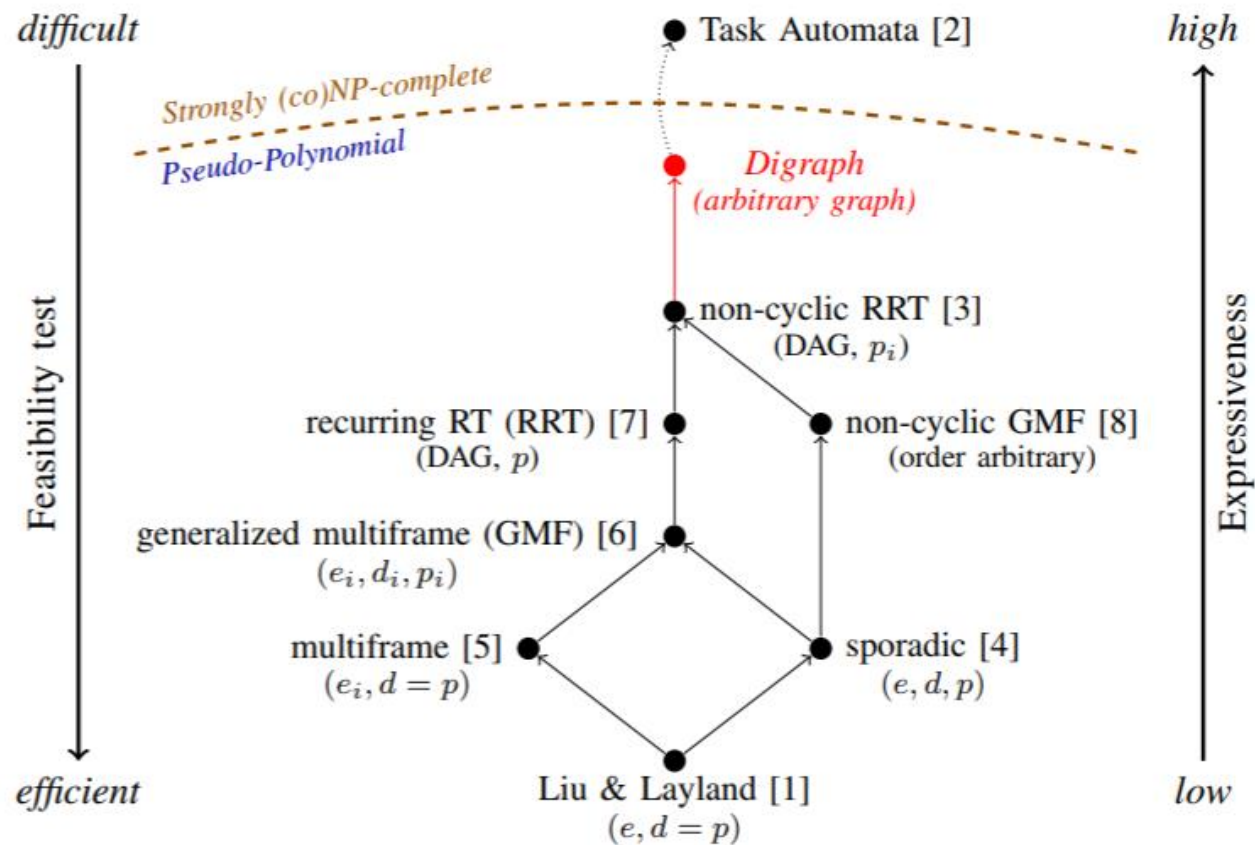
The Hong Kong Polytechnic University

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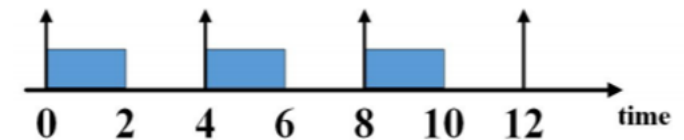
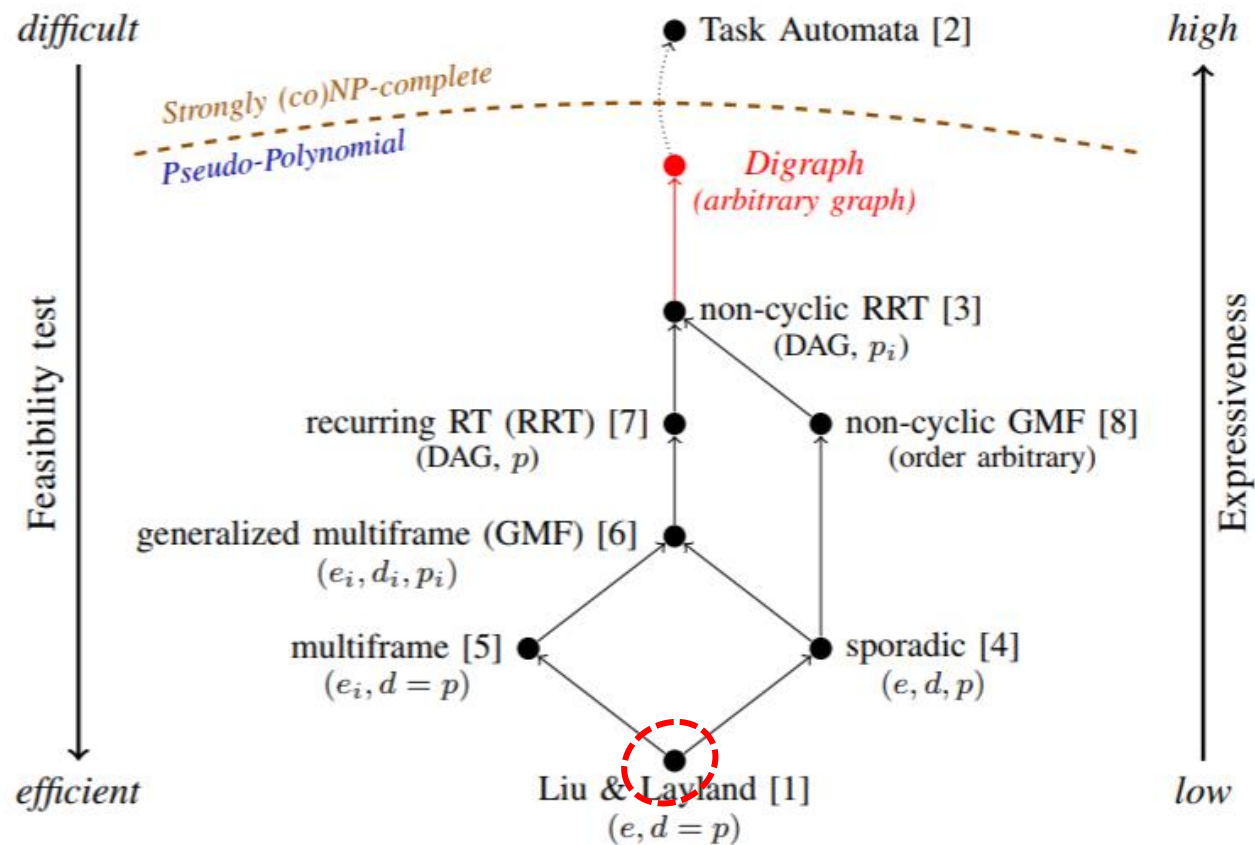
Outline

- Problem
 - What are Timed/Task Automata?
 - Why compute their DBF/RBF?
- A Starting Point

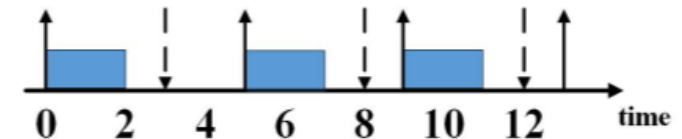
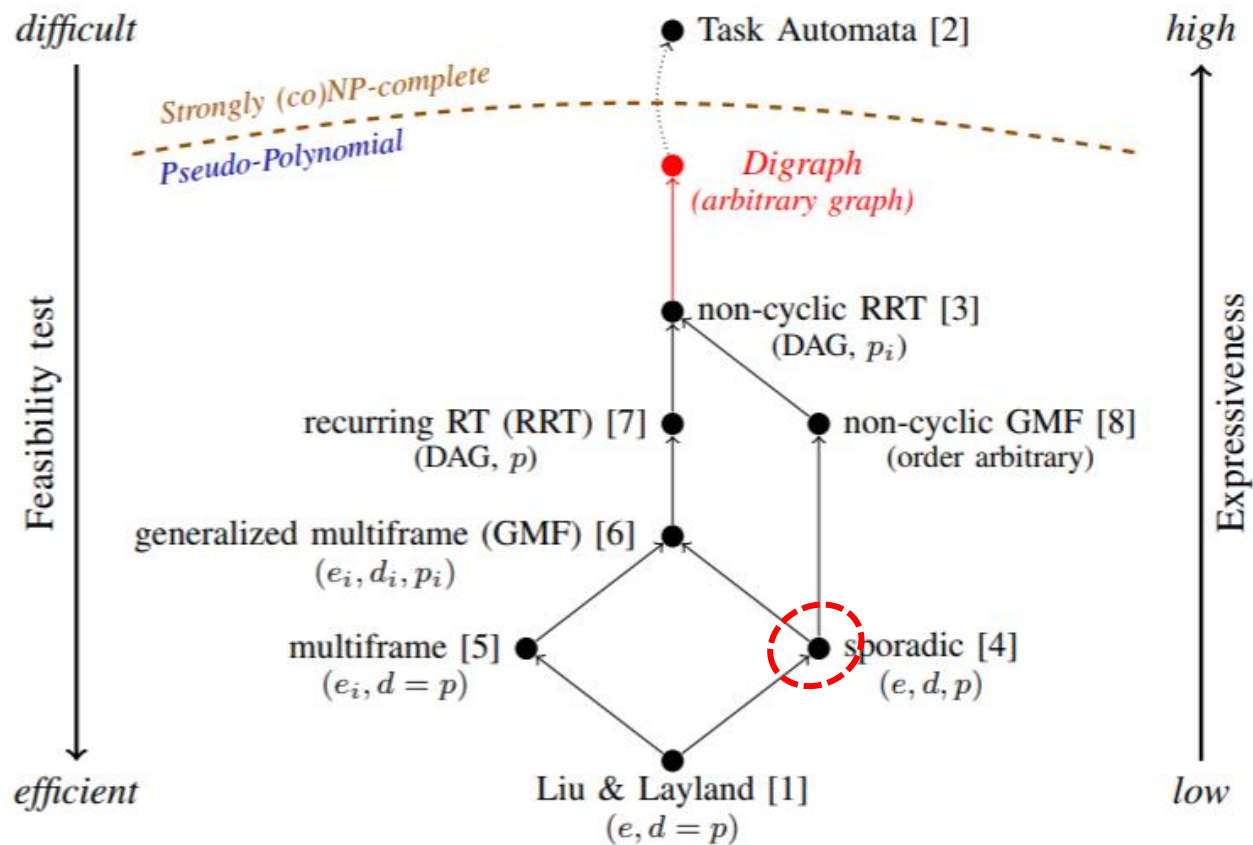
Real-Time Task Models



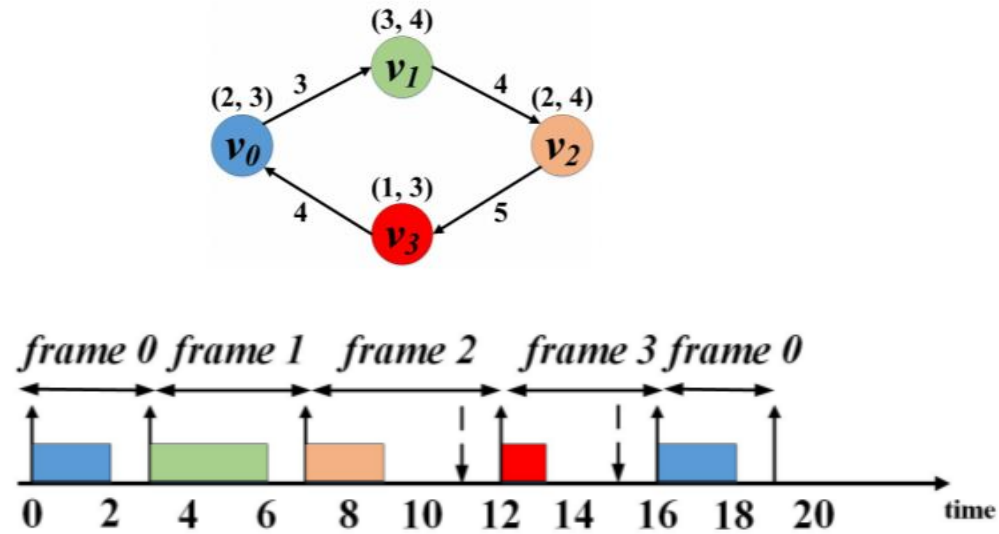
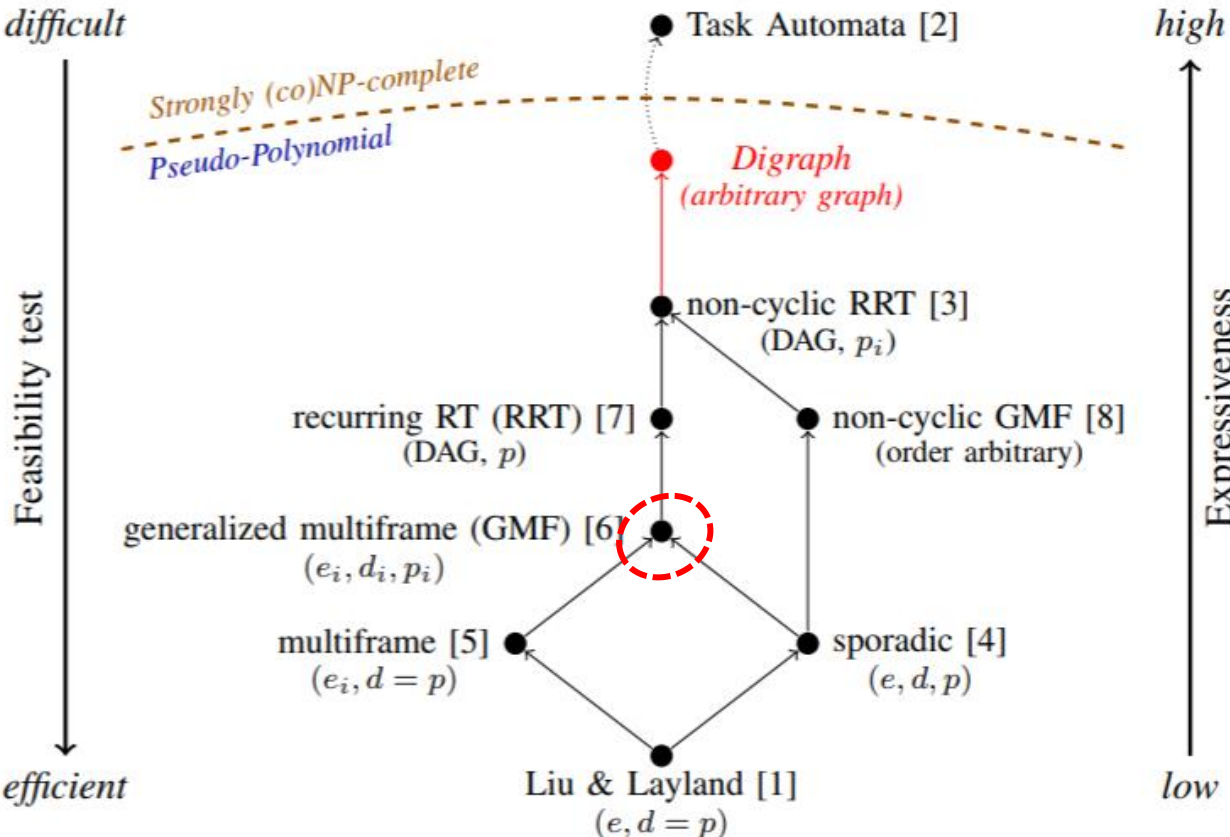
Real-Time Task Models



