Extending Dynamic Software Product Lines with Temporal Constraints

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Adaptive Cloud Environments

• Cloud computing supports construction of customized adaptable environments

"Cloud computing is a model for enabling ubiquitous, convenient, ondemand network access to a shared pool of configurable computing resources (e.g., networks, servers, storage, applications, and services) **that can be rapidly provisioned and released with minimal management effort or service provider interaction**."^[1]

 A cloud environment is a set of cloud services provisioned for running an application

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[1] P. Mell and T. Grance, "The NIST definition of cloud computing," Computer Security Division, Information Technology Laboratory, National Institute of Standards and Technology, Tech. Rep., 2011.



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Cloud Providers Configuration Variability

- Wide range of configurable cloud services
- Complex configuration rules and constraints



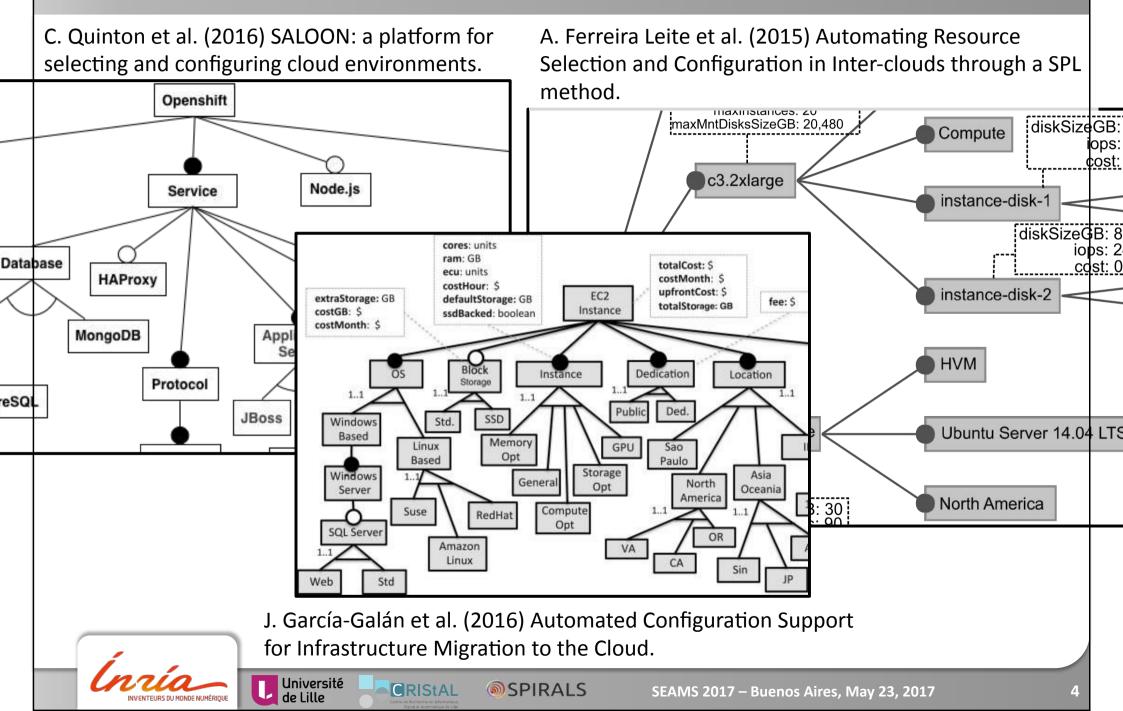
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SPLs for Automated Cloud Configuration



Dynamic Software Product Lines

• High variability with adaptive capabilities







Dynamic Software Product Lines

High variability with adaptive capabilities

• DSPL vs SPL

- Features can be (re)bound at runtime
- Adaptive system vs systems family
- Variability model central to both



Dynamic Software Product Lines

High variability with adaptive capabilities

DSPL vs SPL

- Features can be (re)bound at runtime
- Adaptive system vs systems family
- Variability model central to both
- Adaptation in DSPLs

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- A context change is mapped to a request to include or exclude a set of features from the current configuration
- SPL analysis is used to derive valid configurations

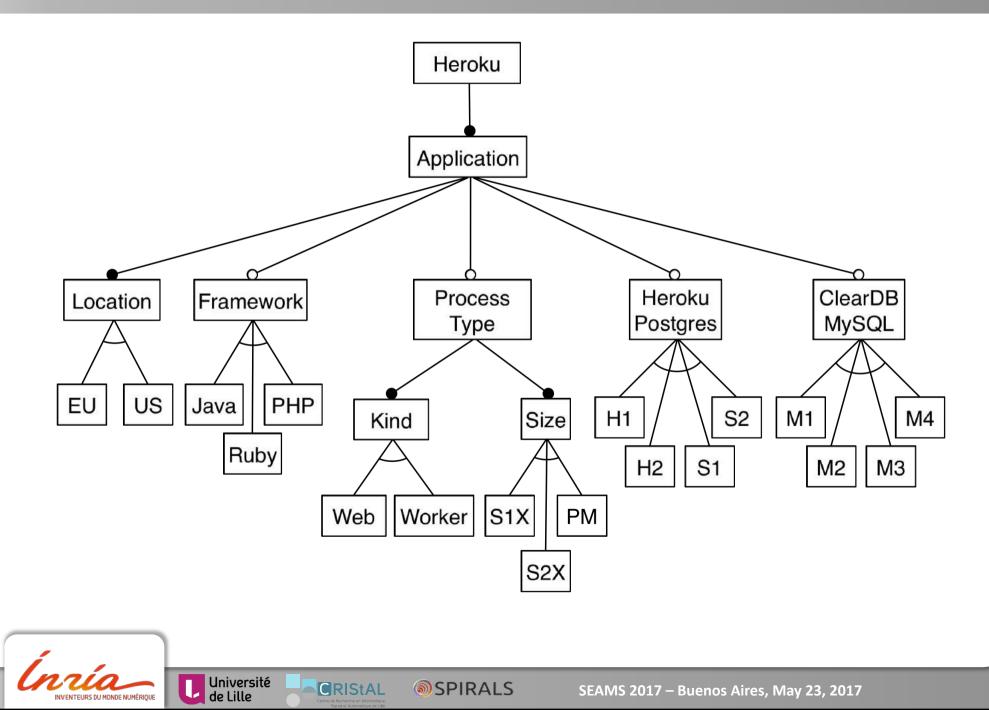
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Cloud Computing Environment