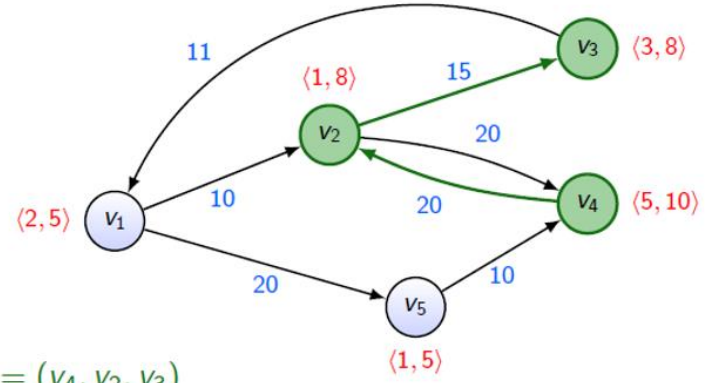
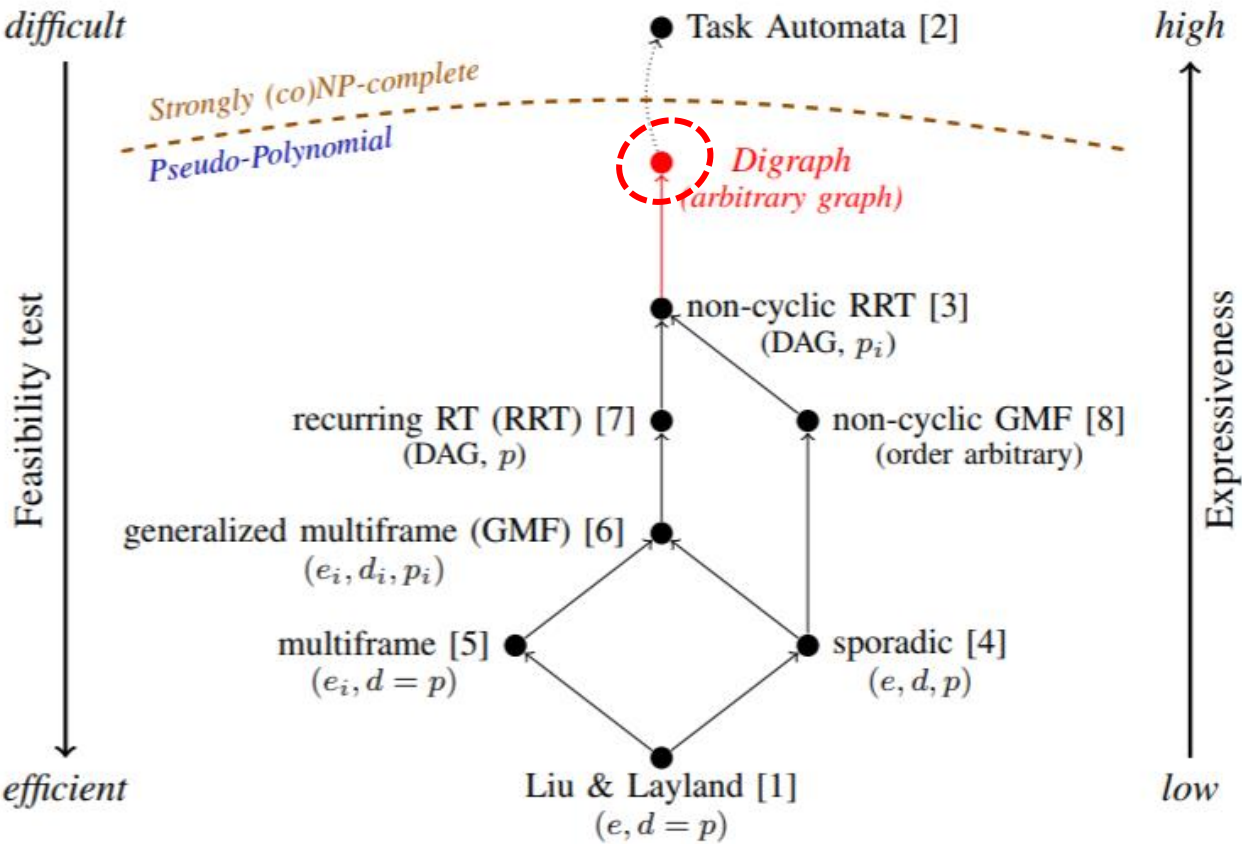
Real-Time Task Models



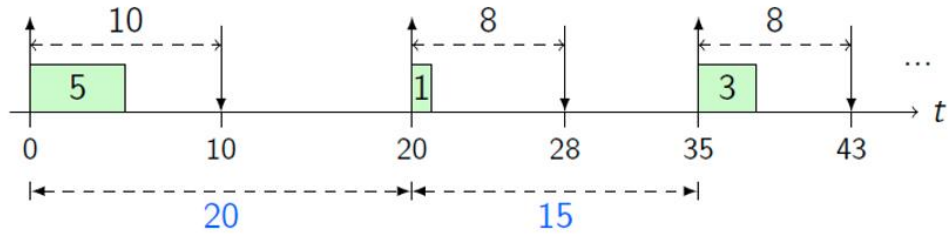
Real-Time Task Models



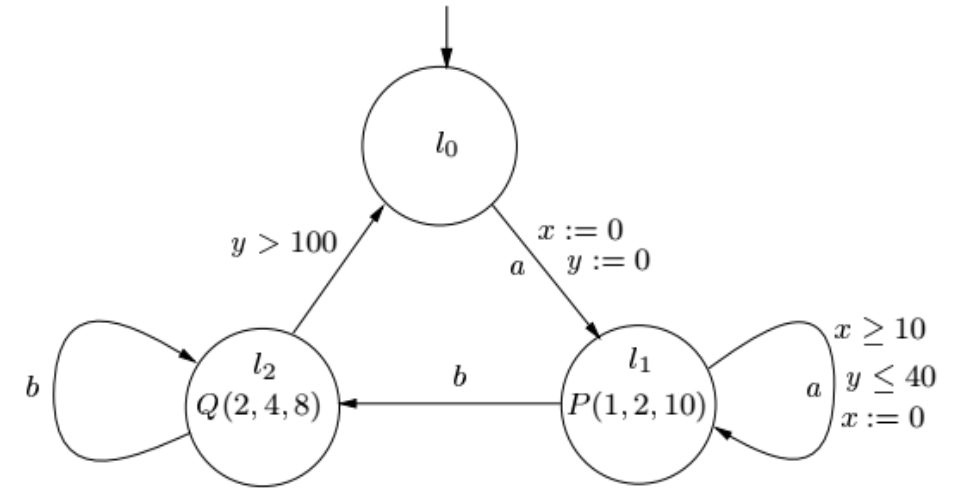
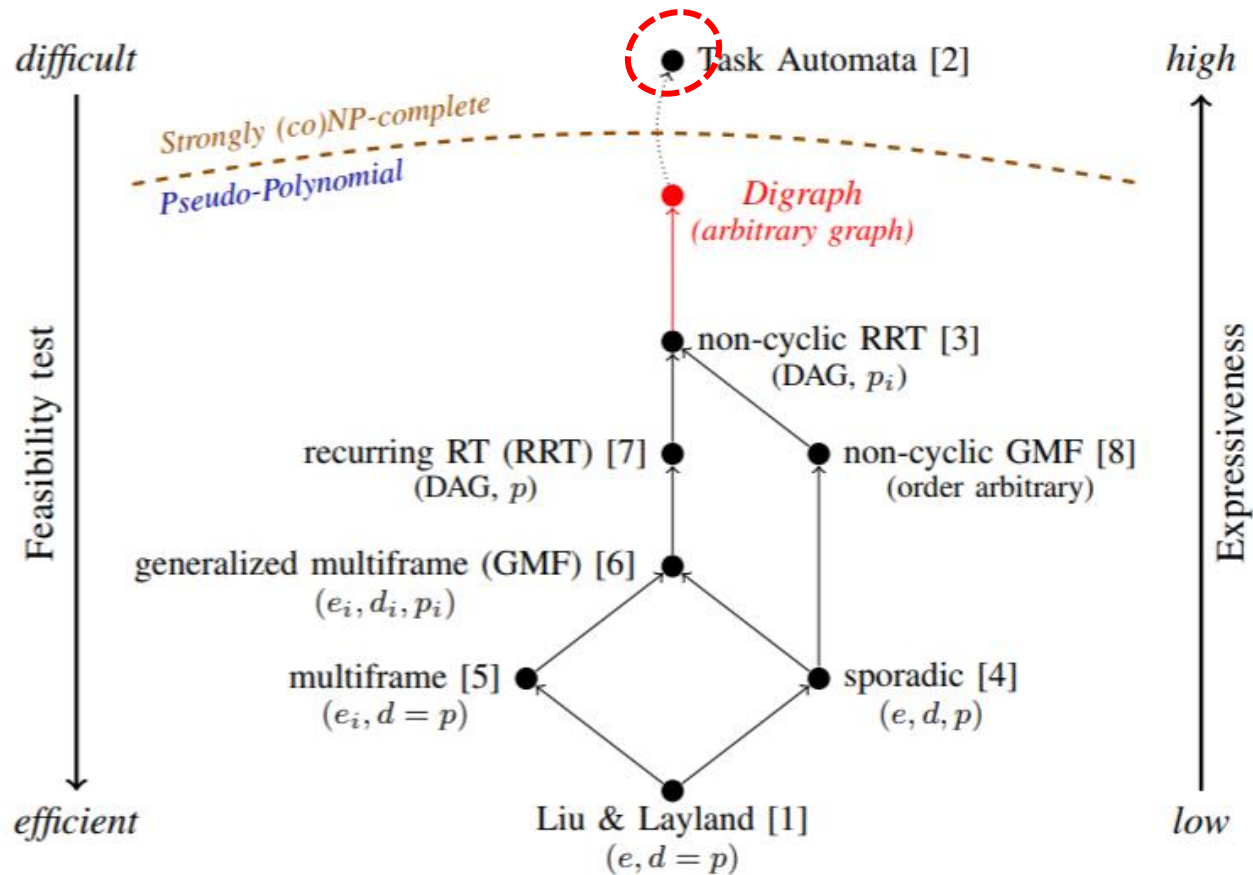
Real-Time Task Models



Path $\pi = (v_4, v_2, v_3)$



Real-Time Task Models

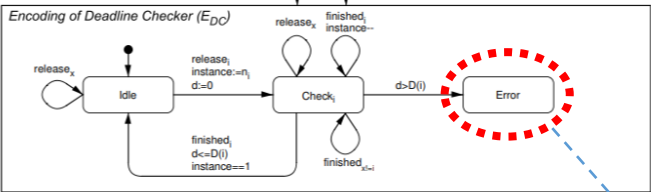
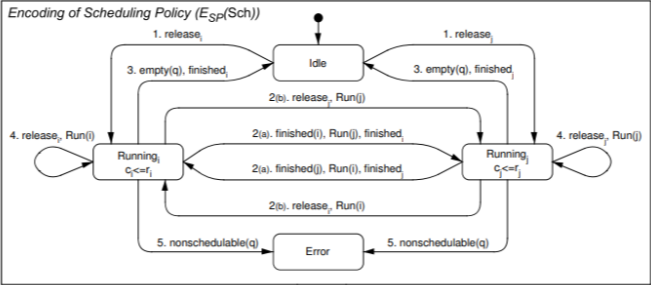


Timed Automata = FSM + clocks

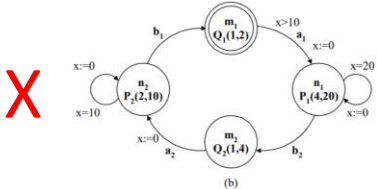
Task Automata = Timed Automata + Task Release

Real-Time Task Models

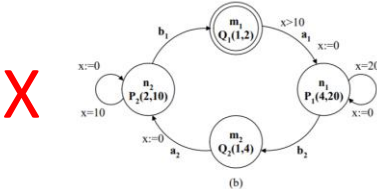
Scheduability as Reachability of TA -- Model Checking



TA for scheduler and deadline checking

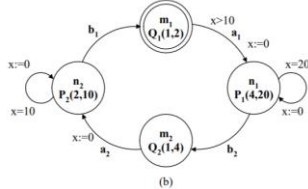


TA for task 1



TA for task 2

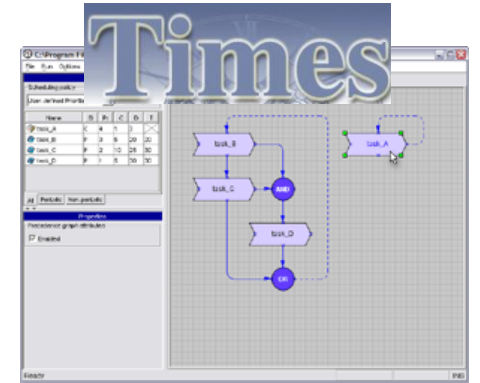
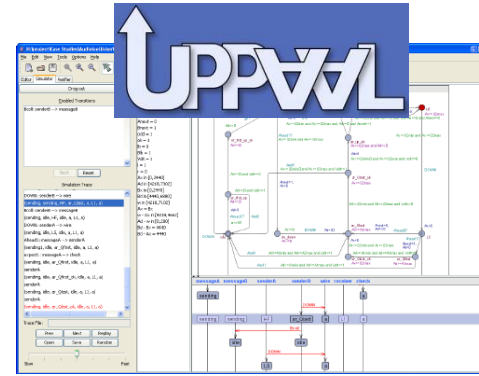
...



TA for task 2

Search the composed state transition graph, to check whether the state representing deadline miss is **reachable**.