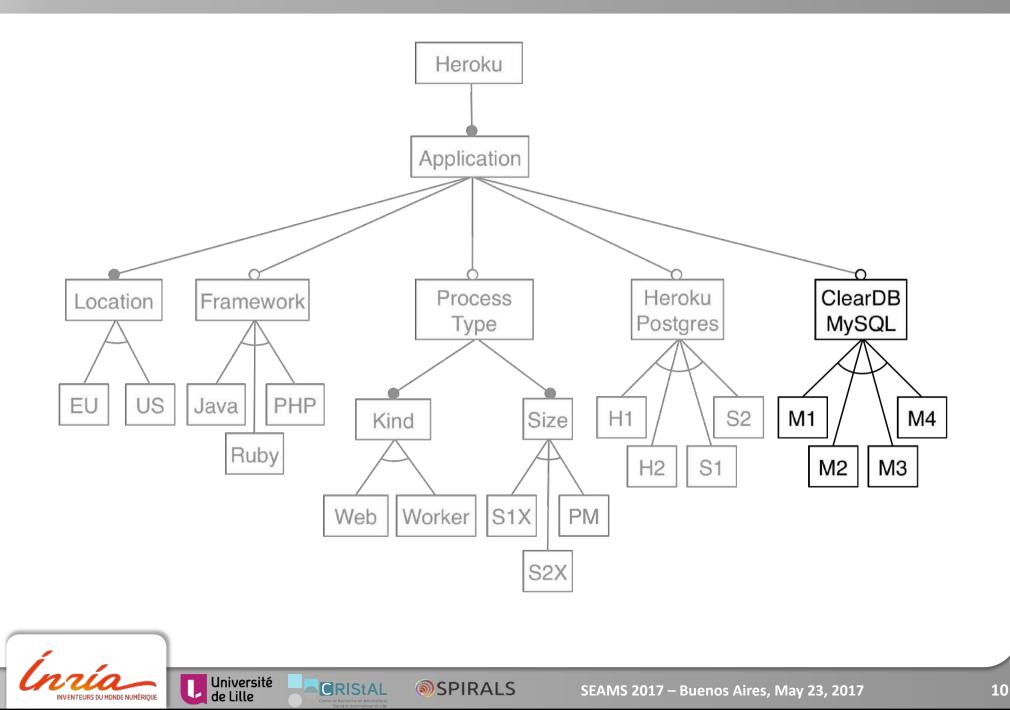
- Reconfiguration mechanisms are provider-dependent and heterogeneous
 - May depend on initial or previous configurations
 - Alternative ways to reconfigure
- Compliance to variability model is not enough
 - Does not ensure valid and safe reconfigurations

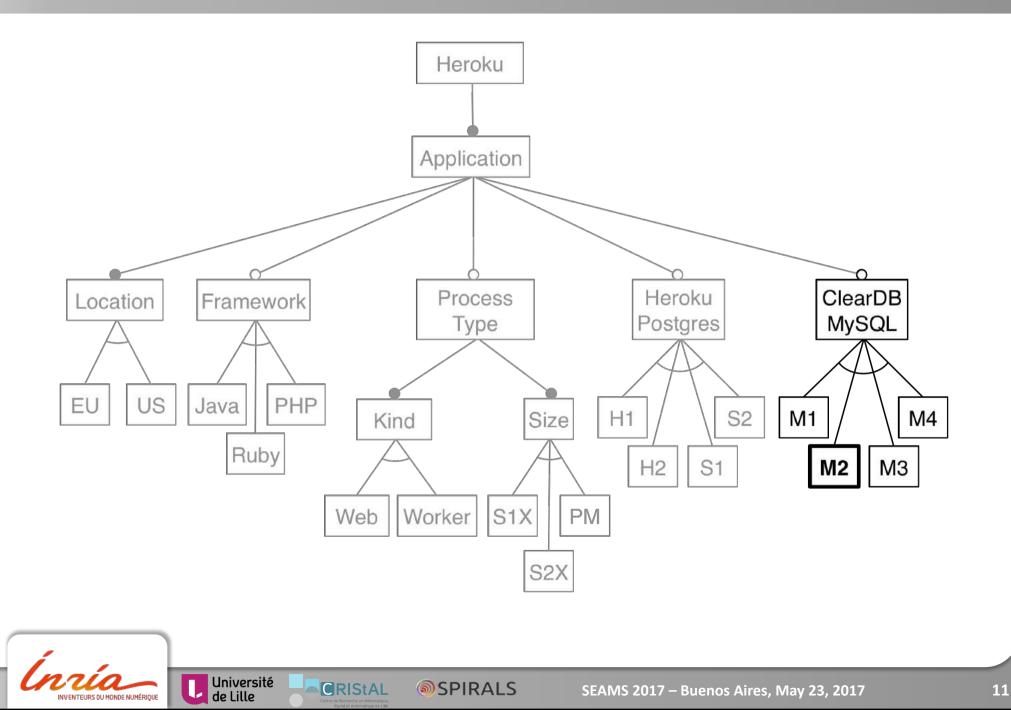


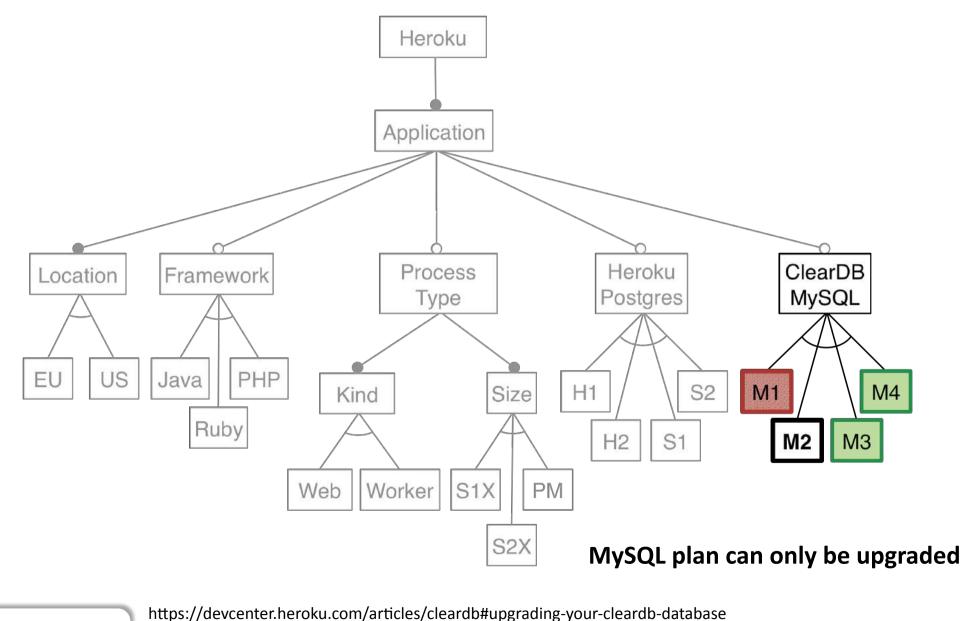




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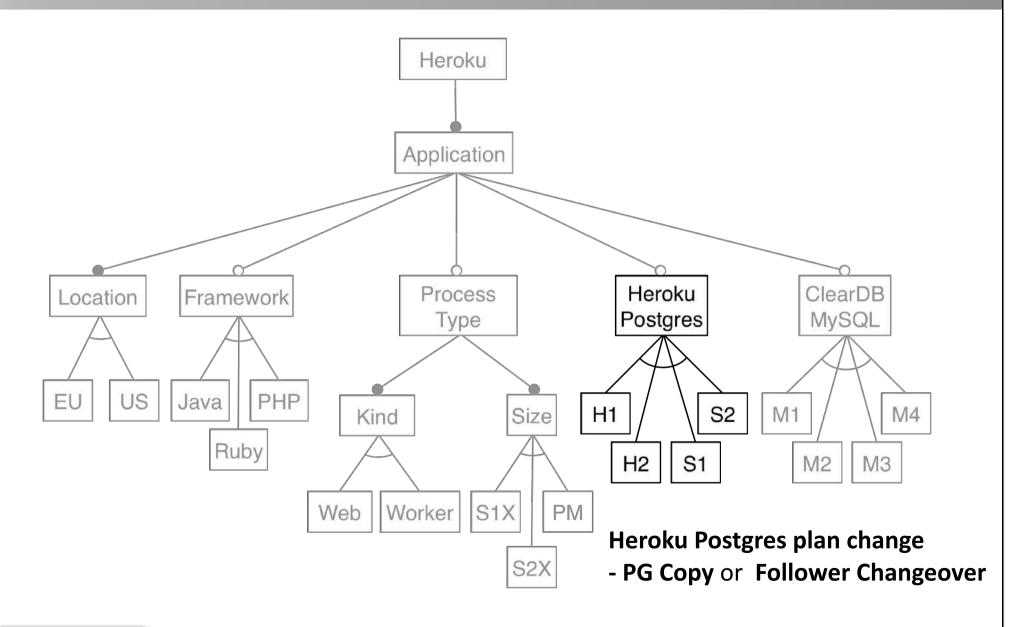