Real-Time Task Models



analysis
difficulty

RTSS
ECRTS
RTAS

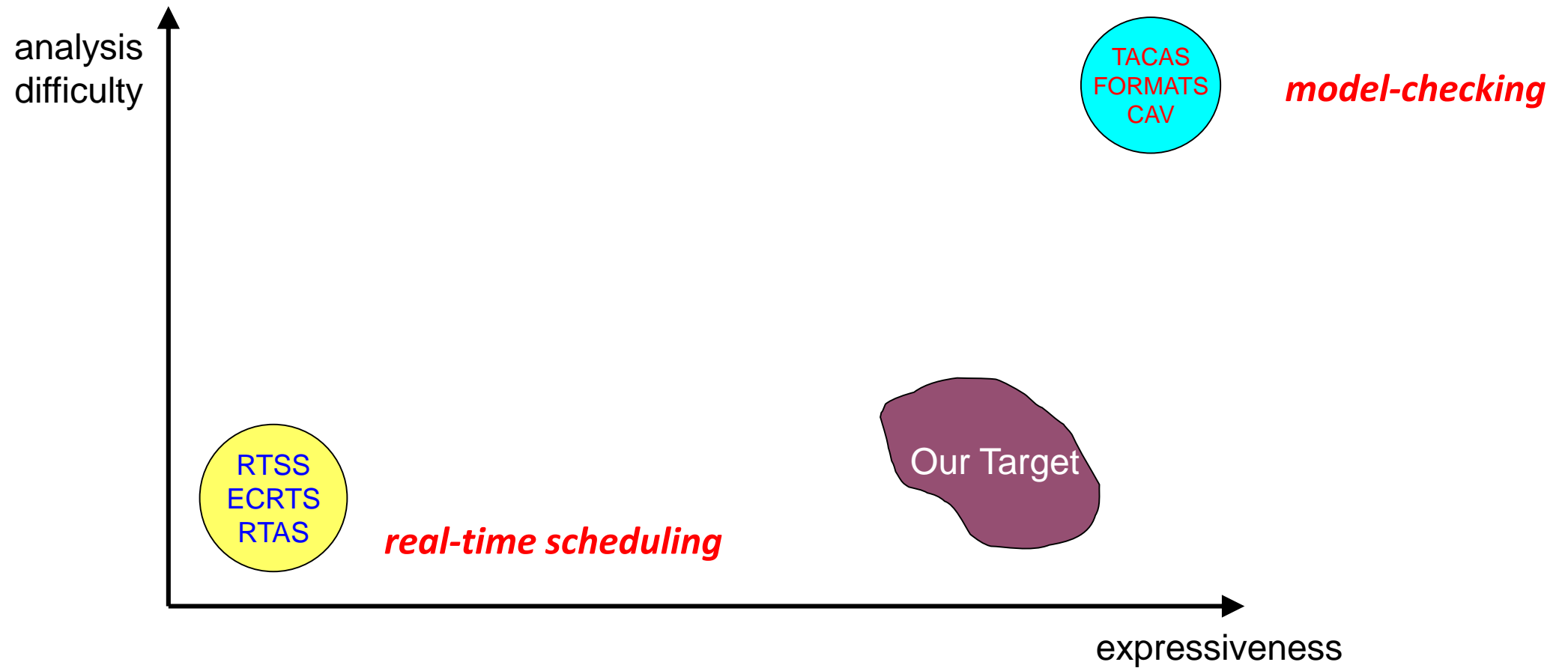
real-time scheduling

TACAS
FORMATS
CAV

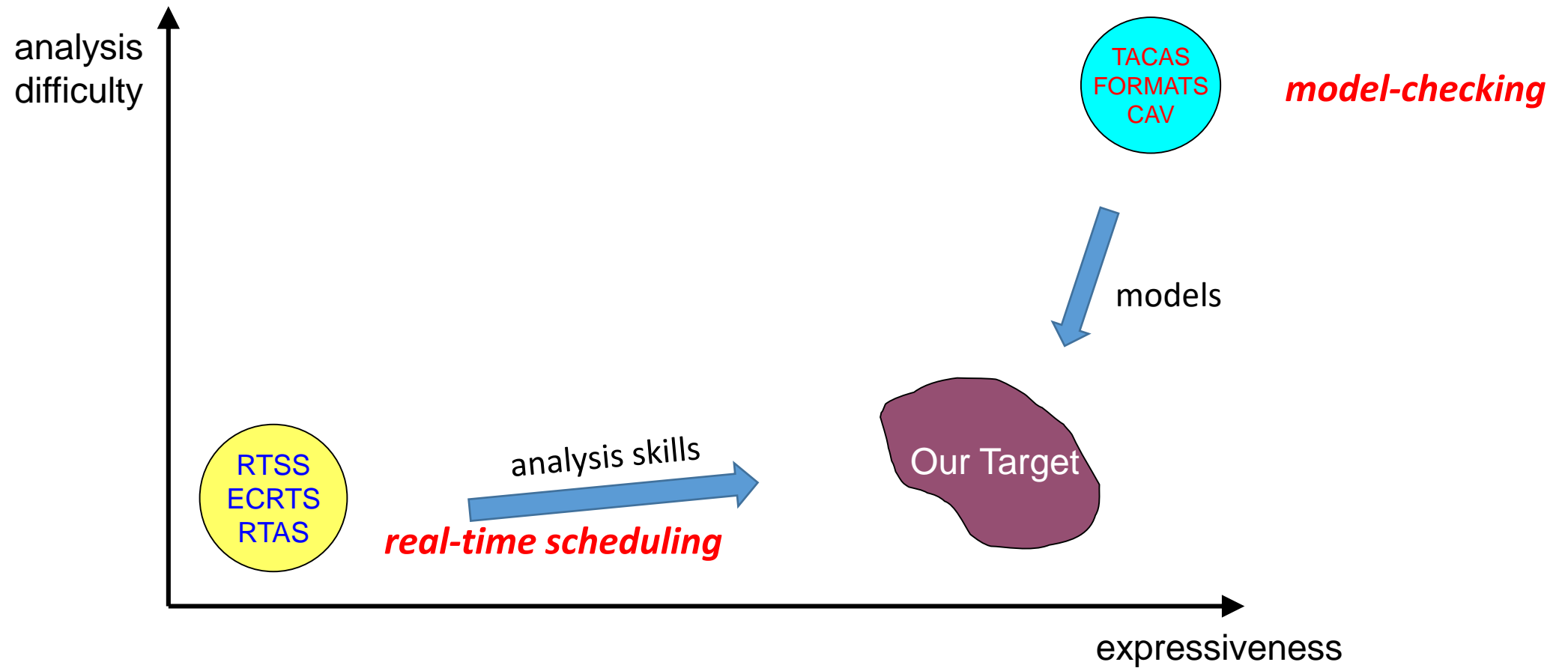
model-checking

expressiveness

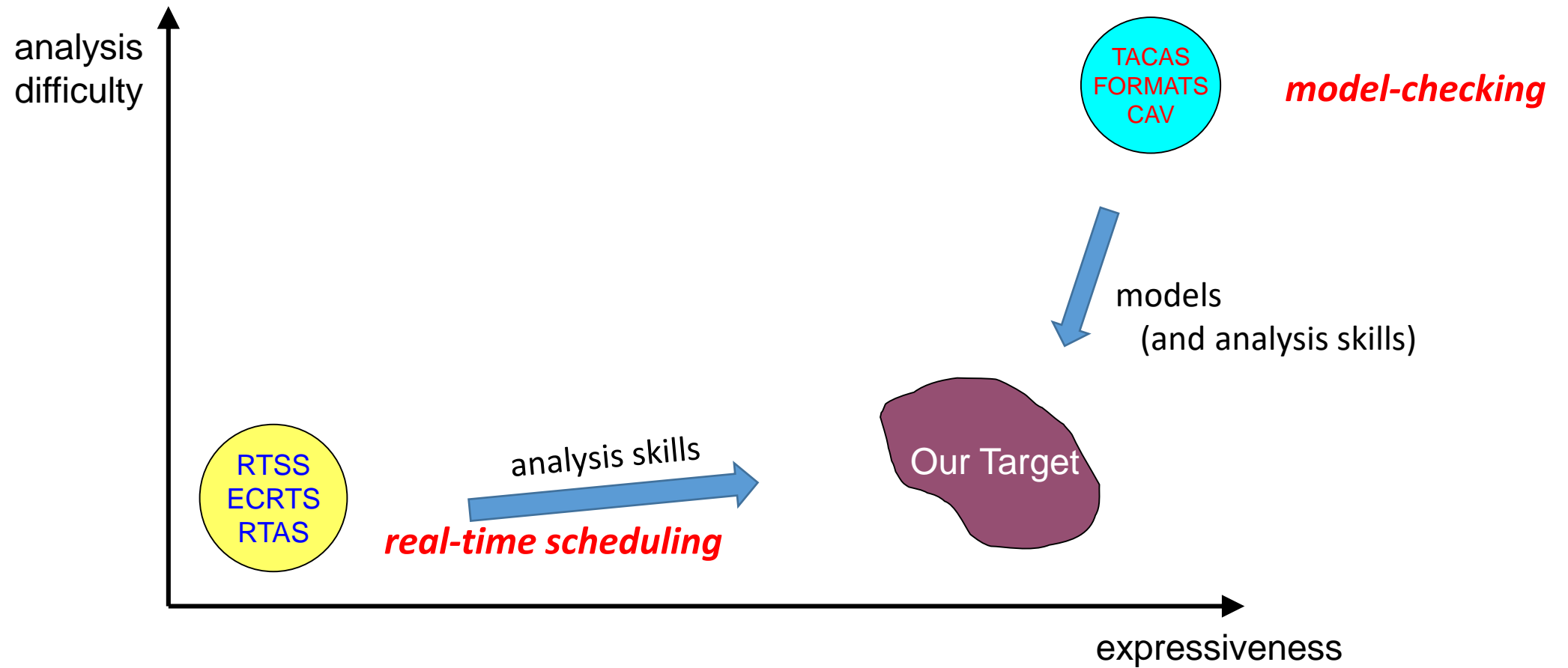
Real-Time Task Models



Real-Time Task Models



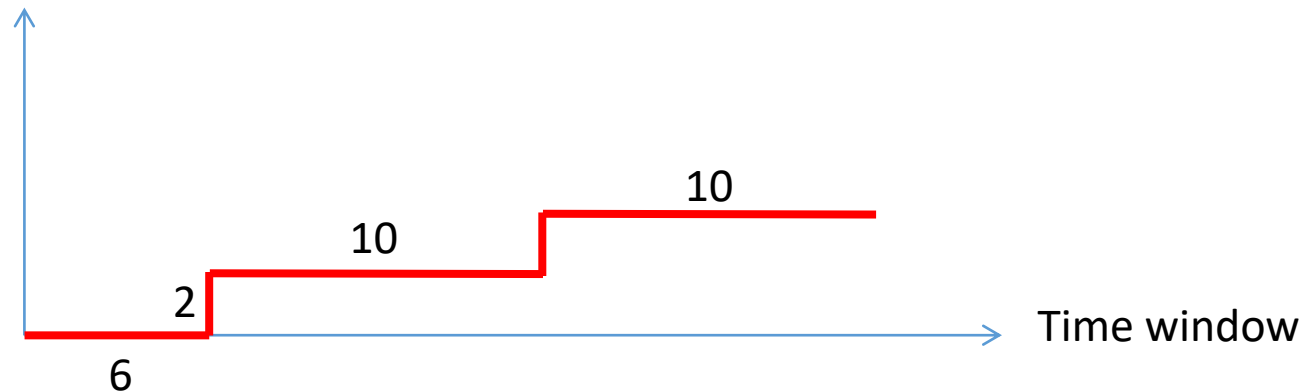
Real-Time Task Models



Demand Bound Function (DBF)

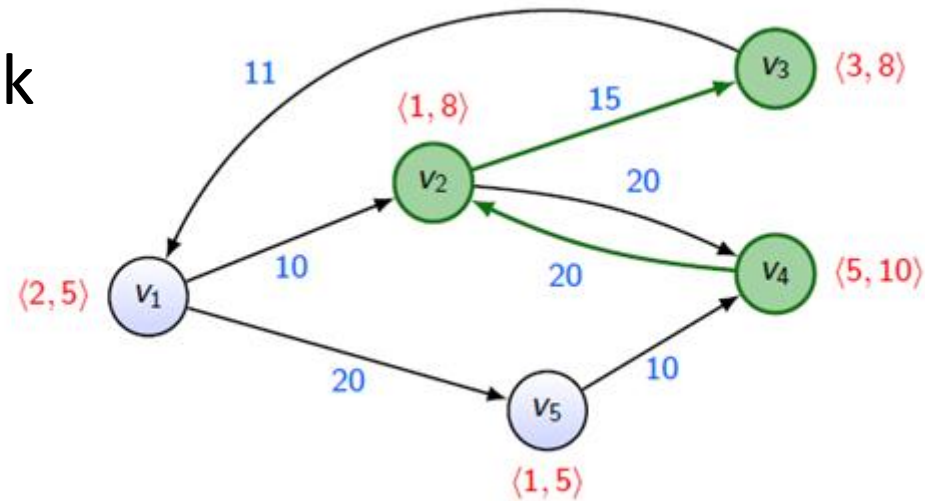
- DBF of a *sporadic* task
 - Period = 10
 - WCET = 2
 - Deadline = 6

Demand Bounds Function (dbf)

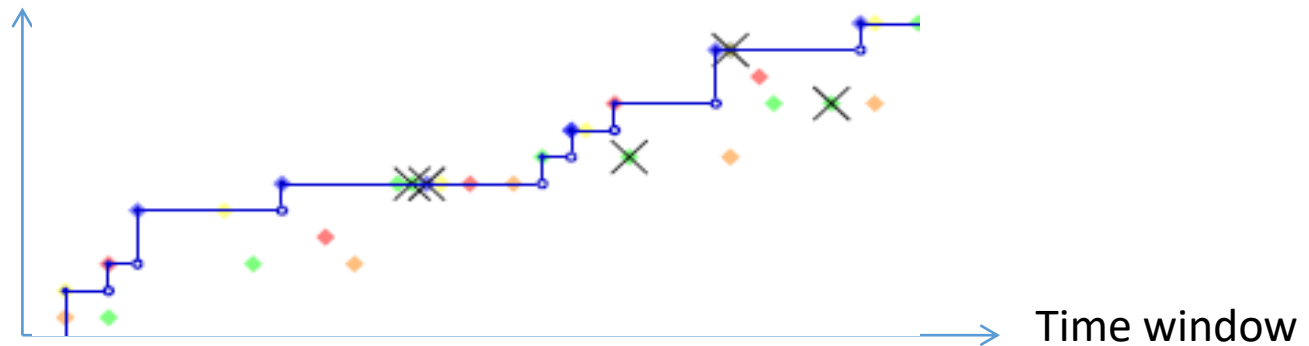


Demand Bound Function (DBF)

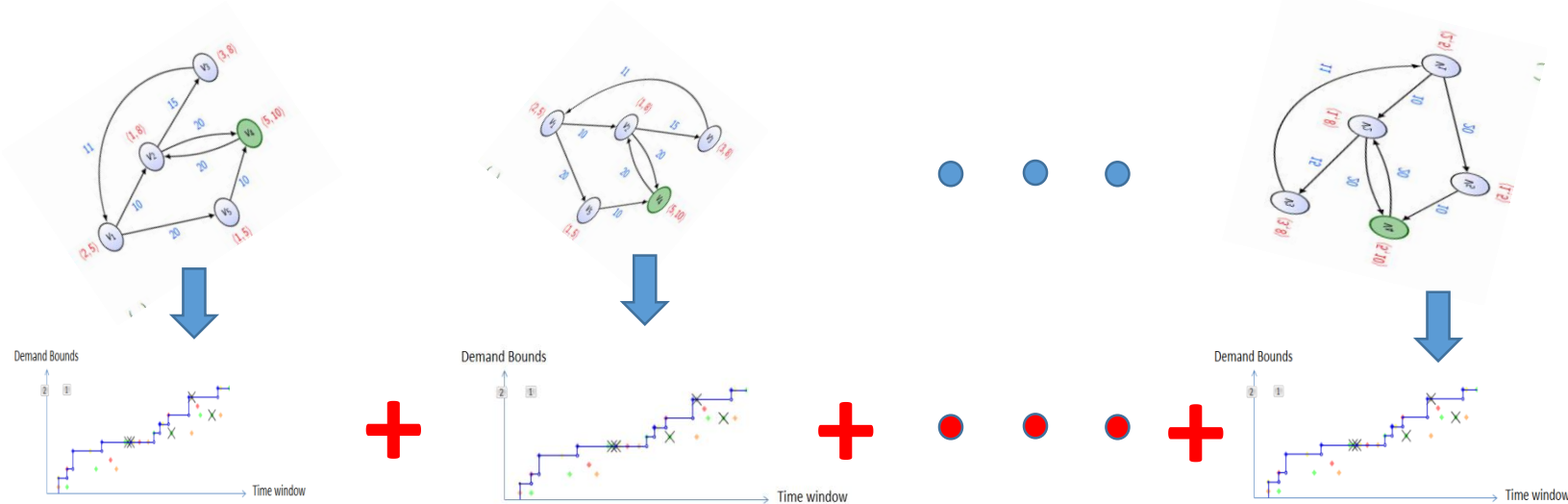
- DBF of a *Digraph* task



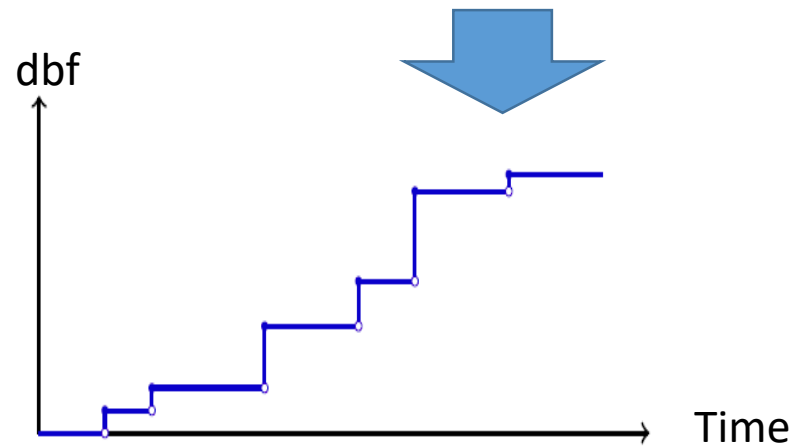
Demand Bounds Function (dbf)