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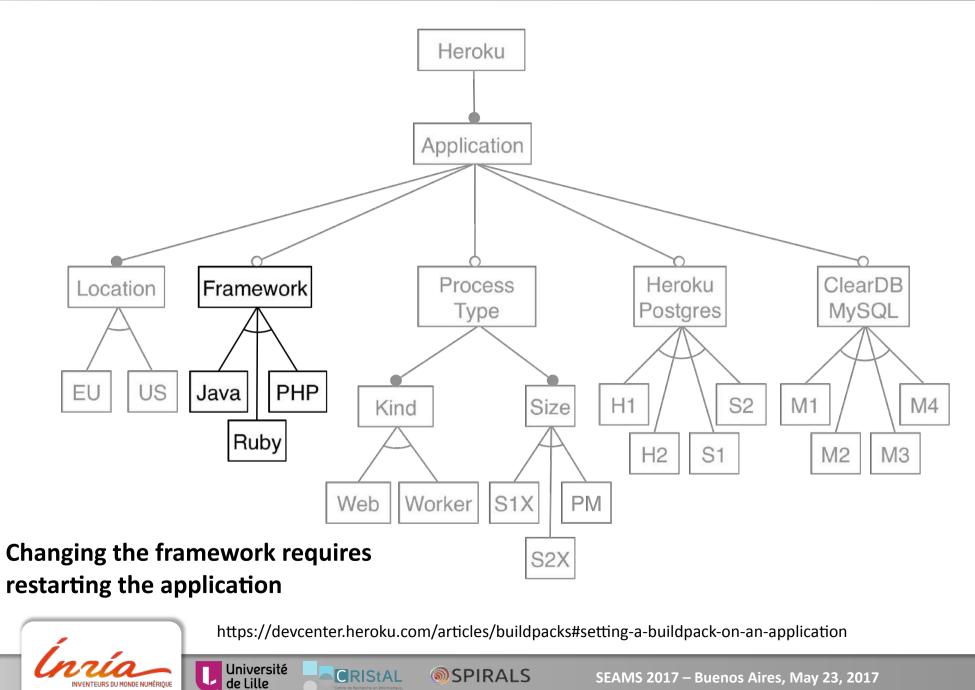
https://devcenter.heroku.com/articles/upgrading-heroku-postgres-databases

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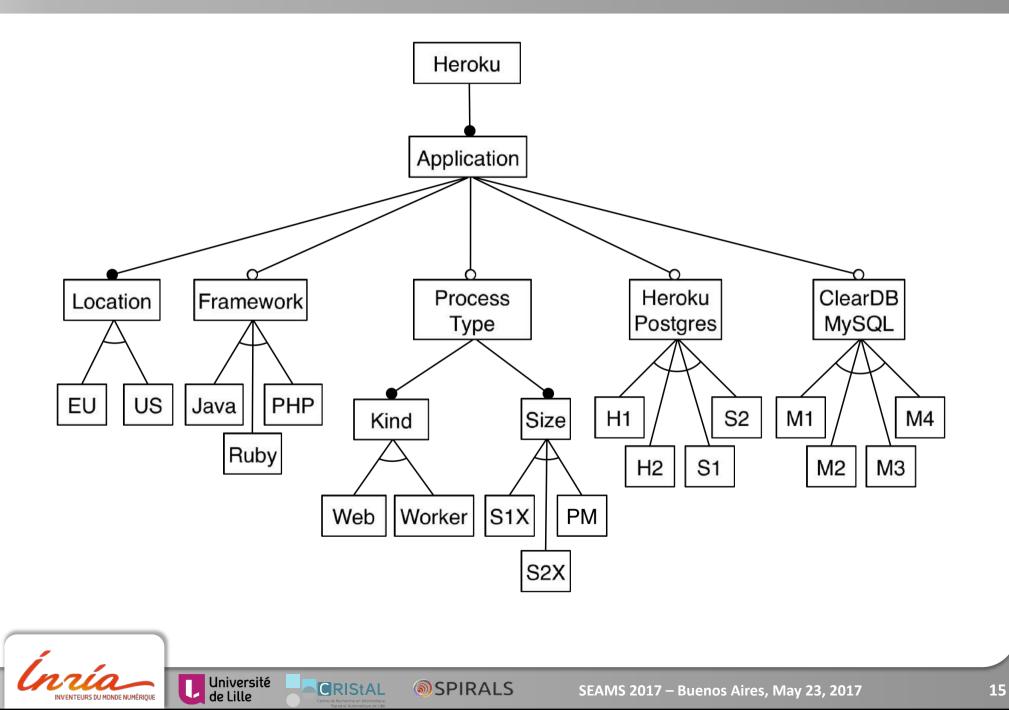


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Limitations in DSPLs

- Seminal works on DSPLs highlight the need for validating transitions between system configurations
 - systems should evolve through safe migration paths^[6]
 - dynamic constraints on allowed transitions must be considered^[7]
- Validation is mostly limited to compliance to a variability model

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[6] B. Morin, O. Barais, J. M. Jezequel, F. Fleurey, and A. Solberg, "Models@run.time to support dynamic adaptation," Computer, vol. 42, no. 10, pp. 44–51, Oct 2009.

[7] A. Hubaux and P. Heymans, "On the evaluation and improvement of feature-based configuration techniques in software product lines," in Proc. 31st Int. Conf. Software Engineering (ICSE'09), Vancouver, Canada, May 2009, pp. 367–370.



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Problem statement

- How to model constraints over the adaptation behavior?
 - Temporal dependencies between features and reconfiguration operations
- How to reason over a variability model with reconfiguration constraints to find reconfigurations that meet a given criteria?
 - e.g. reduced downtime or costs

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Proposed approach

- Combine variability models with temporal constraints and reconfiguration operations
 - Leverage concepts and solutions from model checking





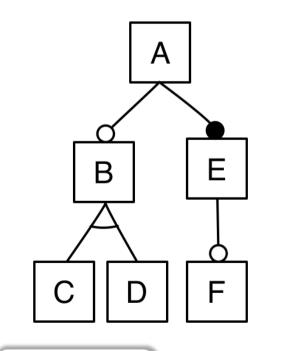
Feature Models and Transition Systems

- Feature model M = (F, C)
 - -F is the set of features

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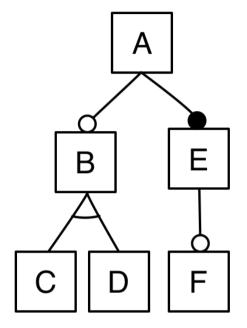
 $-C \subseteq \mathcal{P}(F)$