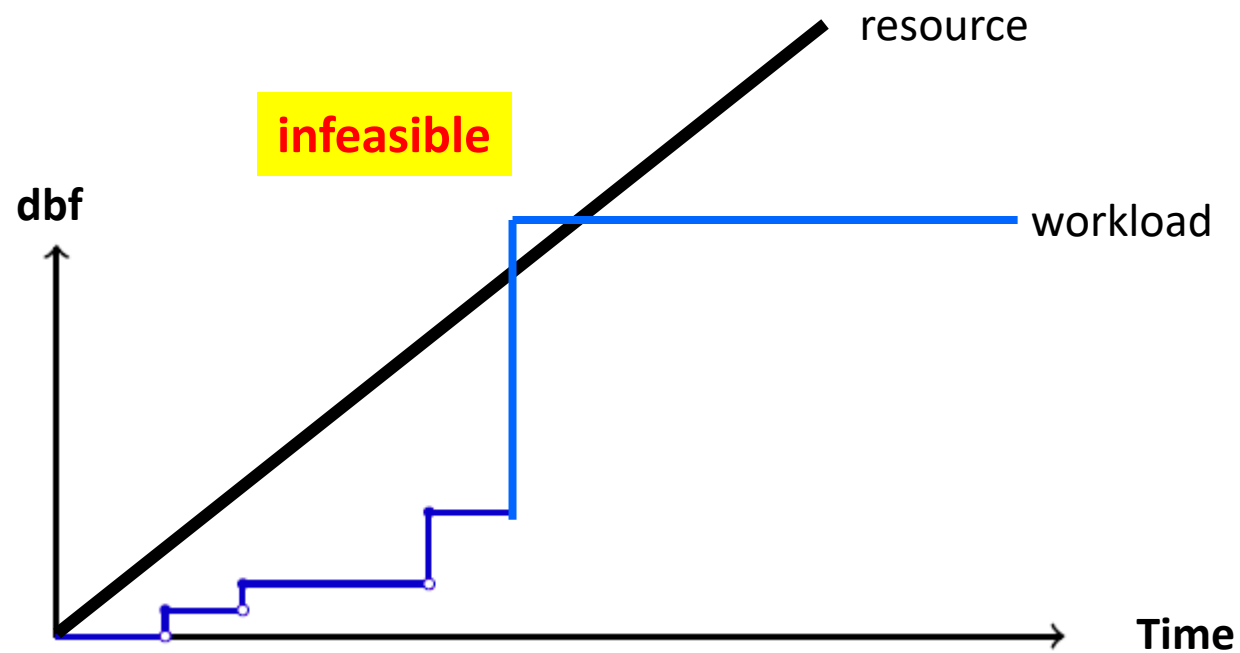
Demand Bound Function (DBF)



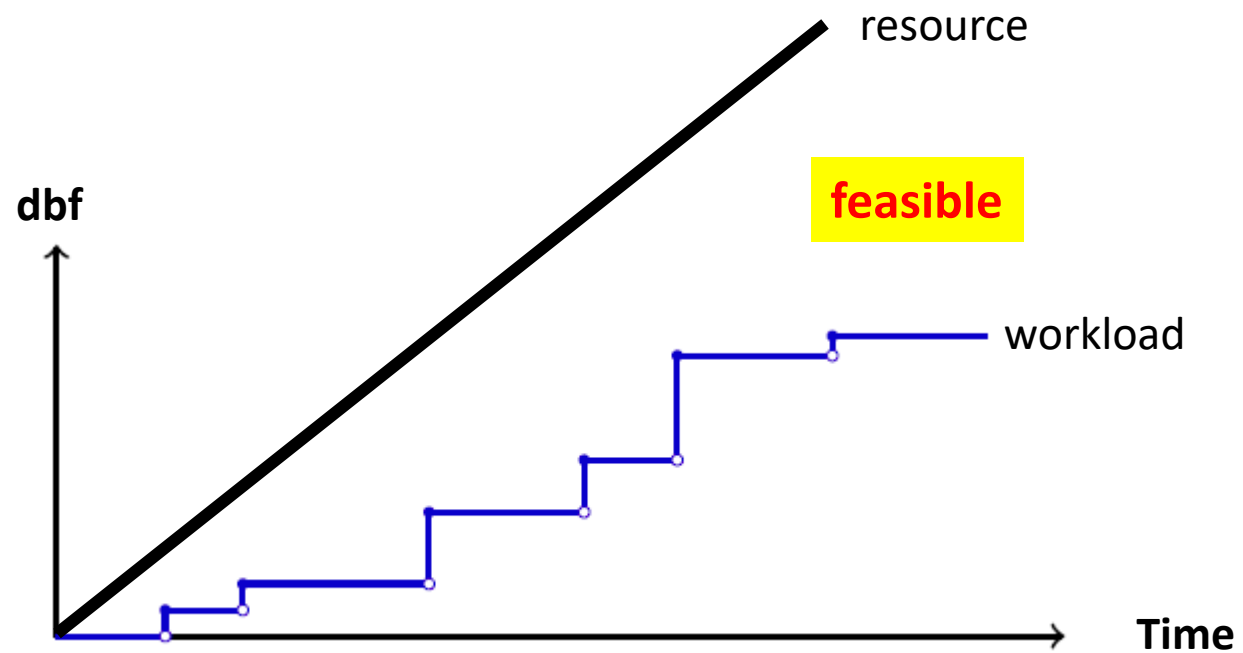
The system workload:



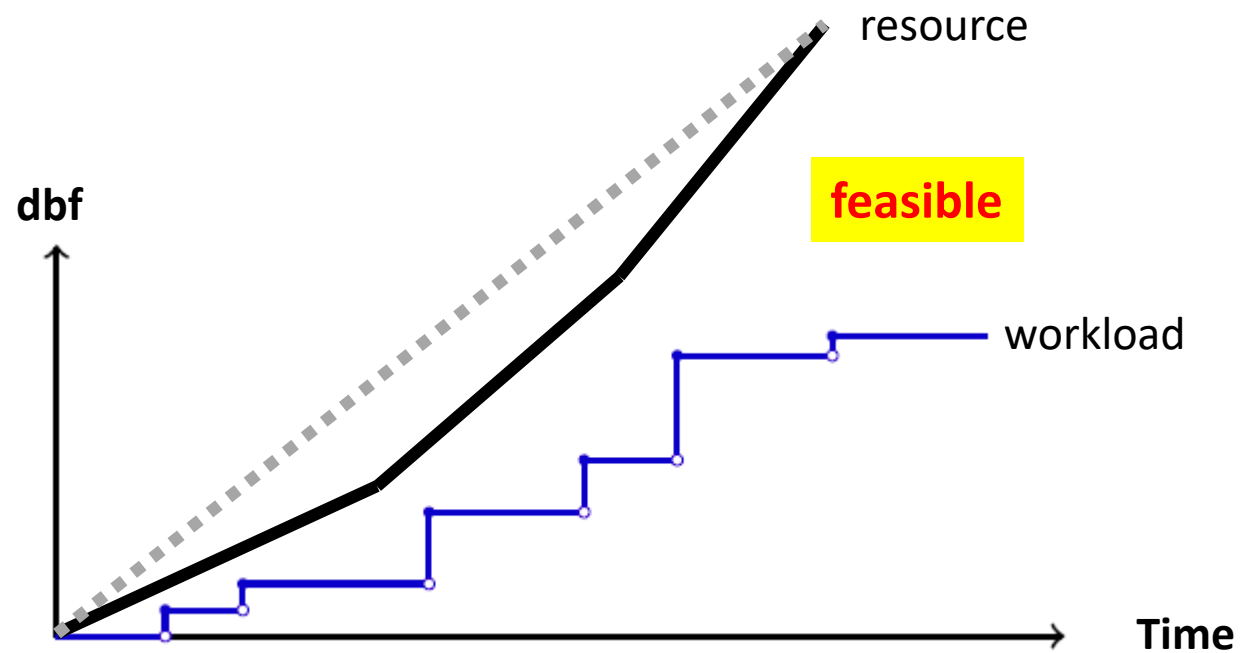
Demand Bound Function (DBF)



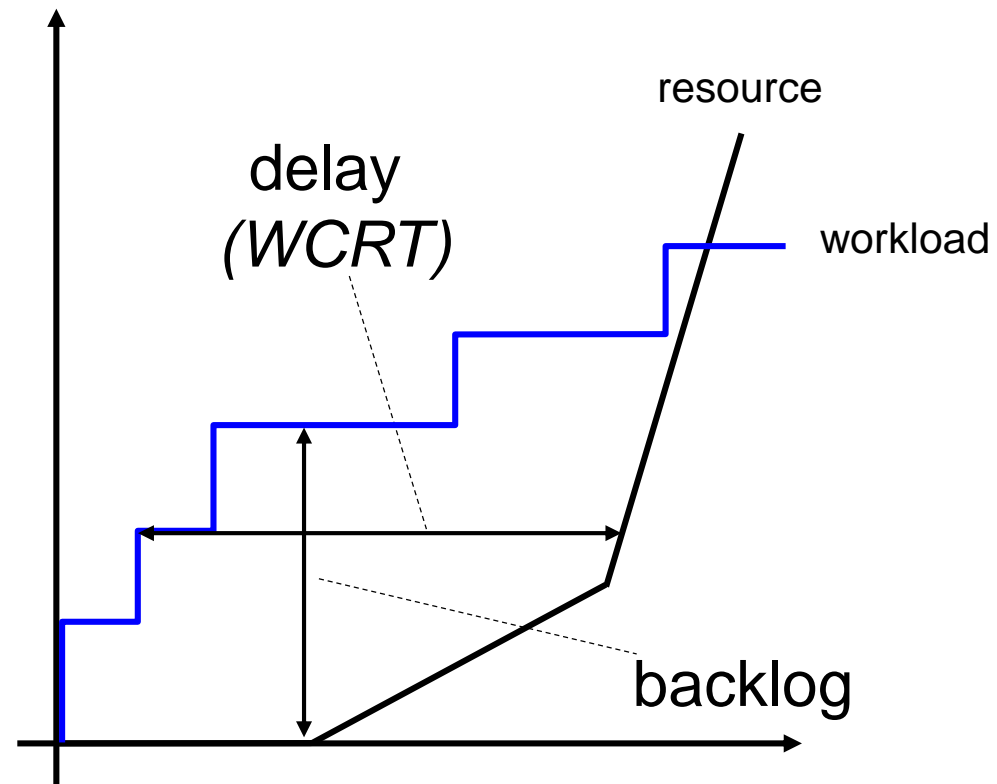
Demand Bound Function (DBF)



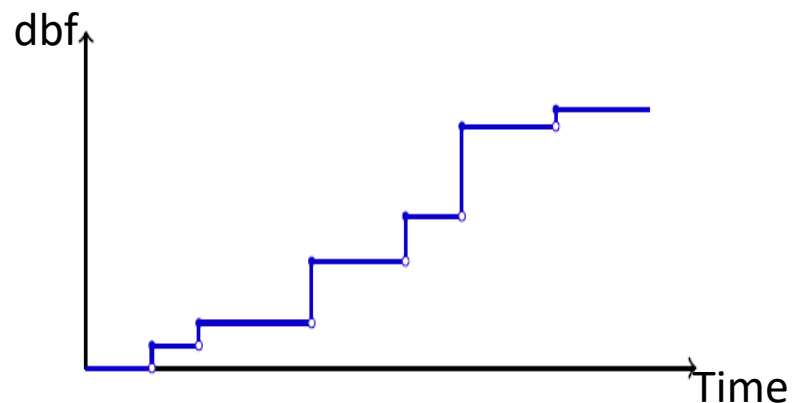
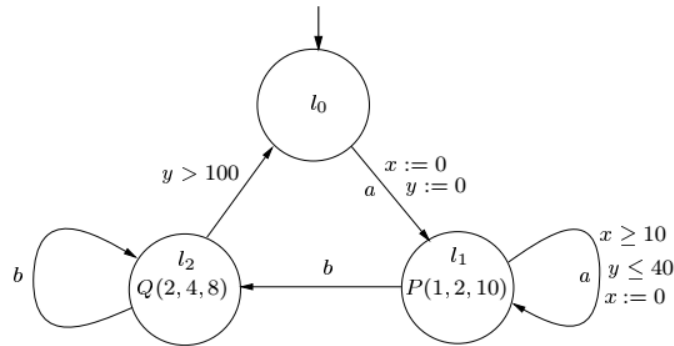
Demand Bound Function (DBF)



Request Bound Function (RBF) / Arrival Curve



Problem: Compute DBF/RBF of a TA (Efficiently)

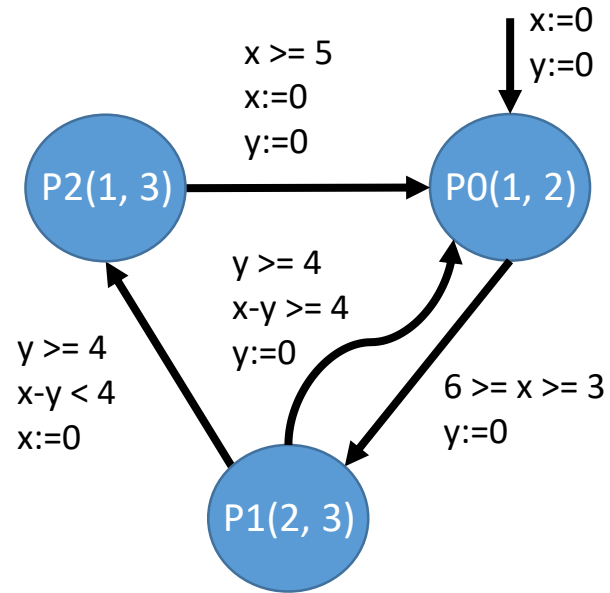


- If solved
 - The first step towards scalable schedulability analysis of general Task Automata models
 - Can efficiently check schedulability of a system composed of **independent** tasks each modeled as a Task Automaton