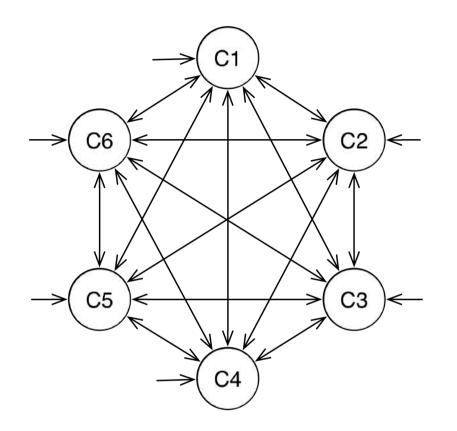


 $C1 = \{A, E\}$ $C2 = \{A, E, F\}$ $C3 = \{A, B, C, E\}$ $C4 = \{A, B, C, E, F\}$ $C5 = \{A, B, D, E\}$ $C6 = \{A, B, D, E, F\}$

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DSPLs as Transition Systems





 $\begin{array}{l} C1 = \{A, E\} \\ C2 = \{A, E, F\} \\ C3 = \{A, B, C, E\} \\ C4 = \{A, B, C, E, F\} \\ C5 = \{A, B, D, E\} \\ C6 = \{A, B, D, E, F\} \end{array}$

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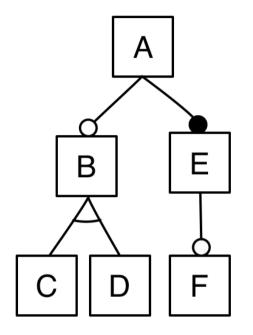
"A DSPL's execution can be abstracted as a highly connected state machine where the states are the possible system configurations and the transitions the migration paths."^[6]

[6] B. Morin, O. Barais, J. M. Jezequel, F. Fleurey, and A. Solberg, "Models@run.time to support dynamic adaptation," Computer, vol. 42, no. 10, pp. 44–51, Oct 2009.



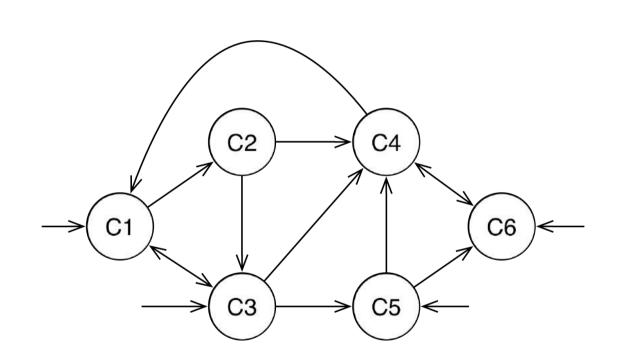
DSPLs as Transition Systems

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Temporal properties

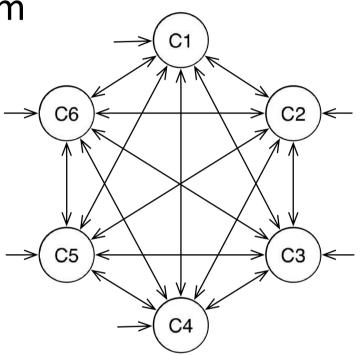
- A temporal property defines a condition over the executions of a transition system
 - Execution:
 - $\rho = s_0 s_1 s_2 s_3 s_4 \dots$
 - $s_i
 ightarrow s_{i+1}$ is a transition

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A property is a set executions



 A system exhibits a property if all its executions are part of the property set

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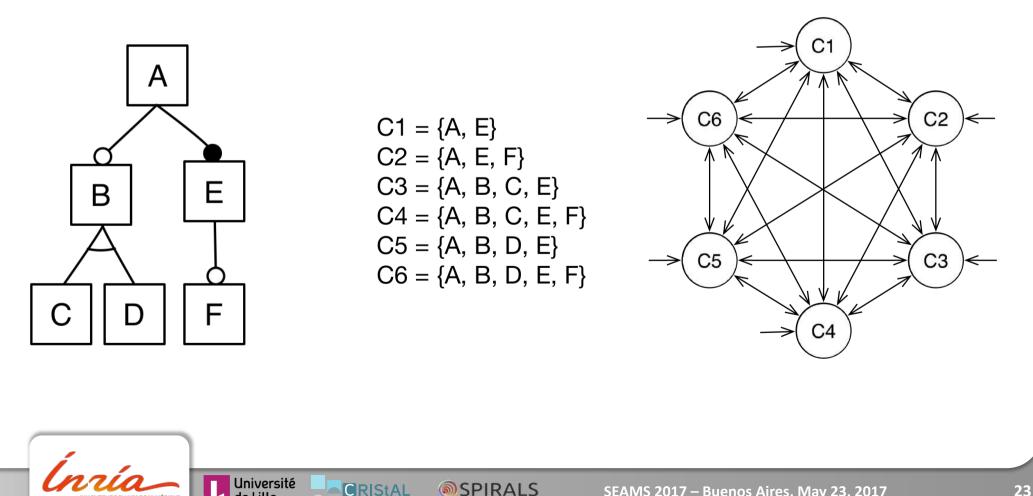


Feature Models and Transition Systems

• Feature model M = (F, C)