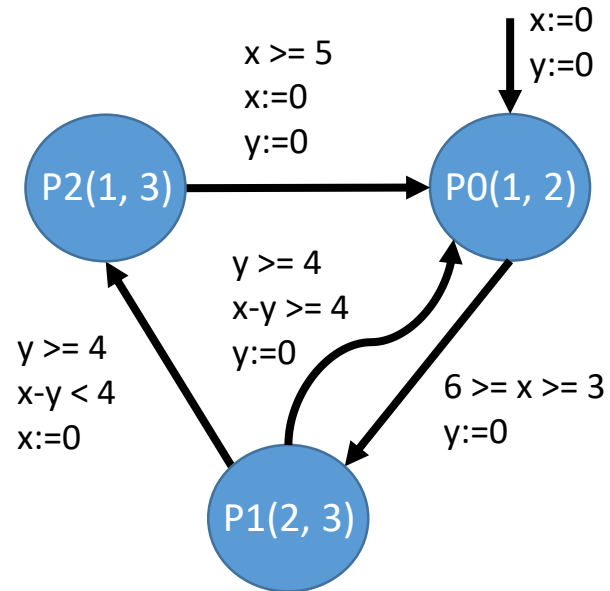
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Semantics of Task Automata

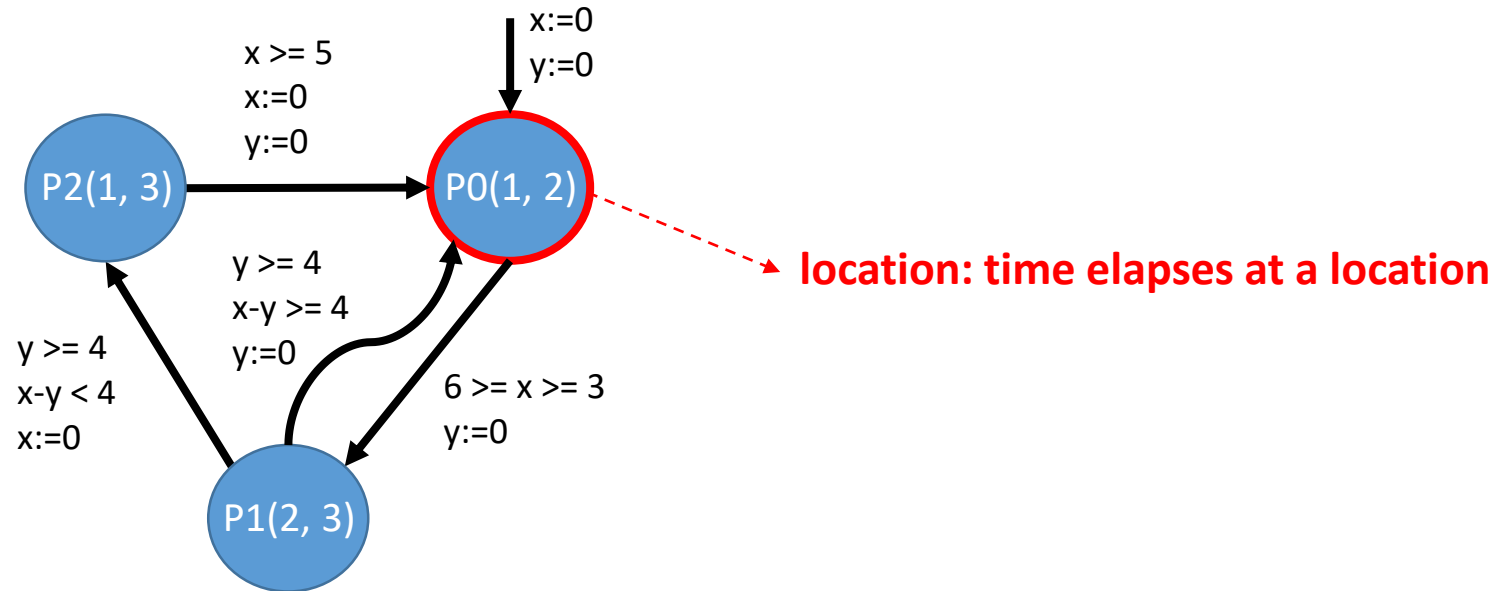


Semantics of Task Automata

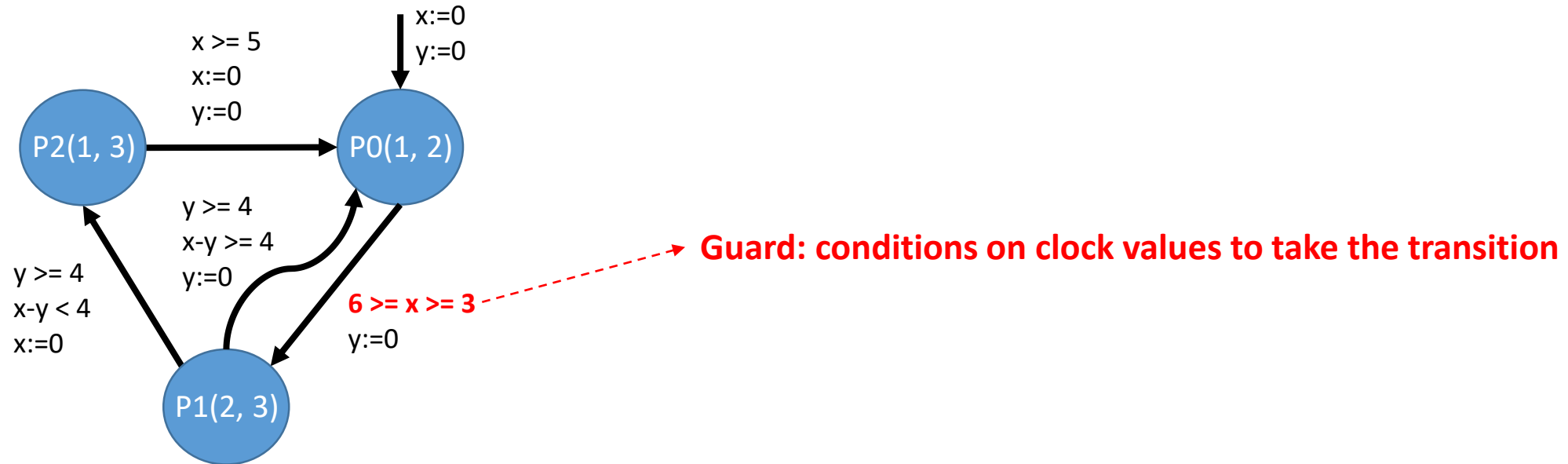


x, y: clocks

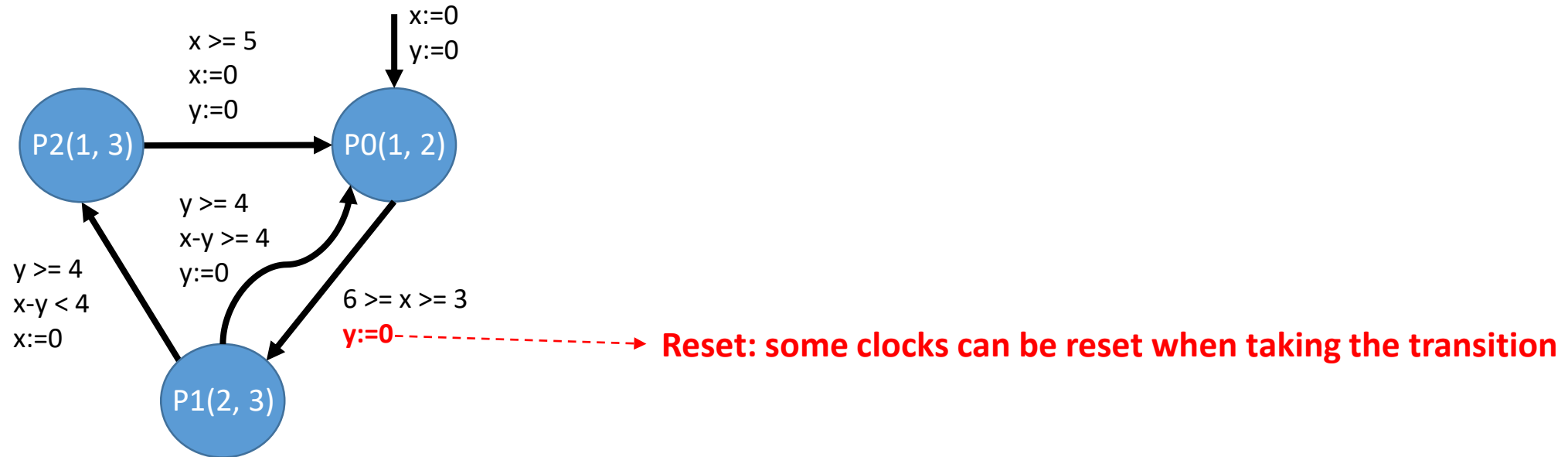
Semantics of Task Automata



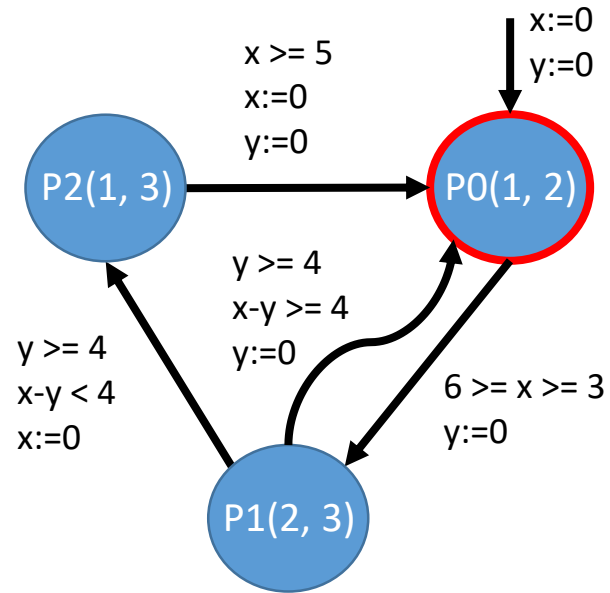
Semantics of Task Automata



Semantics of Task Automata



Semantics of Task Automata



A possible run:

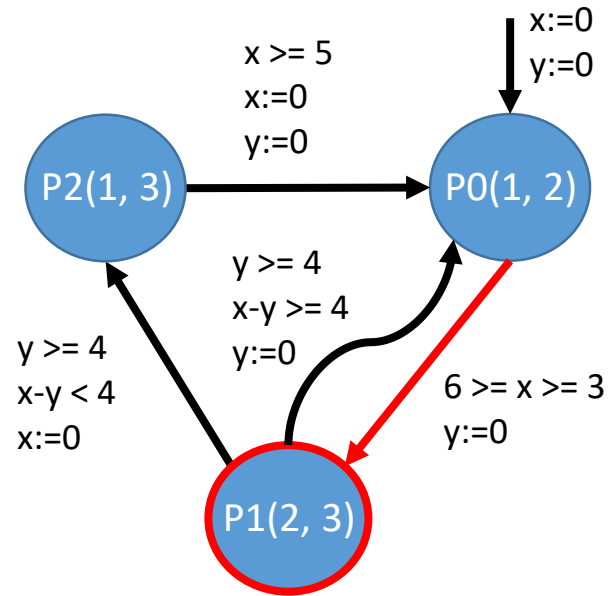


before

after

x=0
y=0

Semantics of Task Automata



A possible run:



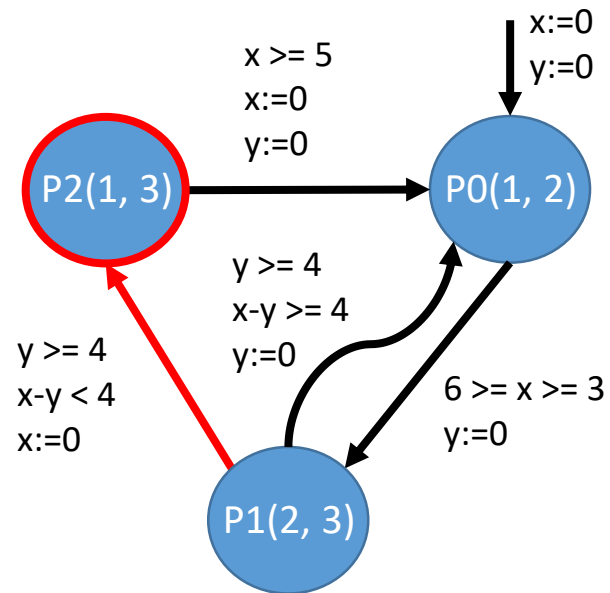
before

$x=4.5$
 $y=4.5$

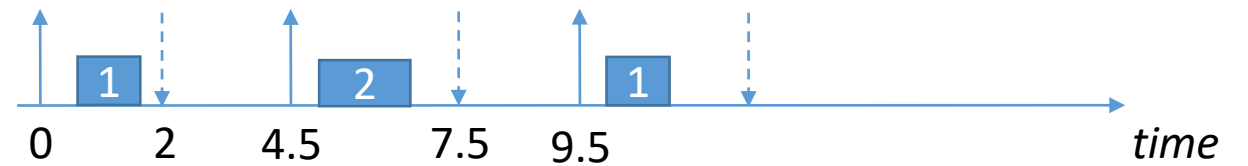
after

$x=0$ $x=4.5$
 $y=0$ $y=0$

Semantics of Task Automata



A possible run:



before

$x=4.5$
 $y=4.5$

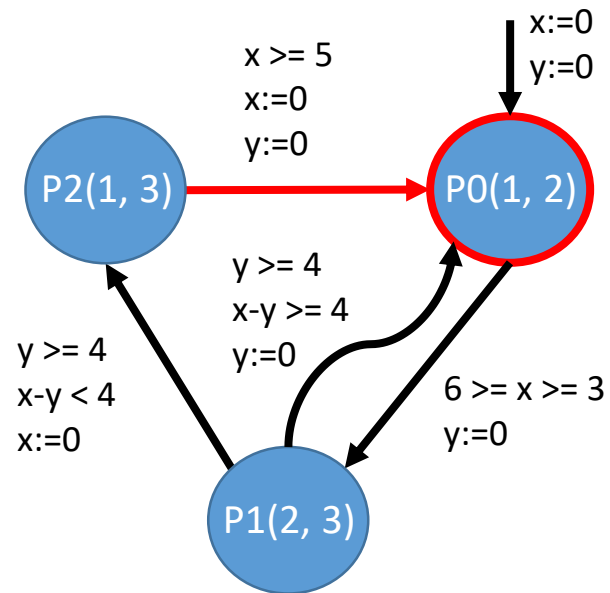
$x=9.5$
 $y=5$

after

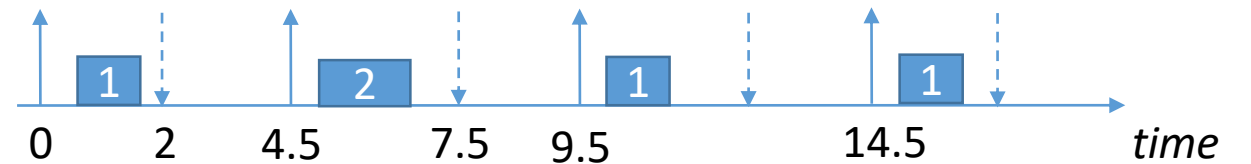
$x=0$
 $y=0$

$x=0$
 $y=5$

Semantics of Task Automata



A possible run:



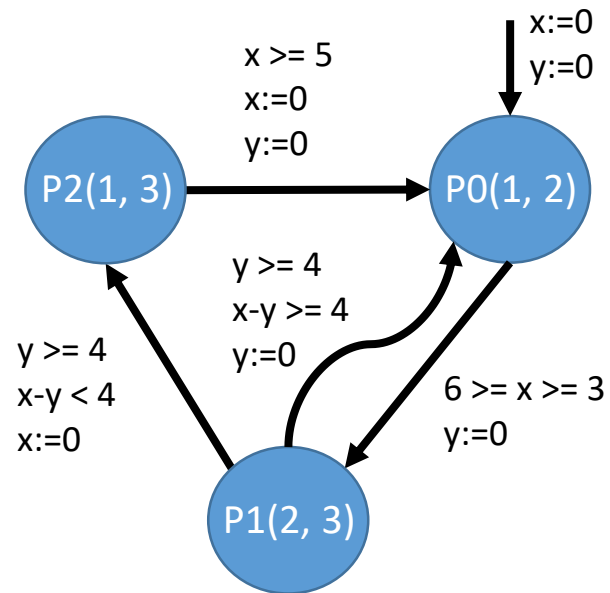
before

$x=4.5$
 $y=4.5$ $x=9.5$
 $y=5$ $x=5$
 $y=10$

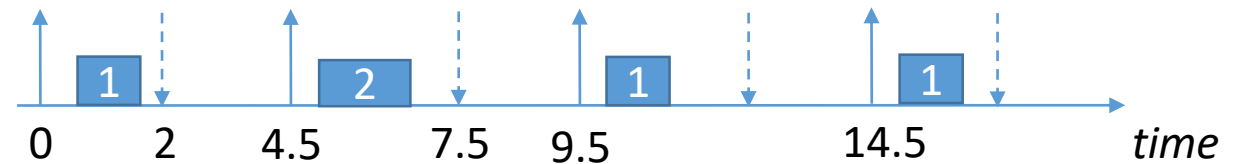
after

$x=0$
 $y=0$ $x=4.5$
 $y=0$ $x=0$
 $y=5$ $x=0$
 $y=0$

Semantics of Task Automata



A possible run:

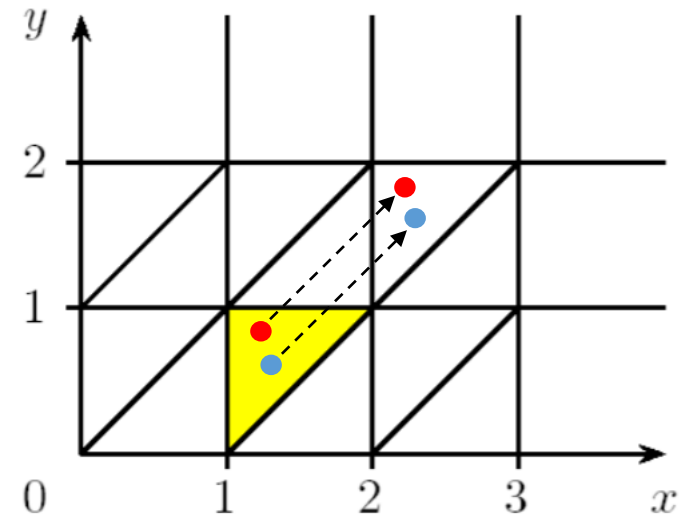
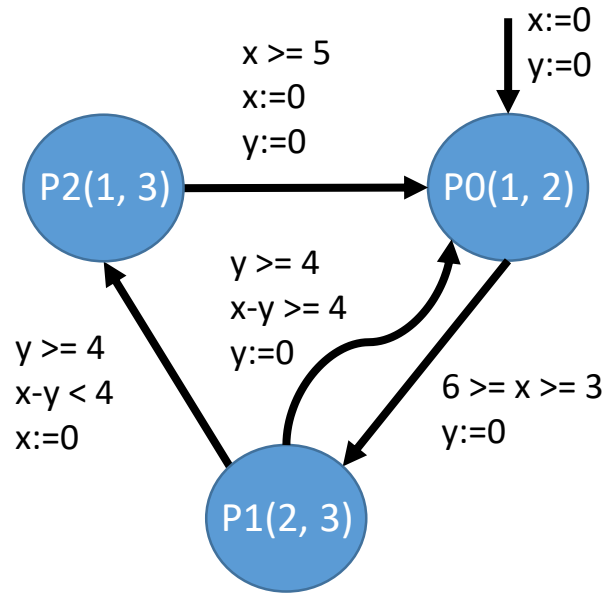


<i>before</i>		$x=4.5$	$x=9.5$	$x=5$
		$y=4.5$	$y=5$	$y=10$
<i>after</i>	$x=0$	$x=4.5$	$x=0$	$x=0$
	$y=0$	$y=0$	$y=5$	$y=0$

Problem: Infinitely many possible runs

Region

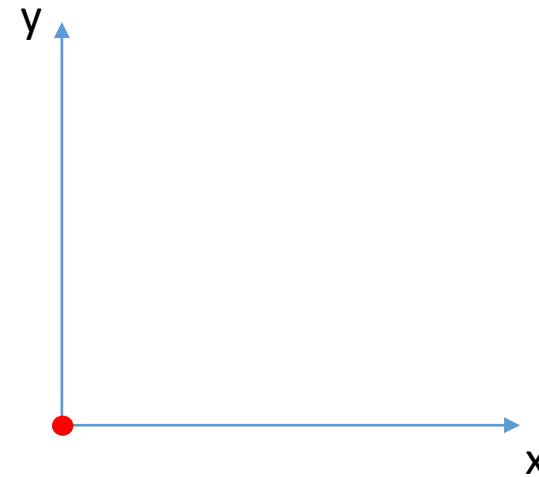
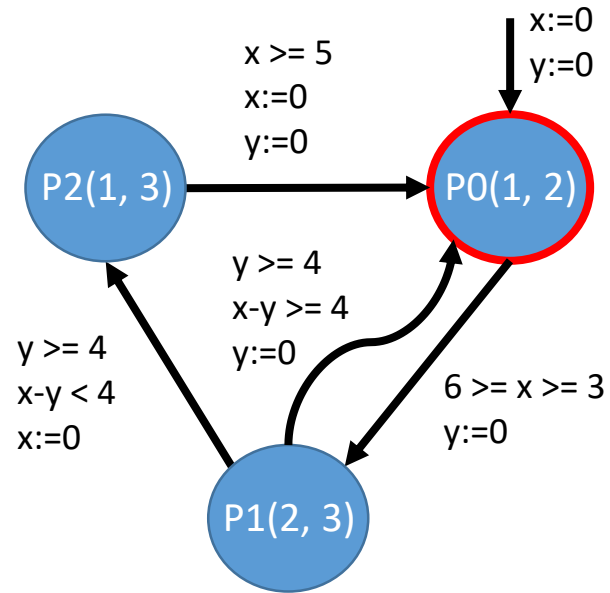
From Infinite to Finite State Space



Region preserves reachability

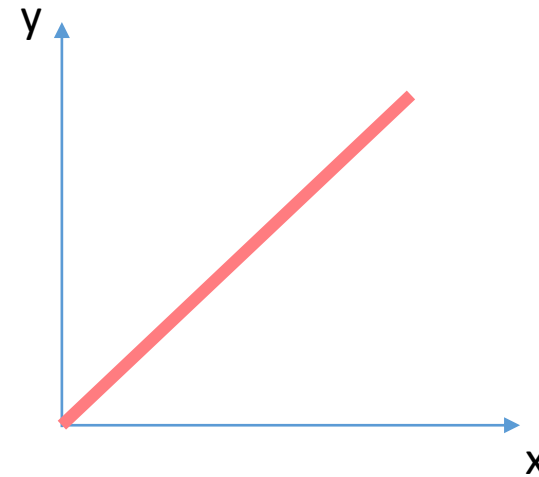
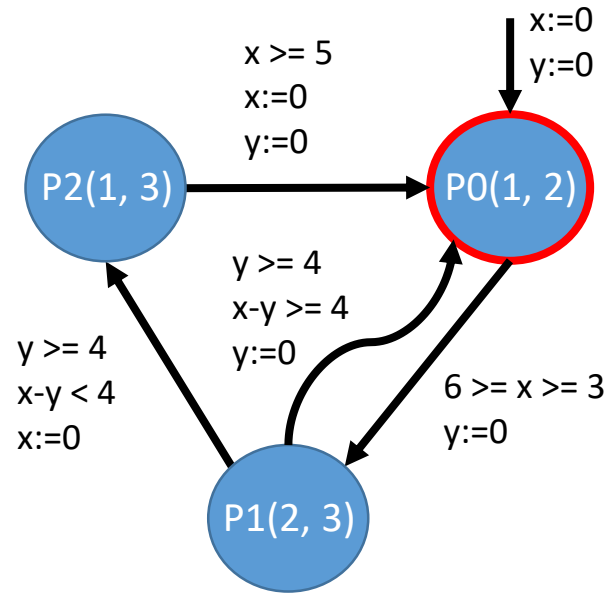
Zone

More compact state space representation



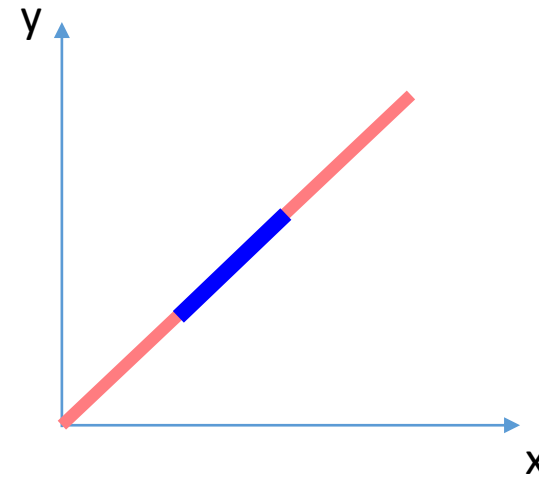
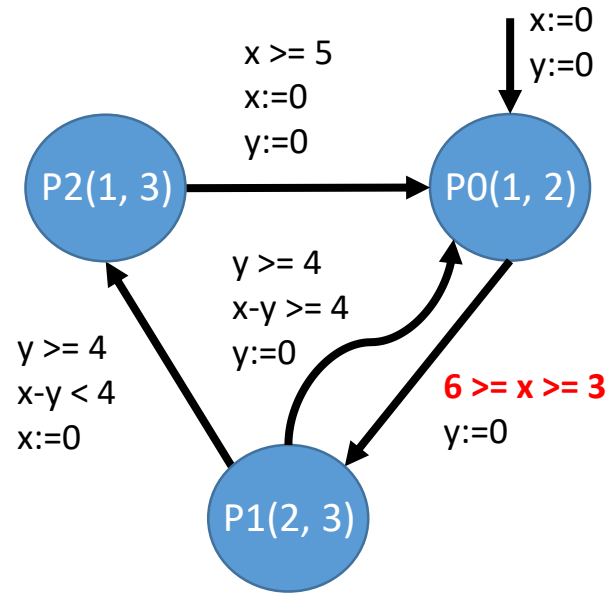
Zone

More compact state space representation



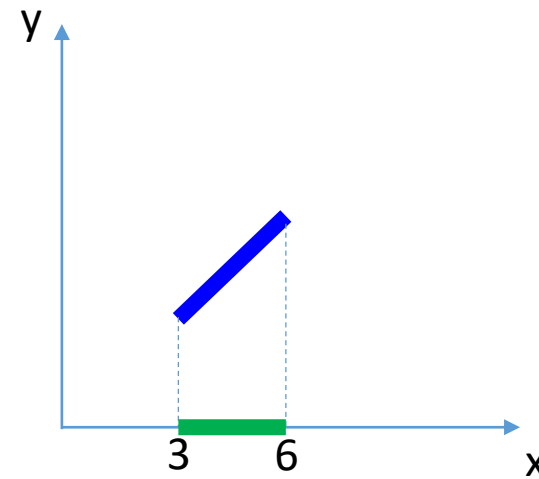
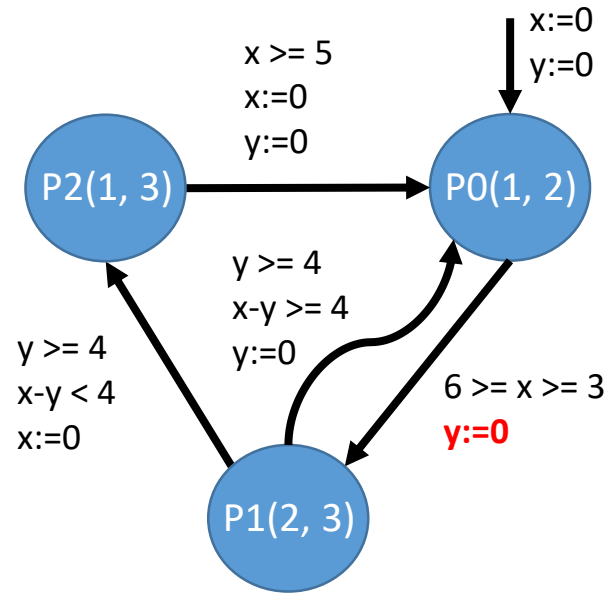
Zone

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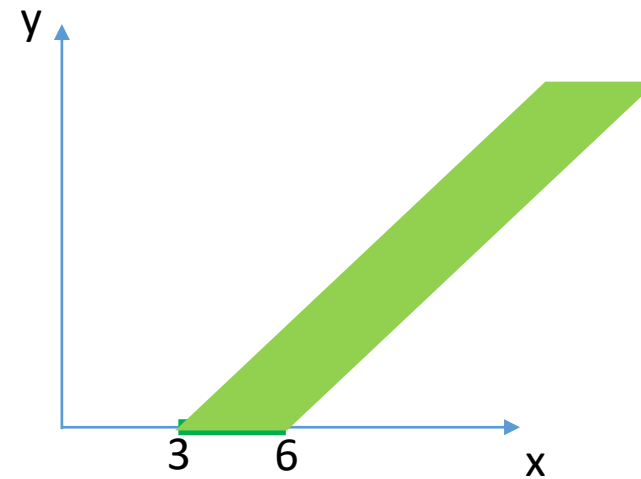
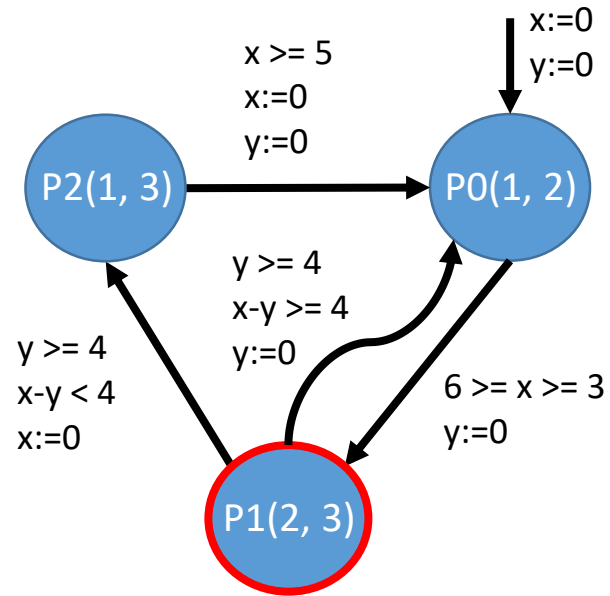
Zone