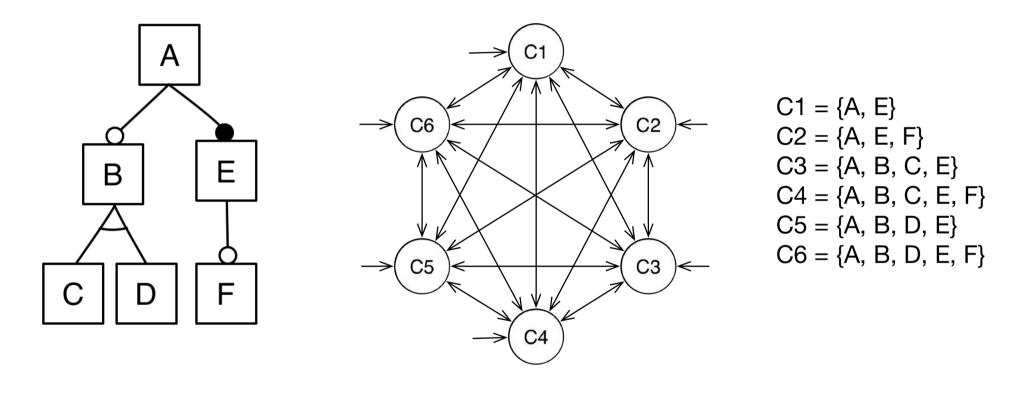
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• Transition system $TS_M = (S, I, R, AP, L)$ -S = I = C, $R = S \times S$, AP = F, L(x) = x



Temporal properties

• A temporal property is a condition over the executions of a transition system



 $P = \{s_0 s_1 s_2 s_3 \dots \mid C \in L(s_i) \leftrightarrow D \notin L(s_{i+1})\}$

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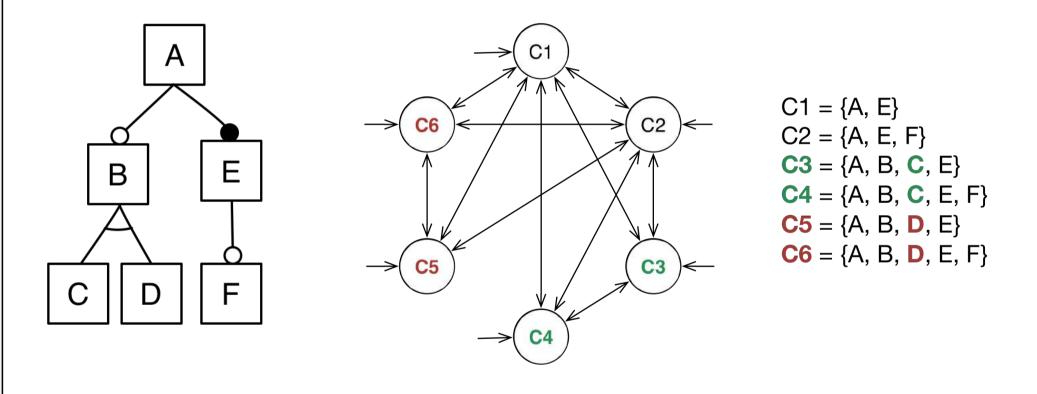


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Temporal properties

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Linear Temporal Logic (LTL)

- Defines temporal properties over transition systems
- Combines propositional logic with temporal operators (always, eventually, until)

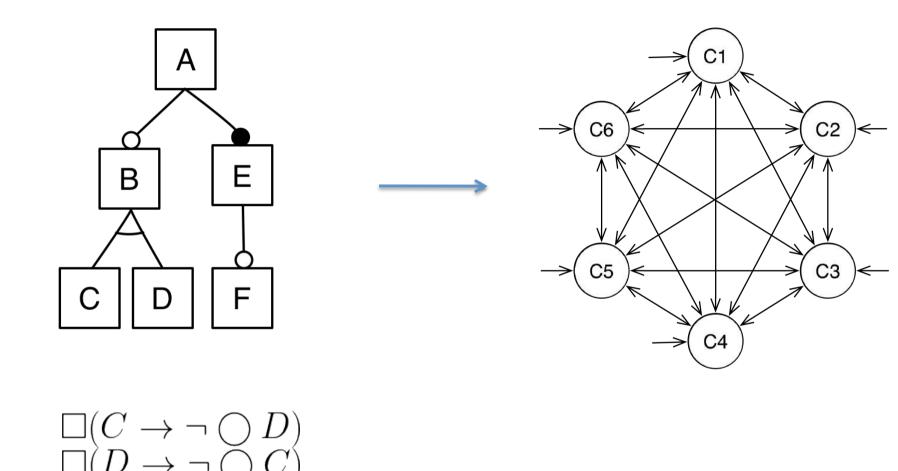
 $-\Box A$

// always A $-\Box(M2 \rightarrow \neg \bigcirc M1)$ // always (M2 is not followed by M1) $-\Box(M2 \rightarrow \neg \diamond M1)$ // after M2, M1 is not allowed





DSPL with temporal properties



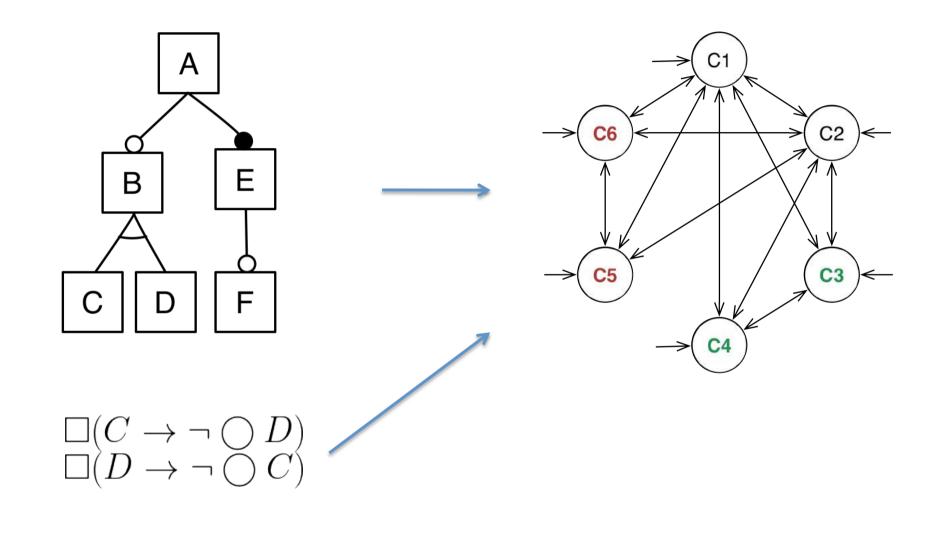
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DSPL with temporal properties



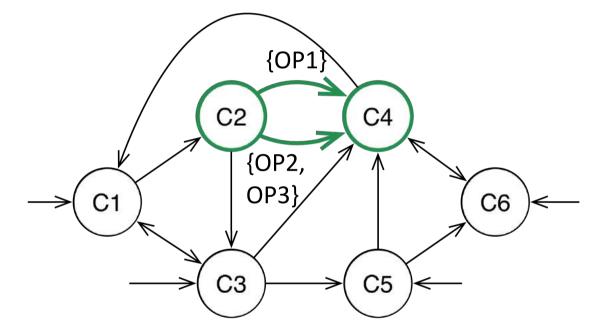
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Reconfiguration operations

• Doubly labeled transition systems^[22]



[22] M. H. ter Beek et al., "An Action/State-Based Model-Checking Approach for the Analysis of Communication Protocols for Service-Oriented Applications," in Proc. 12th Int. Workshop Formal Methods for Industrial Critical Systems (FMICS'07), Berlin, Germany, Jul. 2008, pp. 133–148.

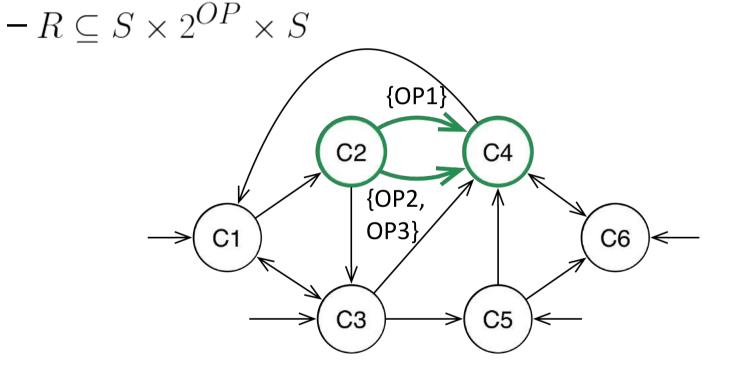
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Reconfiguration operations

- Doubly labeled transition systems^[22] -TS = (S, I, OP, R, AP, L)
 - $\ OP$ is the set of reconfiguration operations in the DSPL



[22] M. H. ter Beek et al., "An Action/State-Based Model-Checking Approach for the Analysis of Communication Protocols for Service-Oriented Applications," in Proc. 12th Int. Workshop Formal Methods for Industrial Critical Systems (FMICS'07), Berlin, Germany, Jul. 2008, pp. 133–148.

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State/Event LTL

- SE-LTL can express temporal expressions over state and transition labels^[23]
 - Can combine reconfiguration operations and features in temporal constraints

[23] S. Chaki et al. "State/Event-Based Software Model Checking," in Proc. 4th Int. Conf. Integrated Formal Methods (IFM'04), Canterbury, UK, Apr. 2004, pp. 128–147.

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Reconfiguration operations

 $OP = \{ActivateF\}$ C1 Α $C1 = \{A, E\}$ C2 **C6** $C2 = \{A, E, F\}$ $C3 = \{A, B, C, E\}$ Ε В $C4 = \{A, B, C, E, F\}$ $C5 = \{A, B, D, E\}$ C3 C5 $C6 = \{A, B, D, E, F\}$ F С $\Box((\neg F \land \bigcirc F) \leftrightarrow ActivateF)$ Université de Lille CRIStAL SPIRALS SEAMS 2017 – Buenos Aires, May 23, 2017

Examples

- Cannot downgrade MySQL plan
 - [](M2 -> !<>M1)
 - [](M3 -> !<>(M1 | M2))
 - [](M4 -> !<>(M1 | M2 | M3))
 - [](Change(ClearDBMySQL) -> UpgradeClearDB)
- Upgrade PostgreSQL
 - [](Change(HerokuPostgres) & (H1 | H2) -> PGCopy)
 - [](Change(HerokuPostgres) -> (PGCopy | FollowerChangeover))
- Change Location
 - [](Change(Location) -> MigrateApp)

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Problem statement

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 - e.g. reduced downtime or costs

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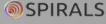