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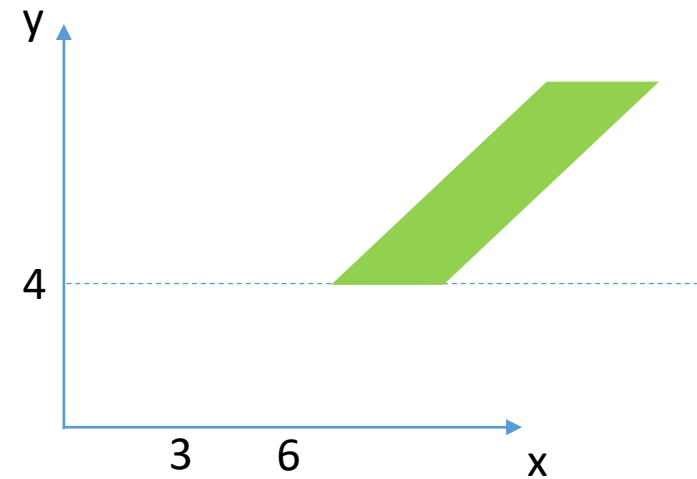
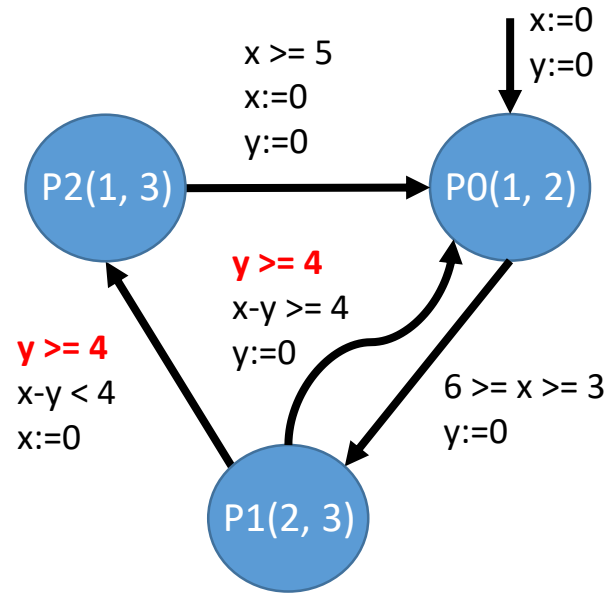
Zone

More compact state space representation



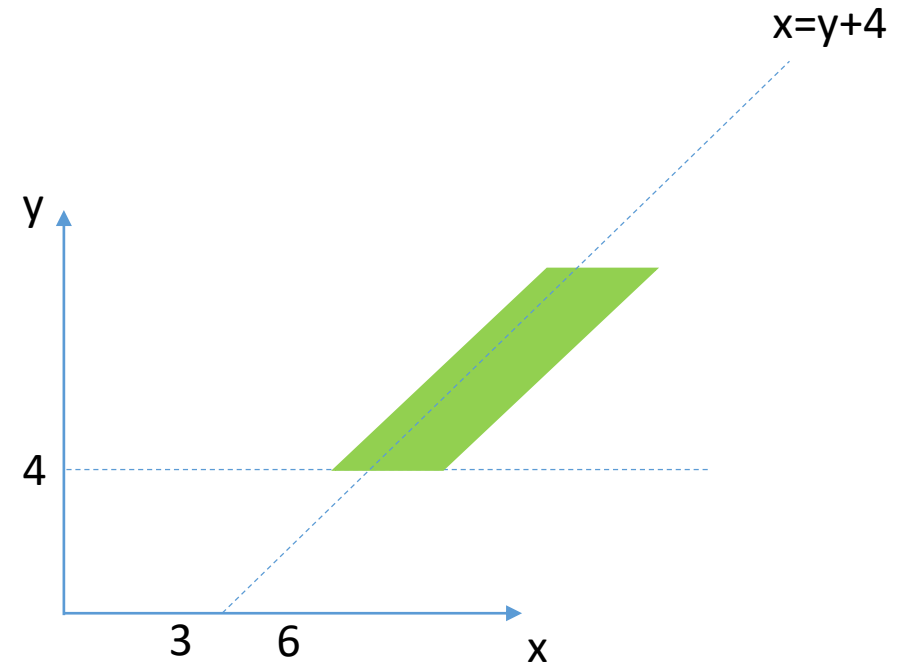
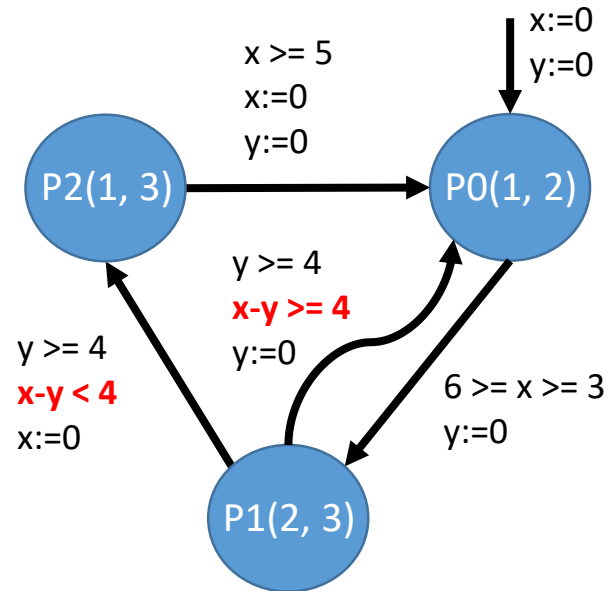
Zone

More compact state space representation



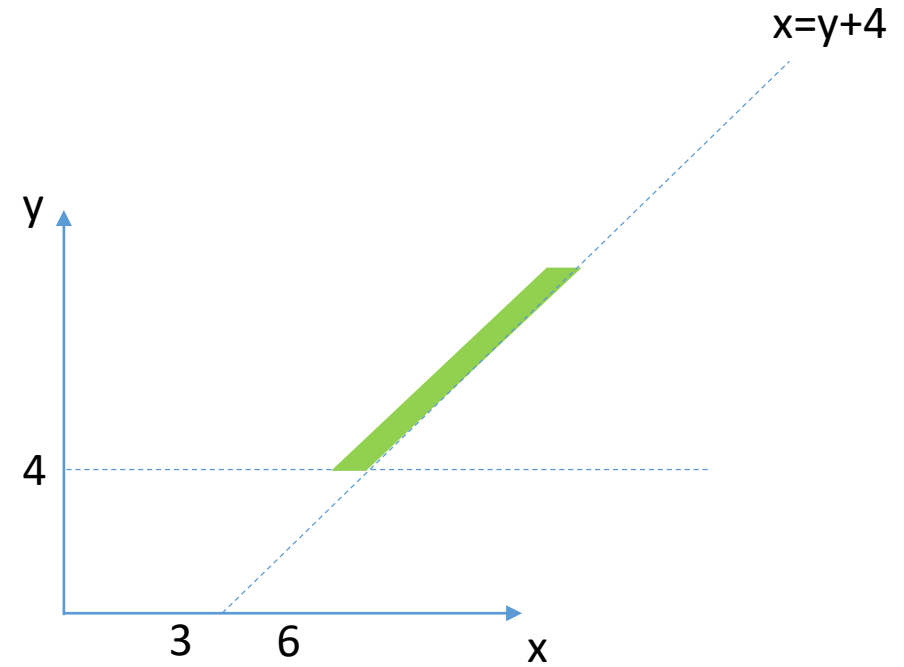
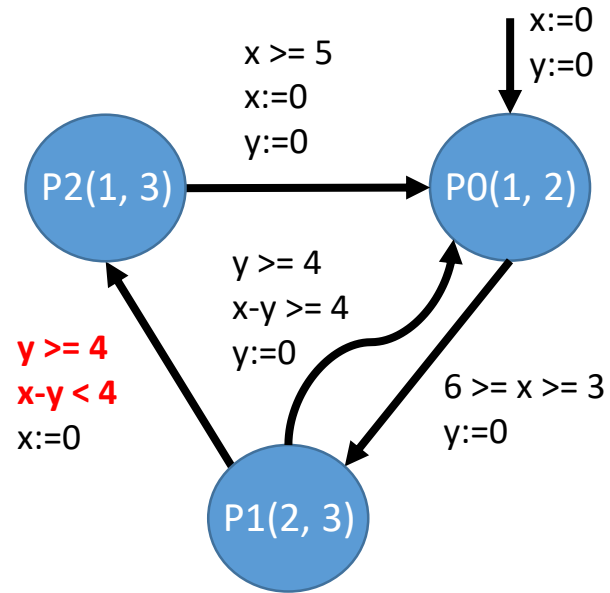
Zone

More compact state space representation



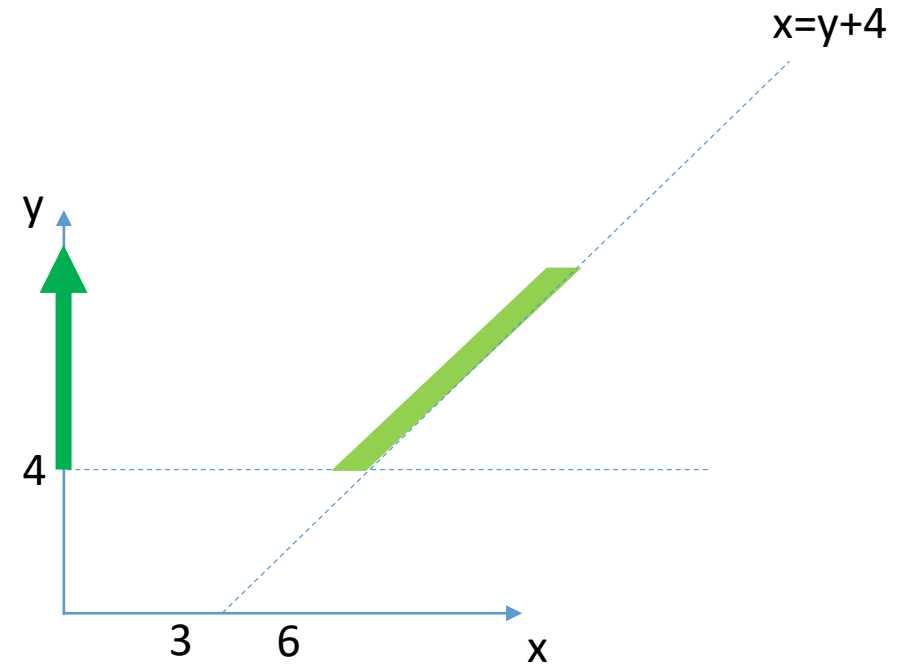
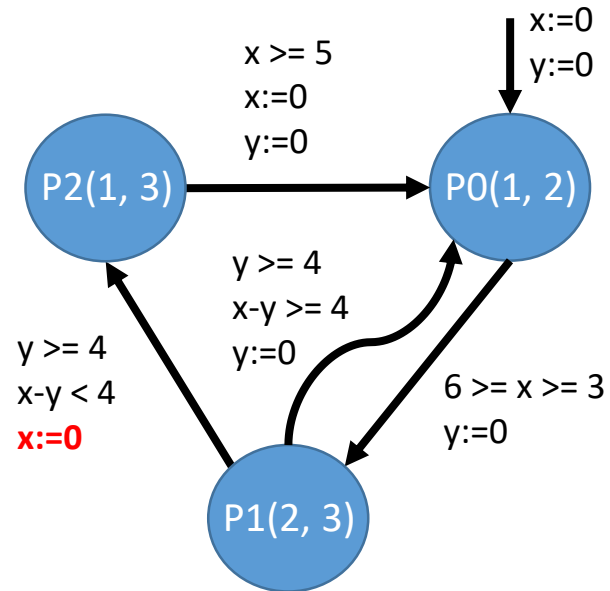
Zone

More compact state space representation



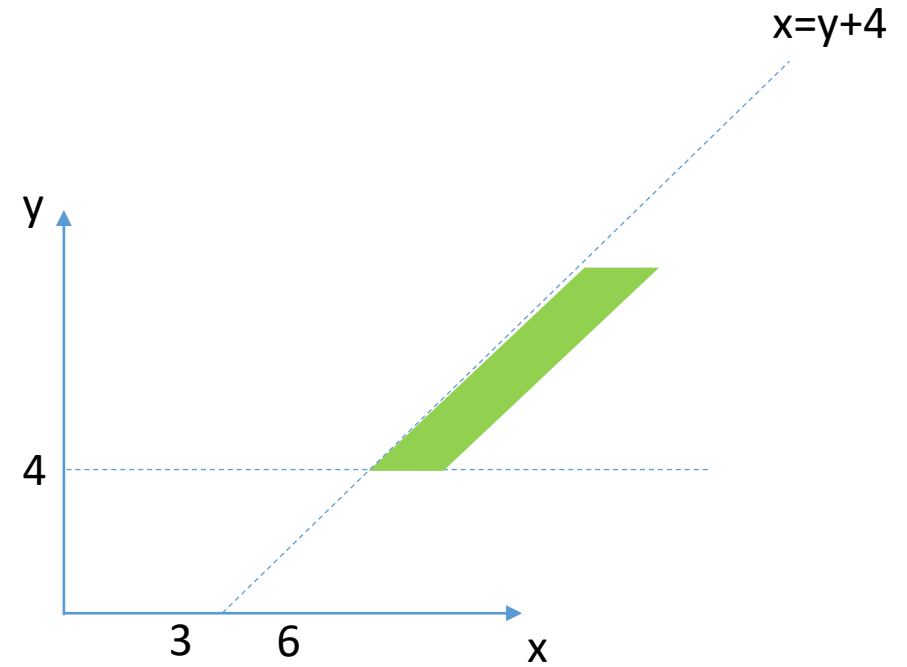
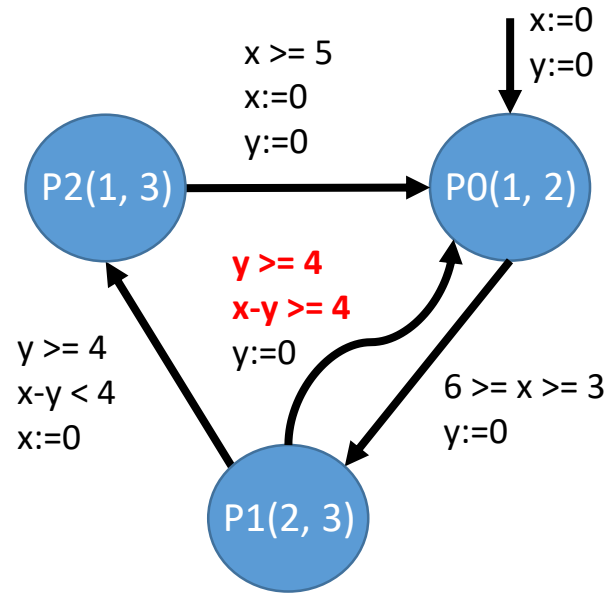
Zone