Reasoning

- Reconfiguration request
 - Features to be included/excluded







Reasoning

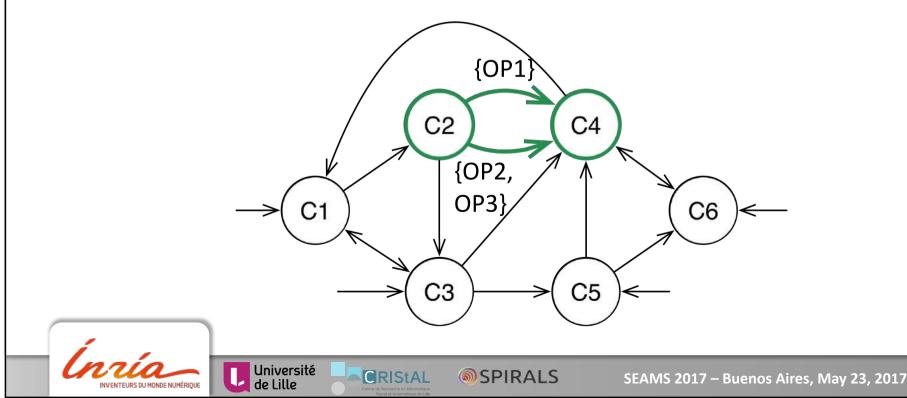
- Reconfiguration request
 - Features to be included/excluded
- Cost-based constraints
 - Reconfiguration time, downtime, financial cost, etc





Reasoning

- Reconfiguration query: $Q = (A, E, \phi)$
 - -A: features to include
 - -E: features to exclude
 - $-\phi$: constraint over costs
- Example query: $q = (\{C\}, \{D\}, \text{downtime} = 0)$



Symbolic Representation

- Building the transition system for a feature model can be unfeasible
 - State-explosion problem
- Represent a transition system as a propositional formula
 - Use SAT solver to solve reconfiguration queries



Symbolic Representation

 Feature models^[25] and SE-LTL expressions^[27] can be represented as propositional formulas



- \widetilde{M} and \widetilde{M}' represent the set of possible source and target states (configurations of the feature model M)
- \widetilde{x} is the conjunction of LTL expressions

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– \widetilde{s} represents the current state

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– \widetilde{q} represents the reconfiguration query (pseudo-boolean encoding)

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[25] D. Batory, "Feature Models, Grammars, and Propositional Formulas," in Proc. 9th Int. Conf. Software Product Lines (SPLC'05), Rennes, France, Sep. 2005, pp. 7–20.

[27] A. Cimatti, M. Pistore, M. Roveri, and R. Sebastiani, "Improving the Encoding of LTL Model Checking into SAT," in Proc. 3rd Int. Workshop Model Checking and Abstract Interpretation (VMCAI'02), Venice, Italy, Jan. 2002, pp. 196–207.



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- Case study on Heroku PaaS
 - feasibility for modeling reconfiguration constraints
 - performance of reasoning





- Case study on Heroku PaaS
- Feature Model extracted from documentation
 - 7 available regions, 11 programming frameworks, 6 container sizes
 - reconfiguration constraints
 - 161 addon services (data storage, networking, security, ...)
 - 1036 features, 134 cross-tree constraints, 124 temporal constraints





- Case study on Heroku PaaS
- Feature Model extracted from documentation
- Simulate context changes
 - 4 adaptation scenarios
 - 5 reconfiguration queries

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- 3 utilization profiles
- 12 executions



- Case study on Heroku PaaS
- Feature Model extracted from documentation
- Simulate context changes
- Adaptation scenarios
 - Change in database size requires new database plan
 - Request for a new feature not available in current region

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Change in programming framework and database

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Scaling up application container

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- Case study on Heroku PaaS
- Feature Model extracted from documentation
- Simulate context changes
- Adaptation scenarios
- Reconfiguration queries
 - No constraints
 - Constraints over price
 - Constraints over downtime/price
 - Optimize on price
 - Optimize on downtime/price

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- Case study on Heroku PaaS
- Feature Model extracted from documentation
- Simulate context changes
- Adaptation scenarios
- Reconfiguration queries
- Application utilization profiles

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- Database size, application size, startup time, etc...



- Case study on Heroku PaaS
- Feature Model extracted from documentation
- Simulate context changes
- Adaptation scenarios
- Reconfiguration queries
- Application utilization profiles

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- DBSize: 10GB, AppSize: 100 MB, AppStartUp: 15
- DBSize: 100GB, AppSize: 200 MB, AppStartUp: 30s

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DBSize: 2TB, AppSize: 500 MB, AppStartUp: 60s

Process step	Execution time (ms)				1
	Avg	StdDev	Min	Max	#Exec
Build Trans System	8777.31	303.71	8262	10308	720
- Process FM	244.75	28.57	191	552	
- Process LTL	8533.57	291.81	8025	10023	
All Queries	183.34	50.20	118	389	720
- Build	83.05	50.69	27	227	
- Solve	100.29	37.13	5	200	
Q wo/ Constraints	140.97	12.93	118	198	144
- Build	32.73	4.56	27	48	
- Solve	108.24	11.58	90	153	
Q w/ Constraints	224.23	55.55	128	389	288
- Build	140.41	27.65	80	227	211 kg (2122
- Solve	83.82	53.13	5	200	
Q w/ Optimization	163.63	13.26	136	230	288
- Build	50.85	7.11	38	77	
- Solve	112.77	10.21	94	166	

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- Temporal constraints enhance modeling of DSPLs
 - Compact notation for constraints over transitions
 - Support for reasoning over reconfiguration operations
- Performance is acceptable in the cloud context
 Implementation can be improved
- Threats to validity

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Case study is not exhaustive and considers only cloud computing



Conclusion & Perspectives

- Temporal constraints in DSPL
 - Better modeling of adaptive behavior
 - Reasoning over adaptation alternatives





Conclusion & Perspectives

- Temporal constraints in DSPL
 - Better modeling of adaptive behavior
 - Reasoning over adaptation alternatives
- Cardinality-based feature models

• Multi-cloud environment adaptation



Questions

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More information

http://researchers.lille.inria.fr/sousa/seams2017/