More compact state space representation



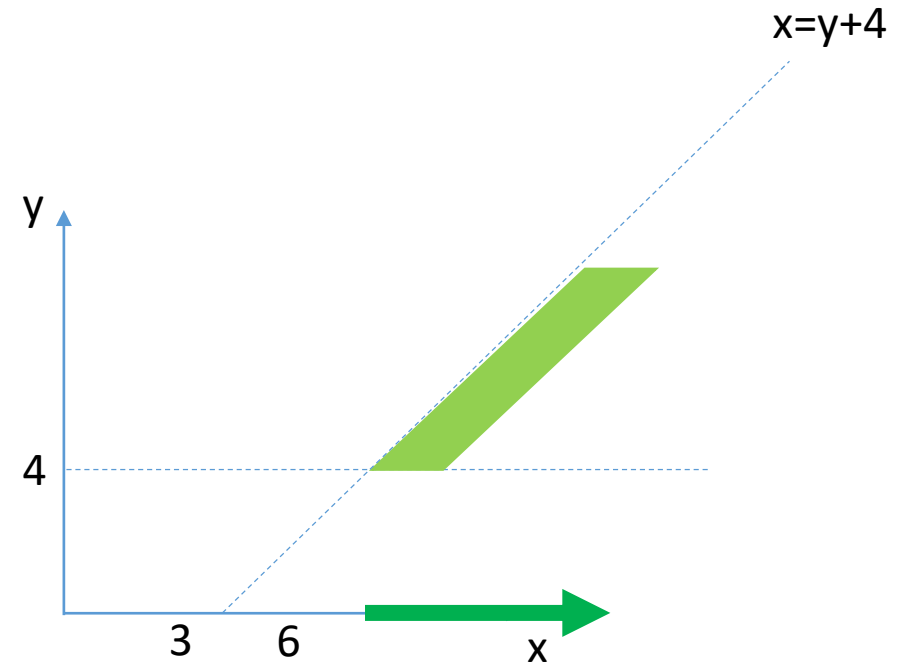
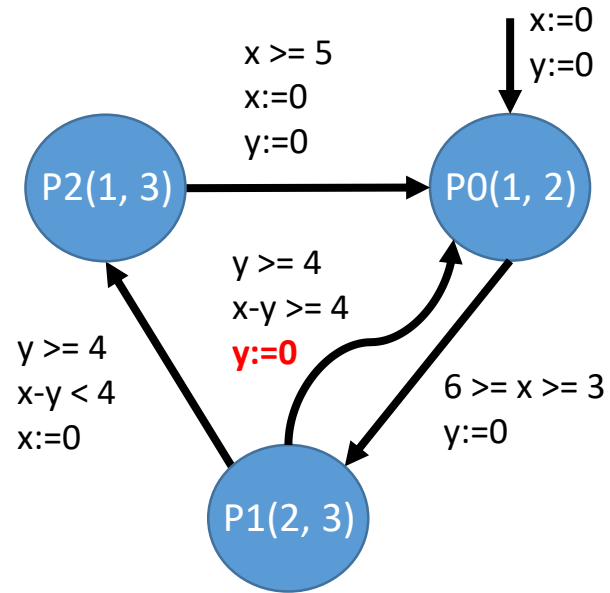
Zone

More compact state space representation



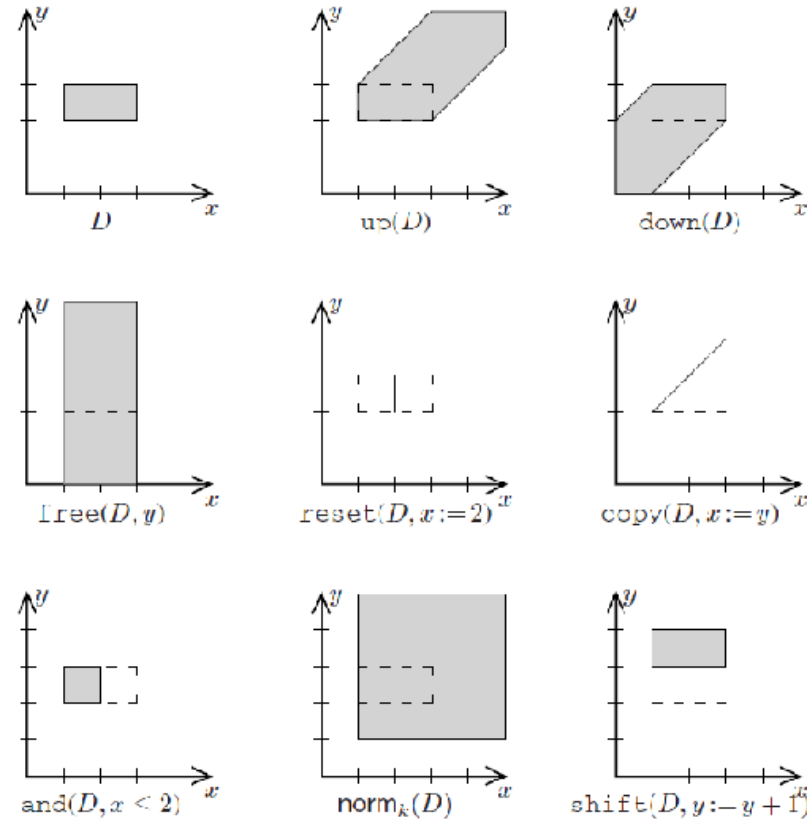
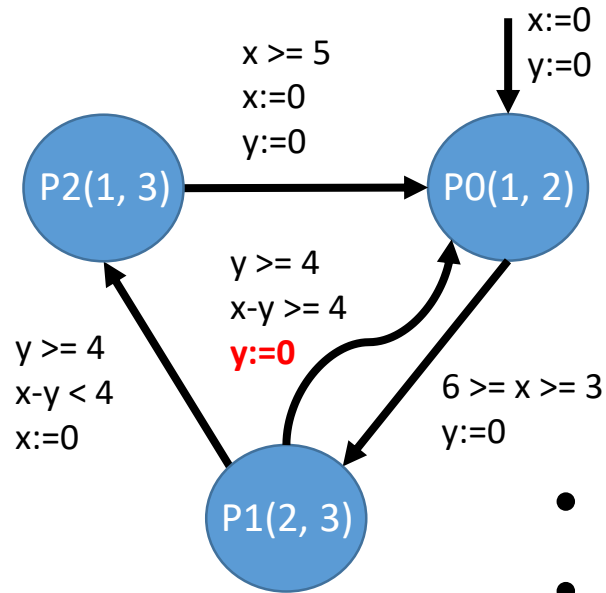
Zone

More compact state space representation



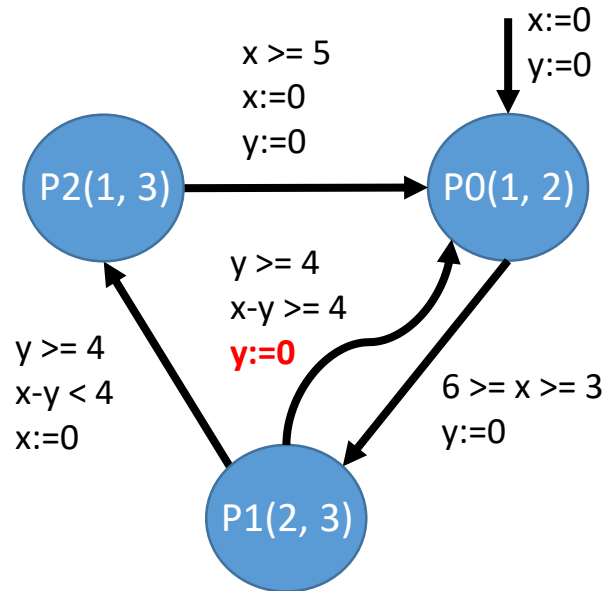
Zone

More compact state space representation

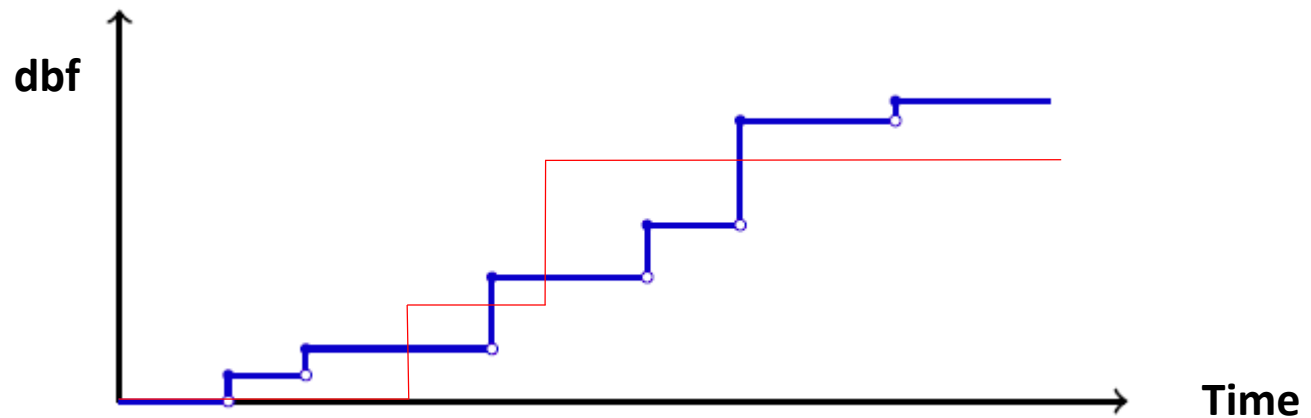


- Zone also preserves reachability
- A Zone is the solution set of a clock constraint
 - maximal set of clock assignments satisfying the constraint
 - Such sets can be efficiently represented and stored in memory as DBMs (Difference Bound Matrices)

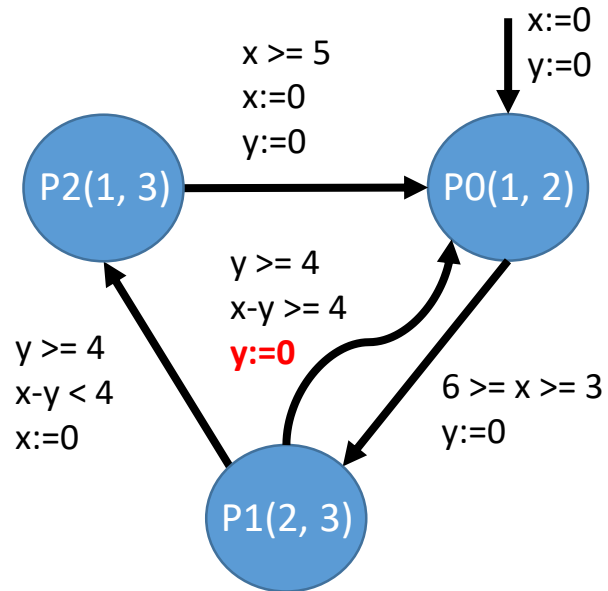
Computing DBF based on Zone



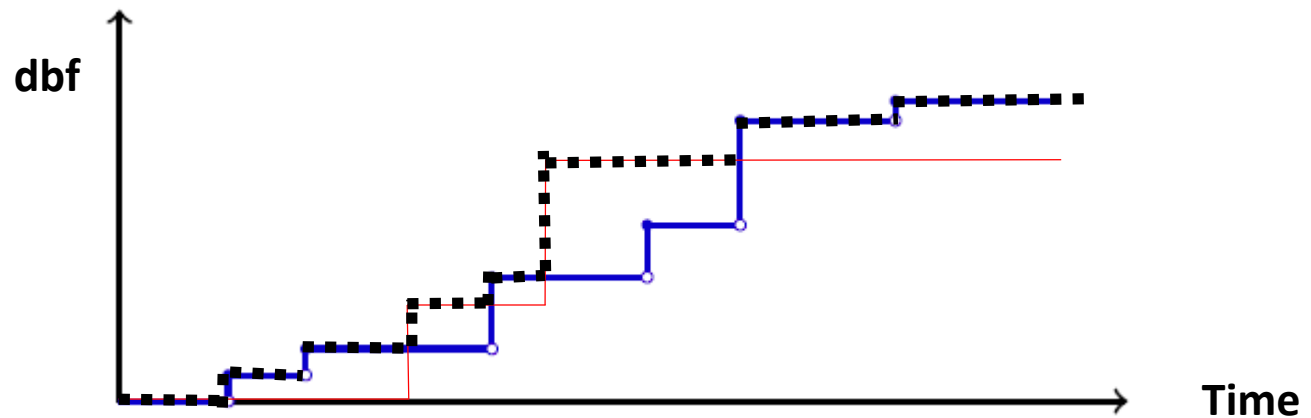
- Traverse the graph with Zone
- Split a Zone upon “branches”
- Each traversal trace generates a “demand trace”
 - the “worst case” must be on the border of Zones
- DBF = Maximization of all “demand functions”



Computing DBF based on Zone



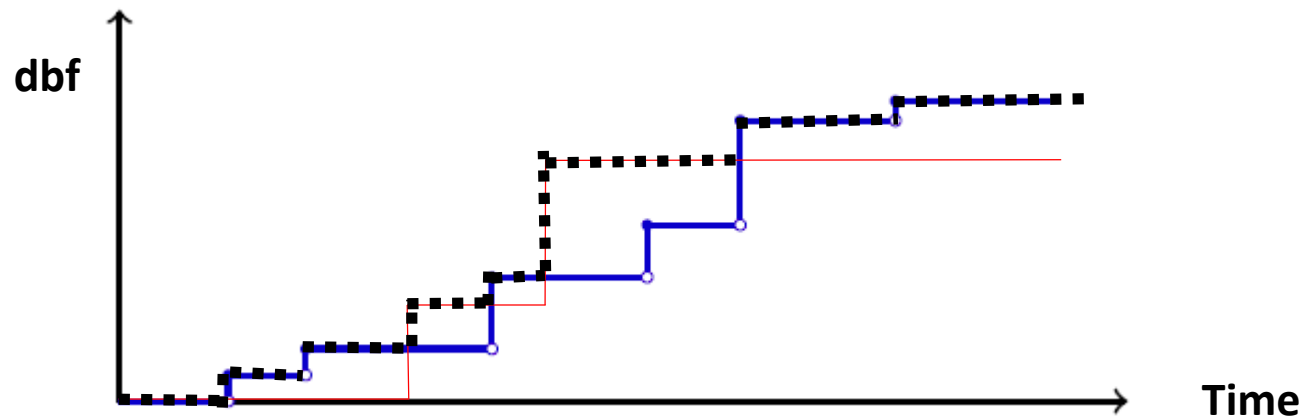
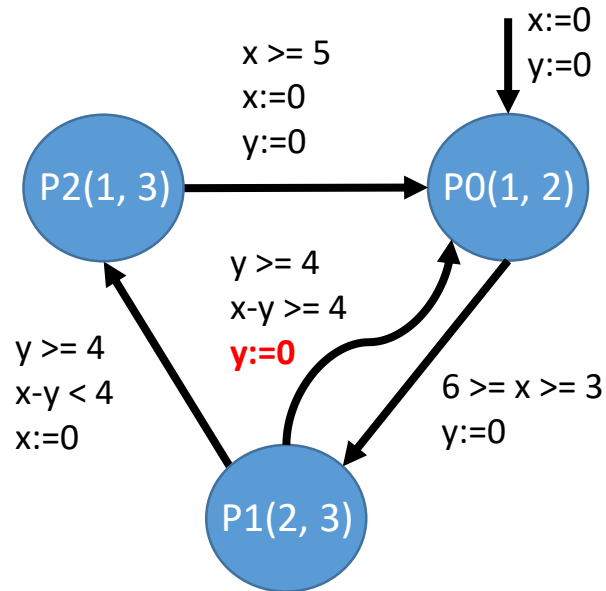
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Computing DBF based on Zone

Exponential

- Traverse the graph with Zone
- Split a Zone upon “branches”
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- DBF = Maximization of all “demand functions”



Summary

- Ultimate goal
 - Utilizing efficient abstraction and analysis techniques in scheduling theory to analyze the TA systems
- As the first step
 - Computing the DBF/RBF of an individual TA
- Starting point
 - Zone is a good starting point as state-space representation
 - Efficient storage and operations are known